

Accurate Solution of Structured Linear Systems and Least Squares Problems through Rank Revealing Decompositions

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Joint work with
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Outline

- 1 High relative accuracy and structured matrices
- 2 RRD and accurate computations
- 3 Algorithms and error analysis
- 4 Perturbation theory
- 5 Numerical Experiments

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Problem

Given a **structured and potentially bad conditioned** (Band, Toeplitz, Hankel, Vandermonde, Cauchy, Quasiseparable, Graded, etc ...) matrix A , is it possible to solve

- a linear system

$$Ax = b, \quad A \in \mathbb{C}^{n \times n}, \exists A^{-1}$$

- or, to compute the minimal length solution of a least squares problem

$$Ax \approx b, \quad A \in \mathbb{C}^{m \times n}, \quad \text{rank}(A) = r \leq n \leq m$$

with **high relative accuracy (hra)**?

High relative accuracy

We say that the solution x of a problem is computed with **high relative accuracy (hra)** if

$$\frac{\|\hat{x} - x\|}{\|x\|} = O(\mathbf{u})$$

with \hat{x} the computed solution and \mathbf{u} the unit roundoff ($\mathbf{u} = 1.11 \cdot 10^{-16}$), independently of the condition number of A .

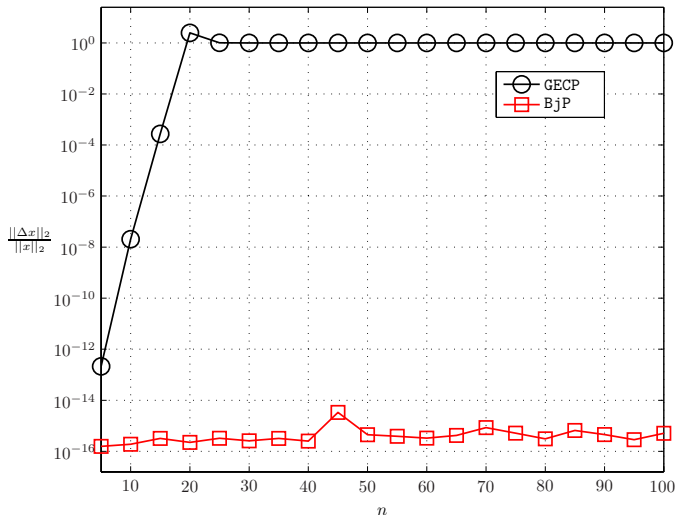
TP Vandermonde matrices

$O(n^2)$ algorithm: [1970, Björk&Pereyra; 1987, Higham]

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Forward error vs size. TP Vandermonde matrices.

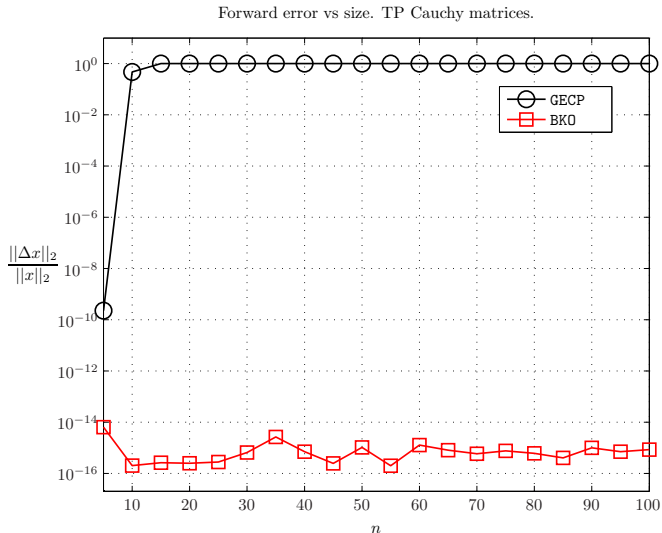


TP Cauchy matrices

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New methods

The methods we are going to present here

- Solve with **high relative accuracy** linear systems $Ax = b$ and least squares problems $Ax \approx b$,
- for “**almost any**” rhs b ,
- it applies to matrices for which there is a **Accurate Rank Revealing Decomposition**.

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- Solve with **high relative accuracy** linear systems $Ax = b$ and least squares problems $Ax \approx b$,
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- it applies to matrices for which there is a **Accurate Rank Revealing Decomposition**.

Accurate Rank Revealing Decompositions are known to exist for many classes of structured matrices (in particular, Cauchy, Vandermonde, Graded).

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Rank Revealing Decomposition (RRD)

- RRD

$$A = XDY = \begin{bmatrix} \times & \times \\ \times & \times \\ \times & \times \\ \times & \times \\ \times & \times \end{bmatrix} \begin{bmatrix} \times & \\ & \times \end{bmatrix} \begin{bmatrix} \times & \times & \times \\ \times & \times & \times \end{bmatrix}$$

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- $A \in \mathbb{C}^{m \times n}$, $\text{rank}(A) = r$,

$$X \in \mathbb{C}^{m \times r}, \quad Y \in \mathbb{C}^{r \times n}, \quad D \in \mathbb{C}^{r \times r} \text{ diagonal}$$

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- $\kappa(X), \kappa(Y) \gtrsim 1, \quad \kappa(D) \approx \kappa(A)$

Structured matrices and accurate RRDs

- RRDs can be computed in practice from Gaussian Elimination with Complete Pivoting (GECP) (i.e., direct methods).

Structured matrices and accurate RRDs

- **RRDs** can be computed in practice from Gaussian Elimination with Complete Pivoting (GECP) (i.e., direct methods).
- But **accurate RRDs** can only be computed for matrices with special structures:
 - Scaled-Cauchy, Vandermonde (DFT + GECP). [Demmel]
 - Diagonally Dominant M-Matrices. [Demmel and Koev]
 - Polynomial Vandermonde. [Demmel and Koev]
 - Well Scaled Positive Definite. [Demmel and Veselić]
 - Acyclic Matrices (include bidiagonal). [Demmel and Gragg]
 - Diagonally Dominant. [Qiang Ye, Dopico and Koev]
 - DSTU. [Demmel]
 - Graded Matrices. [Dopico, M, Ceballos]

Rank Revealing Decomposition

Theorem (Demmel et al. 1999)

An SVD with high relative accuracy can be computed for a matrix A if and only if an accurate RRD can be computed.

Rank Revealing Decomposition

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An SVD with high relative accuracy can be computed for a matrix A if and only if an accurate RRD can be computed.

We say that the computed factors \hat{X} , \hat{Y} and \hat{D} , produce an **accurate RRD** of $A = XDY$ if they obey

$$\frac{\|\hat{X} - X\|}{\|X\|} = O(\mathbf{u}), \frac{\|\hat{Y} - Y\|}{\|Y\|} = O(\mathbf{u}) \quad \text{and} \quad \frac{|\hat{D} - D|}{|D|} = O(\mathbf{u})$$

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- **Linear systems of equations**
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- **Least squares problems**
[2011, Castro&Ceballos&Dopico&M]

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Algorithm: Linear Systems

- **Input:** $A \in \mathbb{R}^{n \times n}$, $b \in \mathbb{R}^n$
 - **Output:** x solution of $Ax = b$
- ① Compute an accurate RRD of $A = XDY$
 - ② Solve the three systems

$$Xs = b \quad \longrightarrow \quad s$$

$$Dw = s \quad \longrightarrow \quad w$$

$$Yx = w \quad \longrightarrow \quad x$$

$O(n^3)$ complexity

Algorithm: Least squares Problems

- **Input:** $A \in \mathbb{C}^{m \times n}$, $b \in \mathbb{C}^m$
 - **Output:** $x_0 = A^\dagger b$
- 1 Compute an accurate RRD of $A = XDY$
 - 2 Compute the solution x_1 of $\min_{x \in \mathbb{C}^n} \|b - Xx\|_2$ using QR Householder.
 - 3 Solve, for x_2 , the diagonal linear system,
 $Dx_2 = x_1$.
 - 4 Compute the minimal 2-norm solution x_0 of $Yx = x_2$ using QR Householder on Y^* .

$O(n^3)$ complexity

Backward Error

Theorem

The computed minimal 2-norm solution of $XDYx \approx b$, with roundoff error \mathbf{u} , *is the exact* minimal 2-norm solution of

$$(X + \Delta X)(D + \Delta D)(Y + \Delta Y)\hat{x} \approx b$$

with

$$\|\Delta X\| \leq O(\mathbf{u})\|X\|, \quad \|\Delta Y\| \leq O(\mathbf{u})\|Y\|, \quad |\Delta D| \leq O(\mathbf{u})|D|$$

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$$\begin{aligned}\tilde{A} &= (X + \Delta X)(D + \Delta D)(Y + \Delta Y) \\ &= (I + \Delta X X^\dagger)X(I + \Delta D D^{-1})DY(I + Y^\dagger \Delta Y)\end{aligned}$$

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where $E = \Delta X X^\dagger + X \Delta D D^{-1} X^\dagger + \Delta X \Delta D D^{-1} X^\dagger$ and $F = Y^\dagger \Delta Y$.

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where $\|E\| \leq O(\mathbf{u})\kappa(X)$ and $\|F\| \leq O(\mathbf{u})\kappa(Y)$, if $\Delta X, \Delta Y$ come from previous error analysis.

Moore-Penrose Pseudoinverse Multiplicative Perturbation Theory

Theorem

Let $A \in \mathbb{C}^{m \times n}$ and $\tilde{A} = (I + E)A(I + F)$ where $I + E$ and $I + F$ are nonsingular.

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$$\tilde{A}^\dagger = (I + F)^{-1}A^\dagger(I + E)^{-1}$$

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Theorem

Let $A \in \mathbb{C}^{m \times n}$ and $\tilde{A} = (I + E)A(I + F)$ where $I + E$ and $I + F$ are nonsingular. Then

$$\tilde{A}^\dagger = P_{\tilde{A}^*} (I + F)^{-1} A^\dagger (I + E)^{-1} P_{\tilde{A}}$$

where $P_{\tilde{A}} = \tilde{A}\tilde{A}^\dagger$ and $P_{\tilde{A}^*} = \tilde{A}^\dagger\tilde{A}$.

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where $P_{\tilde{A}} = \tilde{A}\tilde{A}^\dagger$ and $P_{\tilde{A}^*} = \tilde{A}^\dagger\tilde{A}$.

$$\tilde{A}^\dagger = \left(I + (I - P_{\tilde{A}^*})F^* - P_{\tilde{A}^*}\hat{F} \right) A^\dagger \left(I + E^*(I - P_{\tilde{A}}) - \hat{E}P_{\tilde{A}} \right)$$

where $\hat{E} = (I + E)^{-1}E$ and $\hat{F} = (I + F)^{-1}F$.

RRD Perturbation Theory

Theorem

Let

- x_0 be the minimum 2-norm solution of

$$\min_{x \in \mathbb{C}^n} \|b - XDY x\|_2,$$

- \tilde{x}_0 be the minimum 2-norm solution of

$$\min_{x \in \mathbb{C}^n} \|b + h - (X + \Delta X)(D + \Delta D)(Y + \Delta Y) x\|_2$$

$\|\Delta X\|_2 \leq \epsilon \|X\|_2$, $\|\Delta Y\|_2 \leq \epsilon \|Y\|_2$, $|\Delta D| \leq \epsilon |D|$, and $\|h\|_2 \leq \epsilon \|b\|_2$.

Then:

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Then:

$$\frac{\|\tilde{x}_0 - x_0\|_2}{\|x_0\|_2} \leq \epsilon \left(P_1 \kappa_2(Y) + P_2 \kappa_2(X) \frac{\|A^\dagger\|_2 \|b\|_2}{\|x_0\|_2} \right) + O(\epsilon^2).$$

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RRD Perturbation Theory

$$\begin{aligned} \kappa_f(X, D, Y, b) &:= \\ &:= \lim_{\epsilon \rightarrow 0} \sup \left\{ \frac{\|\tilde{x} - x\|}{\epsilon \|x\|} : (X + \Delta X)(D + \Delta D)(Y + \Delta Y)\tilde{x} = b + h, \right. \\ &\quad \left. \|h\| \leq \epsilon \|b\|, \|\Delta X\| \leq \epsilon \|X\|, |\Delta D| \leq \epsilon |D|, \|\Delta Y\| \leq \epsilon \|Y\| \right\} \end{aligned}$$

Theorem

$$\kappa_f(X, D, Y, b) \leq \left(\kappa(Y) + [1 + 2 \kappa(X)] \frac{\|A^{-1}\| \|b\|}{\|x\|} \right) \leq 3 \kappa_f(X, D, Y, b)$$

Perturbing the right-hand side

The minimal length solutions of the least squares problems

$$Ax \approx b \quad \text{and} \quad A(x + \Delta x) \approx b + \Delta b$$

are related as

$$\frac{\|\Delta x_0\|}{\|x_0\|} \leq \frac{\|A^\dagger\| \|\Delta b\|}{\|x_0\|} = \kappa(A, b) \frac{\|\Delta b\|}{\|b\|}$$

with

$$\kappa(A, b) \equiv \frac{\|A^\dagger\| \|b\|}{\|x_0\|} \leq \|A^\dagger\| \|A\| \frac{\|b\|}{\|P_A(b)\|} = \frac{\kappa(A)}{\cos \theta}.$$

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Example in MATLAB:

```
>> m=30;n=15;v= vander(randn(m,1)); v=fliplr(v);  
>> v=v(:,1:n);b=randn(m,1);  
>> cond(v)= 1.6234e+07  
>> cond(v)*norm(b)/norm(v*x) = 2.3898e+07  
>> norm(x) = 230.23  
>> norm(b) = 5.1154  
>> norm(pinv(v))*norm(b)/norm(x) = 4.181
```


Perturbing the right-hand side

If $A \in \mathbb{C}^{n \times n}$ is invertible and

$$Ax = b \quad \text{and} \quad A(x + \Delta x) = b + \Delta b$$

then

$$\frac{\|\Delta x\|}{\|x\|} \leq \frac{\|A^{-1}\| \|\Delta b\|}{\|x\|} = \kappa(A, b) \frac{\|\Delta b\|}{\|b\|}$$

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with

$$\kappa(A, b) \equiv \frac{\|A^{-1}\| \|b\|}{\|x\|} \leq \|A^{-1}\| \|A\| = \kappa(A).$$

Example in MATLAB:

```
>> v= vander(randn(20,1)); b=randn(20,1);
```

```
>> cond(v)= 7.1021e+11
```

```
>> norm(inv(v)) = 1.4857e+11
```

```
>> norm(x) = 5.4507e+10
```

```
>> norm(b) = 3.0934
```

```
>> norm(inv(v))*norm(b)/norm(x) = 8.4317
```

Effective conditioning

Theorem (Chan & Foulser, 1988)

If $A \in \mathbb{C}^{m \times n}$, $\text{rank}(A) = r \leq n \leq m$

$$A = U\Sigma V^H = [\tilde{U} \ U_k] \begin{bmatrix} \tilde{\Sigma} & 0 \\ 0 & \Sigma_k \end{bmatrix} \begin{bmatrix} \tilde{V} \\ V_k \end{bmatrix}$$

is a (partitioned) reduced SVD of A , and $x = A^\dagger b$, then

$$\kappa_2(A, b) = \frac{\|b\|_2}{\sigma_r \|x\|_2} \leq \frac{\sigma_{r+1-k}}{\sigma_r} \frac{\|b\|_2}{\|P_k b\|_2}, \quad \forall k = 1 : r$$

where $P_k = U_k U_k^*$ is the orthogonal projector onto the subspace spanned by the last k left singular vectors of A .

Effective conditioning

For the perturbation problem $A(x + \Delta x) \approx b + \Delta b$ this means

$$\frac{\|\Delta x_0\|_2}{\|x_0\|_2} \leq \kappa_2(A, b) \frac{\|\Delta b\|_2}{\|b\|_2} \leq \kappa_{eff}(A, b) \frac{\|\Delta b\|_2}{\|b\|_2}$$

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$$\frac{\|\Delta x_0\|_2}{\|x_0\|_2} \leq \kappa_2(A, b) \frac{\|\Delta b\|_2}{\|b\|_2} \leq \kappa_{eff}(A, b) \frac{\|\Delta b\|_2}{\|b\|_2}$$

with

$$\kappa_{eff}(A, b) = \min_{k=1:r} \left\{ \frac{\sigma_{r+1-k}}{\sigma_r} \frac{\|b\|_2}{\|P_k b\|_2} \right\} \leq \frac{\|b\|_2}{\|P_1 b\|_2}$$

Forward Error

Theorem

If \hat{x}_0 is the computed solution of $Ax \approx b$ using previous algorithm, with roundoff error \mathbf{u} , then,

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$$\frac{\|\hat{x}_0 - x_0\|}{\|x_0\|} \leq \mathbf{u} \left(P_1 \kappa_2(\hat{Y}) + P_2 \kappa_2(\hat{X}) \frac{\|A^\dagger\| \|b\|}{\|x_0\|} \right) + O(\mathbf{u}^2)$$

Forward Error

Theorem

If \hat{x}_0 is the computed solution of $Ax \approx b$ using previous algorithm, with roundoff error \mathbf{u} , then, for $k = 1 : r$,

$$\frac{\|\hat{x}_0 - x_0\|}{\|x_0\|} \leq \mathbf{u} \left(P_1 \kappa_2(\hat{Y}) + P_2 \kappa_2(\hat{X}) \frac{\sigma_{r+1-k}}{\sigma_r} \frac{\|b\|}{\|P_k b\|} \right) + O(\mathbf{u}^2)$$

Forward Error

Theorem

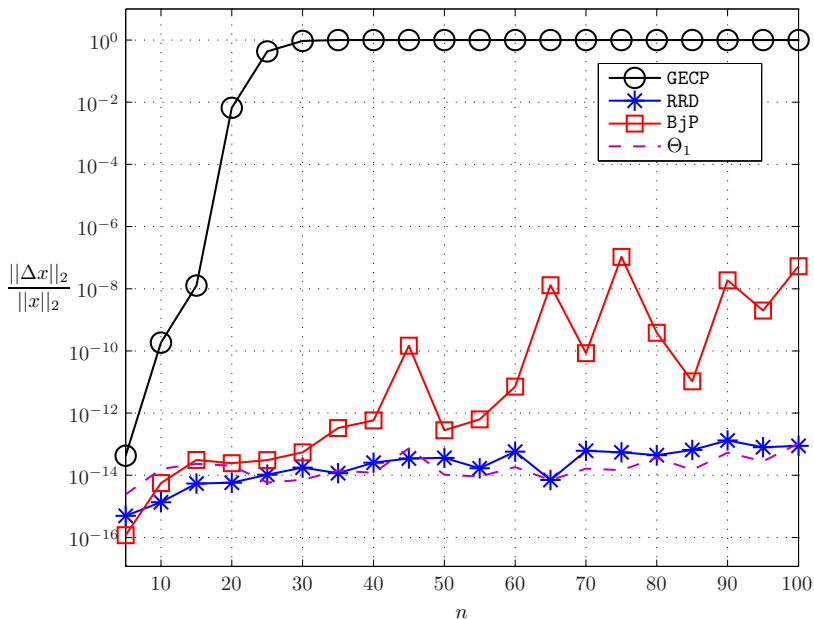
If \hat{x}_0 is the computed solution of $Ax \approx b$ using previous algorithm, with roundoff error \mathbf{u} , then,

$$\frac{\|\hat{x}_0 - x_0\|}{\|x_0\|} \leq \mathbf{u} \left(P_1 \kappa_2(\hat{Y}) + P_2 \kappa_2(\hat{X}) \frac{\|b\|}{\|P_1 b\|} \right) + O(\mathbf{u}^2)$$

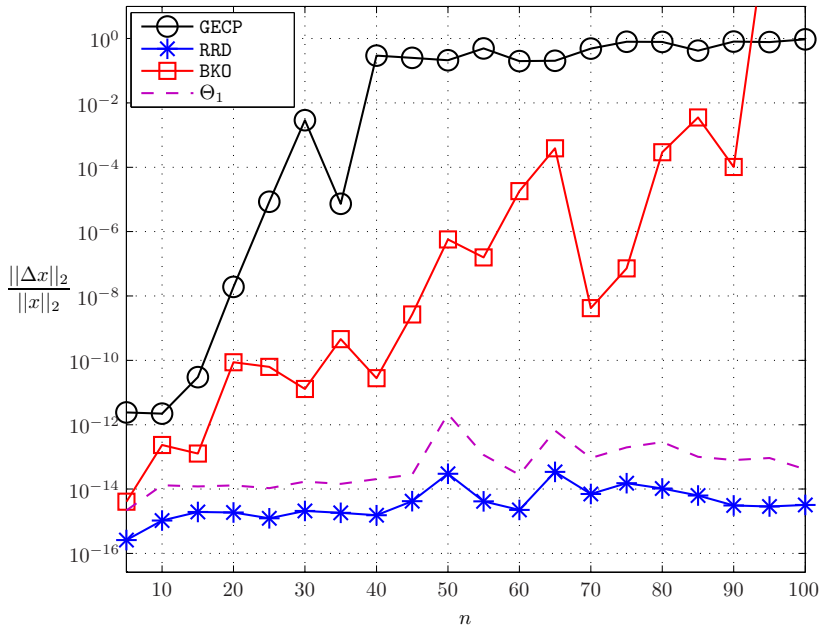
Outline

- 1 High relative accuracy and structured matrices
- 2 RRD and accurate computations
- 3 Algorithms and error analysis
- 4 Perturbation theory
- 5 Numerical Experiments**

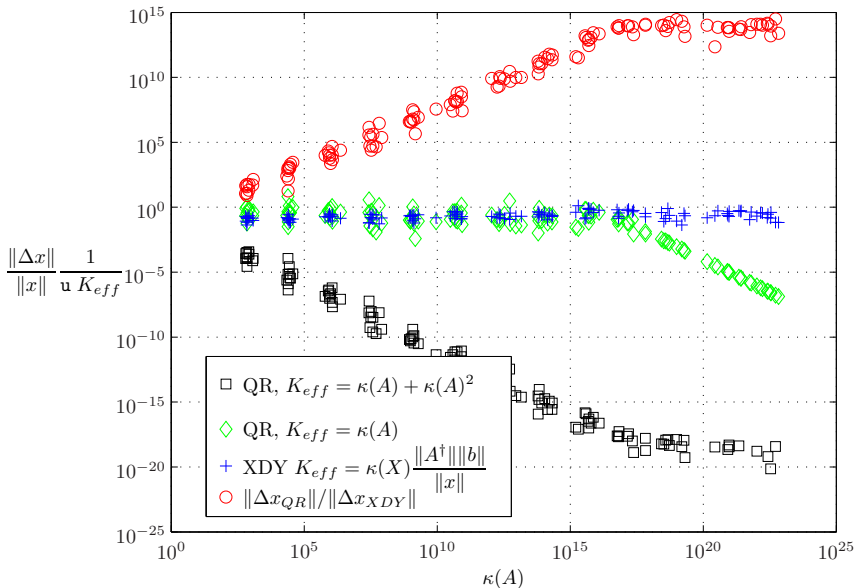
Forward error vs size. Random Vandermonde matrices.



Forward error vs size. Random Cauchy matrices.



Vandermonde Matrices. $Ax \approx b$, $m = 50, n = 5 : 2 : 30$.



Conclusions

- It is possible now to compute **accurately** solutions of **structured** systems with huge condition numbers.
- It is necessary to take advantage of the structure of the matrix.
- New **multiplicative perturbation** results are necessary.
- **RRD** play a central and important role.

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Thanks for your attention