

Strong linearizations of rational matrices with polynomial part expressed in an orthogonal basis

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Abstract

We construct a new family of strong linearizations of rational matrices considering the polynomial part of them expressed in a basis that satisfies a three term recurrence relation. For this purpose, we combine the theory developed by Amparan et al., *SIAM J. Matrix Anal. Appl.*, 39(4) (2018), and the new linearizations of polynomial matrices introduced by Faßbender and Saltenberger, *Linear Algebra Appl.*, 525 (2017). In addition, we present a detailed study of how to recover eigenvectors of a rational matrix from those of its linearizations in this family. We complete the paper by discussing how to extend the results when the polynomial part is expressed in other bases, and by presenting strong linearizations that preserve the structure of symmetric or Hermitian rational matrices. A conclusion of this work is that the combination of the results in this paper with those in Amparan et al., *SIAM J. Matrix Anal. Appl.*, 39(4) (2018), allows us to use essentially all the strong linearizations of polynomial matrices developed in the last fifteen years to construct strong linearizations of any rational matrix by expressing such a matrix in terms of its polynomial and strictly proper parts.

Keywords: rational matrix, rational eigenvalue problem, strong block minimal bases pencil, strong linearization, recovery of eigenvectors, symmetric strong linearization, Hermitian strong linearization

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1. Introduction

In recent years, the interest in solving the rational eigenvalue problem (REP) has grown as it arises in many applications, either directly or as an approximation of other nonlinear eigenvalue problems, see [9, 17, 18, 24, 27, 30]. There are several algorithms for its numerical resolution and, since the appearance of [30], using linearizations is one of the most competitive methods for solving REPs nowadays [12, 31]. This has led to a rigorous development of the theory of linearizations for rational matrices.

There are two different approaches in order to give a notion of linearization of a rational matrix. On the one hand, Alam and Behera give in [1] a definition based on the fact that any rational matrix $G(\lambda)$ admits a right coprime matrix fraction description $G(\lambda) = N(\lambda)D(\lambda)^{-1}$, where $N(\lambda)$ and $D(\lambda)$ are polynomial matrices. These linearizations preserve the finite pole and zero structure of the original matrix. In contrast, Amparan et al. give in [4] a new notion of linearization that not only preserves the finite but also the infinite structure of poles and zeros. These linearizations are called strong linearizations in [4]. This definition and other notions about rational matrices will be reviewed in Section 2.

Throughout this work the fundamental fact is that any rational matrix $G(\lambda)$ can be uniquely written as $G(\lambda) = D(\lambda) + G_{sp}(\lambda)$ where $D(\lambda)$ is a polynomial matrix, called the polynomial part of $G(\lambda)$, and $G_{sp}(\lambda)$ is a strictly proper rational matrix, called the strictly proper part of $G(\lambda)$. Thanks to this property, infinitely many strong linearizations of rational matrices are constructed in [4] considering “strong block minimal bases pencils” associated to their polynomial parts, see [13]. These strong block minimal bases pencils are strong linearizations of the polynomial part and according to [8] include, modulo permutations, all the Fiedler-like linearizations of the polynomial part. Although strong block minimal bases pencils are one of the most important classes of strong linearizations of polynomial matrices, the question whether or not other strong linearizations of rational matrices can be constructed based on other kinds of strong linearizations of the polynomial part arises naturally.

For answering the question posed in the previous paragraph, we construct in this paper strong linearizations of a rational matrix by using strong linearizations of its polynomial part $D(\lambda)$ that belong to the other important family of strong linearizations of polynomial matrices (which are not strong block minimal bases pencils in general), i.e., the so-called vector spaces of linearizations, originally introduced in [26], further studied in [10, 28], and recently extended in [15]. In particular, we consider in this paper strong linearizations of $D(\lambda)$ that belong to the ansatz spaces $\mathbb{M}_1(D)$ or $\mathbb{M}_2(D)$, developed by Faßbender and Saltenberger in [15]. Therefore, the results in this paper are of interest when rational matrices with nontrivial polynomial part are considered, that is, rational matrices in which the polynomial part has degree greater than or equal to two.

As a consequence of the discussion above, we emphasize the following main conclusion of this work: the combination of the results in this paper and those in [4] allows us to construct very easily infinitely many strong linearizations of rational matrices via the following three-step strategy: (1) express the rational matrix as the sum of its polynomial and strictly proper parts; (2) construct *any* of the strong linearizations of the polynomial part known so far; and (3) combine adequately that strong linearization with a minimal state-space realization of the strictly proper part.

Next, another motivation of the results in this paper is discussed. In order to compute the eigenvalues of polynomial matrices from linearizations, the work [23] shows that, for

polynomial matrices of large degree, the use of the monomial basis to express the matrix leads to numerical instabilities. According to the algorithms in [12, 30, 31], it is expected that this instability appears also while computing eigenvalues of REPs when the polynomial part of the rational matrix has large degree and is expressed in terms of the monomial basis. For that reason, it is of interest to consider rational matrices with polynomial parts expressed in other bases as the Chebyshev basis. In particular, in Sections 3 and 4, we construct strong linearizations of rational matrices with polynomial parts expressed in terms of a basis that satisfies a three term recurrence relation. In addition, in Section 9, we briefly discuss how to construct strong linearizations when the polynomial part is expressed in other bases. We emphasize that the construction of these new strong linearizations is a consequence of the theory of strong linearizations developed in [4] together with Lemma 2.7. More precisely, given a strong linearization of a rational matrix, Lemma 2.7 allows us to obtain infinitely many strong linearizations of the rational matrix by using strict equivalence with a certain structure.

The rest of this paper is organized as follows. In Section 5, we show how to recover the eigenvectors of the rational matrix from those of its strong linearizations constructed in Sections 3 and 4. Moreover, given a symmetric rational matrix, in Section 7 we construct strong linearizations that preserve its symmetric structure by using symmetric realizations of the strictly proper part, which are introduced in Section 6, and strong linearizations in the double ansatz space $\mathbb{DM}(D)$ [15] of the polynomial part. In Section 8, we present analogous results for Hermitian rational matrices. Finally, Section 10 is reserved for discussing the conclusions and lines of future work.

2. Preliminaries

$\mathbb{F}[\lambda]$ denotes the ring of polynomials with coefficients in an arbitrary field \mathbb{F} , and $\mathbb{F}(\lambda)$ the field of *rational functions*, i.e., the field of fractions of $\mathbb{F}[\lambda]$. $\mathbb{F}(\lambda)^{p \times m}$, $\mathbb{F}[\lambda]^{p \times m}$ and $\mathbb{F}^{p \times m}$ denote the sets of $p \times m$ matrices with elements in $\mathbb{F}(\lambda)$, $\mathbb{F}[\lambda]$ and \mathbb{F} , respectively. The elements of $\mathbb{F}(\lambda)^{p \times m}$ and $\mathbb{F}[\lambda]^{p \times m}$ are called rational and polynomial matrices, respectively. A polynomial matrix $P(\lambda) = \sum_{i=0}^k \lambda^i P_i$ with $P_i \in \mathbb{F}^{p \times m}$ is said to have *degree* k if $P_k \neq 0$. If $k = 1$ or $k = 0$ then $P(\lambda)$ is said to be a *pencil*. Matrices in $\mathbb{F}[\lambda]^{m \times m}$ with nonzero constant determinant are said to be *unimodular*. Two rational matrices $Q(\lambda), R(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ are said to be *unimodularly equivalent* if there exist unimodular matrices $U(\lambda) \in \mathbb{F}[\lambda]^{p \times p}$ and $V(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ such that $U(\lambda)Q(\lambda)V(\lambda) = R(\lambda)$. Moreover, $Q(\lambda)$ and $R(\lambda)$ are said to be *strictly equivalent* if $UQ(\lambda)V = R(\lambda)$ with $U \in \mathbb{F}^{p \times p}$ and $V \in \mathbb{F}^{m \times m}$ invertible matrices.

A rational matrix $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ is said to be *regular* or *nonsingular* if $p = m$ and its determinant, $\det G(\lambda)$, is not identically equal to zero. Otherwise, $G(\lambda)$ is said to be *singular*. Given any rational matrix $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$, the *(finite) eigenvalues of $G(\lambda)$* are defined as the scalars $\lambda_0 \in \overline{\mathbb{F}}$ (the algebraic closure of \mathbb{F}) such that $G(\lambda_0) \in \overline{\mathbb{F}}^{p \times m}$ and $\text{rank } G(\lambda_0) < \max_{\mu \in \overline{\mathbb{F}}} \text{rank } G(\mu)$. The *rational eigenvalue problem* (REP) consists of finding the eigenvalues of $G(\lambda)$. If $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ is regular, which is the most common case in applications of REPs, the REP is equivalent to the problem of finding scalars $\lambda_0 \in \overline{\mathbb{F}}$ such that there exist nonzero constant vectors $x \in \overline{\mathbb{F}}^{m \times 1}$ and $y \in \overline{\mathbb{F}}^{m \times 1}$ satisfying

$$G(\lambda_0)x = 0 \quad \text{and} \quad y^T G(\lambda_0) = 0,$$

respectively. The vectors x are called *right eigenvectors associated to* λ_0 , and the vectors y *left eigenvectors*. Although it is not common in the literature, if $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ is singular, we call in this paper right and left eigenvectors of $G(\lambda)$ associated to an eigenvalue λ_0 to any nonzero vectors $x \in \overline{\mathbb{F}}^{m \times 1}$ and $y \in \overline{\mathbb{F}}^{p \times 1}$ satisfying $G(\lambda_0)x = 0$ and $y^T G(\lambda_0) = 0$, respectively.

The *finite poles* and *zeros* of a rational matrix $G(\lambda)$ are the roots in $\overline{\mathbb{F}}$ of the polynomials that appear on the denominators and numerators, respectively, in its (*finite*) *Smith–McMillan form* (see [4, 29, 32]). Then the *finite eigenvalues* of $G(\lambda)$ are the finite zeros that are not poles.

For solving the REP, and many other problems on rational matrices, it is useful to consider the fact that any rational matrix $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ can be written as

$$G(\lambda) = D(\lambda) + C(\lambda)A(\lambda)^{-1}B(\lambda) \quad (1)$$

for some nonsingular polynomial matrix $A(\lambda) \in \mathbb{F}[\lambda]^{n \times n}$ and polynomial matrices $B(\lambda) \in \mathbb{F}[\lambda]^{n \times m}$, $C(\lambda) \in \mathbb{F}[\lambda]^{p \times n}$ and $D(\lambda) \in \mathbb{F}[\lambda]^{p \times m}$ (see [29]). The polynomial matrix

$$P(\lambda) = \begin{bmatrix} A(\lambda) & B(\lambda) \\ -C(\lambda) & D(\lambda) \end{bmatrix} \quad (2)$$

is called a *polynomial system matrix* of $G(\lambda)$, i.e., $G(\lambda)$ is the Schur complement of $A(\lambda)$ in $P(\lambda)$. Then $G(\lambda)$ is called the *transfer function matrix* of $P(\lambda)$, and $\deg(\det A(\lambda))$ the *order* of $P(\lambda)$, where $\deg(\cdot)$ stands for degree. Moreover, $P(\lambda)$ is said to have *least order*, or to be *minimal*, if its order is the smallest integer for which polynomial matrices $A(\lambda)$ (nonsingular), $B(\lambda)$, $C(\lambda)$ and $D(\lambda)$ satisfying (1) exist. The least order is uniquely determined by $G(\lambda)$ and is denoted by $\nu(G(\lambda))$. It is also called the *least order* of $G(\lambda)$ ([29, Chapter 3, Section 5.1] or [32, Section 1.10]). From [29, Chapter 3, Theorem 4.1], it can be deduced that $\nu(G(\lambda))$ is the degree of the polynomial that results by making the product of the denominators in the (*finite*) Smith–McMillan form of $G(\lambda)$. The regular REP $G(\lambda)x = 0$ is related to the polynomial eigenvalue problem (PEP) $P(\lambda)z = 0$ as is shown in [5, Proposition 3.1].

A rational function $r(\lambda) = \frac{n(\lambda)}{d(\lambda)}$ is said to be *proper* if $\deg(n(\lambda)) \leq \deg(d(\lambda))$, and *strictly proper* if $\deg(n(\lambda)) < \deg(d(\lambda))$. Let us denote $\mathbb{F}_{pr}(\lambda)$ the ring of proper rational functions. Its units are called *biproper rational functions*, i.e., rational functions having the same degree of numerator and denominator. $\mathbb{F}_{pr}(\lambda)^{p \times m}$ denotes the set of $p \times m$ matrices with entries in $\mathbb{F}_{pr}(\lambda)$, which are called *proper matrices*. A *biproper matrix* is a square proper matrix whose determinant is a biproper rational function.

By the division algorithm for polynomials, any rational function $r(\lambda) \in \mathbb{F}(\lambda)$ can be uniquely written as $r(\lambda) = p(\lambda) + r_{sp}(\lambda)$, where $p(\lambda)$ is a polynomial and $r_{sp}(\lambda)$ a strictly proper rational function. Therefore, any rational matrix $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ can be uniquely written as

$$G(\lambda) = D(\lambda) + G_{sp}(\lambda) \quad (3)$$

where $D(\lambda) \in \mathbb{F}[\lambda]^{p \times m}$ is a polynomial matrix and $G_{sp}(\lambda) \in \mathbb{F}_{pr}(\lambda)^{p \times m}$ is a *strictly proper rational matrix*, i.e., the entries of $G_{sp}(\lambda)$ are strictly proper rational functions. As said in the introduction, $D(\lambda)$ is called the polynomial part of $G(\lambda)$ and $G_{sp}(\lambda)$ its strictly proper part.

A polynomial system matrix $P(\lambda)$ of $G(\lambda)$ is said to be a polynomial system matrix in *state-space form* if $A(\lambda) = \lambda I_n - A$, $B(\lambda) = B$ and $C(\lambda) = C$ for some constant matrices $A \in \mathbb{F}^{n \times n}$, $B \in \mathbb{F}^{n \times m}$ and $C \in \mathbb{F}^{p \times n}$. It is known that any strictly proper rational matrix admits *state-space realizations* (see [29] or [22]). This means that for some positive integer n there exist constant matrices $A \in \mathbb{F}^{n \times n}$, $B \in \mathbb{F}^{n \times m}$ and $C \in \mathbb{F}^{p \times n}$ such that $G_{sp}(\lambda) = C(\lambda I_n - A)^{-1}B$ and

$$\begin{bmatrix} \lambda I_n - A & B \\ -C & D(\lambda) \end{bmatrix}$$

is a polynomial system matrix of $G(\lambda)$. Therefore $G(\lambda) = D(\lambda) + C(\lambda I_n - A)^{-1}B$. In addition, the state-space realization may always be taken of least order, or minimal, (i.e., such that the polynomial system matrix in state-space form is of least order).

Rational matrices may also have infinite eigenvalues. In order to define them, we need the notion of reversal.

Definition 2.1. Let $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ be a rational matrix expressed in the form (3). We define the reversal of $G(\lambda)$ as the rational matrix

$$\text{rev } G(\lambda) = \lambda^d G\left(\frac{1}{\lambda}\right)$$

where $d = \deg(D(\lambda))$ if $G(\lambda)$ is not strictly proper, and $d = 0$ otherwise.

Notice that this definition extends the definition of reversal for polynomial matrices (see [11, Definition 2.12] or [26, Definition 2.2]). Moreover, note that $d = 0$ if and only if $G(\lambda)$ is proper. Following the usual definition in polynomial matrices [26, Definition 2.3], we say that $G(\lambda)$ has an *eigenvalue at infinity* if $\text{rev } G(\lambda)$ has an eigenvalue at $\lambda = 0$. If $G(\lambda)$ has an eigenvalue at infinity, we say that z is a *right* (respectively *left*) *eigenvector associated to infinity* if z is a right (respectively left) eigenvector associated to 0 of $\text{rev } G(\lambda)$.

Remark 2.2. *Poles and zeros at infinity* of a rational matrix $G(\lambda)$ are defined as the poles and zeros at $\lambda = 0$ of $G(1/\lambda)$ (see [22]). If $G(\lambda)$ is not proper, i.e., $\deg(D(\lambda)) \geq 1$, $G(\lambda)$ has always a pole at ∞ (see [3]). Thus, if we define the eigenvalues at infinity of $G(\lambda)$ as those zeros that are not poles at infinity, any non-proper $G(\lambda)$ would not have eigenvalues at infinity. In particular, this would happen if $G(\lambda)$ is a polynomial matrix. Therefore, as in the polynomial case, we have considered $\text{rev } G(\lambda)$ in order to define eigenvalues at infinity. \square

Next we present the definition of strong linearization for a rational matrix given in [4]. This definition contains the notion of first invariant order at infinity q_1 of a rational matrix $G(\lambda)$. For any non strictly proper rational matrix this number is $-\deg(D(\lambda))$ where $D(\lambda)$ is the polynomial part of $G(\lambda)$ in the expression (3); otherwise, $q_1 > 0$. More information can be found in [3, 4, 32].

Definition 2.3. [4, Definition 3.4] Let $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$. Let q_1 be its first invariant order at infinity and $g = \min(0, q_1)$. Let $n = \nu(G(\lambda))$. A strong linearization of $G(\lambda)$ is a linear polynomial matrix

$$\mathcal{L}(\lambda) = \begin{bmatrix} A_1\lambda + A_0 & B_1\lambda + B_0 \\ -(C_1\lambda + C_0) & D_1\lambda + D_0 \end{bmatrix} \in \mathbb{F}[\lambda]^{(n+q) \times (n+r)} \quad (4)$$

such that the following conditions hold:

- (a) if $n > 0$ then $\det(A_1\lambda + A_0) \neq 0$, and
- (b) if $\widehat{G}(\lambda) = (D_1\lambda + D_0) + (C_1\lambda + C_0)(A_1\lambda + A_0)^{-1}(B_1\lambda + B_0)$, \widehat{q}_1 is its first invariant order at infinity and $\widehat{g} = \min(0, \widehat{q}_1)$ then:

- (i) there exist nonnegative integers s_1, s_2 , with $s_1 - s_2 = q - p = r - m$, and unimodular matrices $U_1(\lambda) \in \mathbb{F}[\lambda]^{(p+s_1) \times (p+s_1)}$ and $U_2(\lambda) \in \mathbb{F}[\lambda]^{(m+s_1) \times (m+s_1)}$ such that

$$U_1(\lambda) \operatorname{diag}(G(\lambda), I_{s_1}) U_2(\lambda) = \operatorname{diag}(\widehat{G}(\lambda), I_{s_2}), \text{ and}$$

- (ii) there exist biproper matrices $B_1(\lambda) \in \mathbb{F}_{pr}(\lambda)^{(p+s_1) \times (p+s_1)}$ and $B_2(\lambda) \in \mathbb{F}_{pr}(\lambda)^{(m+s_1) \times (m+s_1)}$ such that

$$B_1(\lambda) \operatorname{diag}(\lambda^g G(\lambda), I_{s_1}) B_2(\lambda) = \operatorname{diag}(\lambda^{\widehat{g}} \widehat{G}(\lambda), I_{s_2}).$$

It may seem that the integer $\nu(G(\lambda))$ has to be previously known in order to verify that a linear polynomial matrix as in (4) is a strong linearization of $G(\lambda)$. However, there are conditions to ensure that the size of $A_1\lambda + A_0$ is $n = \nu(G(\lambda))$. We state them in Proposition 2.4.

Proposition 2.4. *Let $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ and let*

$$\mathcal{L}(\lambda) = \begin{bmatrix} A_1\lambda + A_0 & B_1\lambda + B_0 \\ -(C_1\lambda + C_0) & D_1\lambda + D_0 \end{bmatrix} \in \mathbb{F}[\lambda]^{(n+q) \times (n+r)}$$

be a linear polynomial matrix with $n > 0$ and $\det(A_1\lambda + A_0) \neq 0$. Assume that there exist nonnegative integers s_1, s_2 , with $s_1 - s_2 = q - p = r - m$, and unimodular matrices $U_1(\lambda) \in \mathbb{F}[\lambda]^{(p+s_1) \times (p+s_1)}$ and $U_2(\lambda) \in \mathbb{F}[\lambda]^{(m+s_1) \times (m+s_1)}$ such that

$$U_1(\lambda) \operatorname{diag}(G(\lambda), I_{s_1}) U_2(\lambda) = \operatorname{diag}(\widehat{G}(\lambda), I_{s_2}), \quad (5)$$

where $\widehat{G}(\lambda) = (D_1\lambda + D_0) + (C_1\lambda + C_0)(A_1\lambda + A_0)^{-1}(B_1\lambda + B_0)$. Then $n = \nu(G(\lambda))$ if and only if the following conditions hold:

- a) A_1 is invertible, and
- b) $\operatorname{rank} \begin{bmatrix} A_1\mu + A_0 \\ C_1\mu + C_0 \end{bmatrix} = \operatorname{rank} [A_1\mu + A_0 \quad B_1\mu + B_0] = n$ for all $\mu \in \overline{\mathbb{F}}$.

Proof. Condition b) is equivalent to $\mathcal{L}(\lambda)$ being a minimal polynomial system matrix, since $\det(A_1\lambda + A_0) \neq 0$, see [29, Chapters 2 and 3]. By condition (5) and [4, Lemma 2.1], we have that $\nu(G(\lambda)) = \nu(\widehat{G}(\lambda))$. Assume that $n = \nu(G(\lambda))$. Thus, $\nu(\widehat{G}(\lambda)) = n$, and $\deg(\det(A_1\lambda + A_0)) \geq \nu(\widehat{G}(\lambda)) = n$. However, $\deg(\det(A_1\lambda + A_0)) \leq n$. Therefore, $\deg(\det(A_1\lambda + A_0)) = n$ and $\deg(\det(A_1\lambda + A_0)) = \nu(\widehat{G}(\lambda))$, which imply conditions a) and b), respectively. We assume now that conditions a) and b) hold. On the one hand, A_1 being invertible implies that $\deg(\det(A_1\lambda + A_0)) = n$. On the other hand, $\mathcal{L}(\lambda)$ being a minimal polynomial system matrix means that $\deg(\det(A_1\lambda + A_0)) = \nu(\widehat{G}(\lambda))$. Therefore, $n = \nu(\widehat{G}(\lambda)) = \nu(G(\lambda))$. □

It is known [5] that if condition (i) in Definition 2.3 holds, then condition (ii) is equivalent to the existence of unimodular matrices $W_1(\lambda)$ and $W_2(\lambda)$ such that

$$W_1(\lambda) \operatorname{diag} \left(\frac{1}{\lambda^g} G \left(\frac{1}{\lambda} \right), I_{s_1} \right) W_2(\lambda) = \operatorname{diag} \left(\frac{1}{\lambda^{\widehat{g}}} \widehat{G} \left(\frac{1}{\lambda} \right), I_{s_2} \right). \quad (6)$$

In Definition 2.3 it can always be taken $s_1 = 0$ or $s_2 = 0$, according to $p \geq q$ and $m \geq r$ or $q \geq p$ and $r \geq m$. In what follows we will consider $s_1 \geq 0$ and $s_2 = 0$. Notice that with this choice and with the notion of reversal given in Definition 2.1, (6) is equivalent to

$$W_1(\lambda) \operatorname{diag} (\operatorname{rev} G(\lambda), I_{s_1}) W_2(\lambda) = \operatorname{rev} \widehat{G}(\lambda). \quad (7)$$

Remark 2.5. Notice that Definition 2.3 extends the notion of strong linearization of polynomial matrices in the usual sense [26, Definition 2.5]. In particular, if $G(\lambda)$ is a polynomial matrix, then $n = \nu(G(\lambda)) = 0$. Therefore, a strong linearization $\mathcal{L}(\lambda)$ of $G(\lambda)$ is of the form $\mathcal{L}(\lambda) = D_1\lambda + D_0$, with $\widehat{G}(\lambda) = \mathcal{L}(\lambda)$, $g = q_1 = -\deg(G(\lambda))$ and $\widehat{g} = \widehat{q}_1 = -\deg(\mathcal{L}(\lambda))$. \square

Let us denote by $\mathcal{N}_r(G(\lambda))$ and $\mathcal{N}_\ell(G(\lambda))$ the *right* and *left null-spaces* over $\mathbb{F}(\lambda)$ of $G(\lambda)$, respectively, i.e., if $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$,

$$\begin{aligned} \mathcal{N}_r(G(\lambda)) &= \{x(\lambda) \in \mathbb{F}(\lambda)^{m \times 1} : G(\lambda)x(\lambda) = 0\}, \\ \mathcal{N}_\ell(G(\lambda)) &= \{x(\lambda) \in \mathbb{F}(\lambda)^{p \times 1} : x(\lambda)^T G(\lambda) = 0\}. \end{aligned}$$

When $G(\lambda)$ is singular at least one of these null spaces is nontrivial. A spectral characterization of strong linearizations is given in [4, Theorem 3.10]. It says that $\mathcal{L}(\lambda)$ is a strong linearization of $G(\lambda)$ if and only if $\dim \mathcal{N}_r(G(\lambda)) = \dim \mathcal{N}_r(\mathcal{L}(\lambda))$ and $\mathcal{L}(\lambda)$ preserves the finite and infinite structures of poles and zeros of $G(\lambda)$ in the sense of [4, Definitions 3.8 and 3.9]. This characterization is the key property of strong linearizations in the realm of REPs.

Remark 2.6. The equality $\dim \mathcal{N}_r(G(\lambda)) = \dim \mathcal{N}_r(\mathcal{L}(\lambda))$ is equivalent to $\dim \mathcal{N}_\ell(G(\lambda)) = \dim \mathcal{N}_\ell(\mathcal{L}(\lambda))$. Consider $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ and $\mathcal{L}(\lambda) \in \mathbb{F}(\lambda)^{(p+q) \times (m+q)}$ with $q \geq 0$. By the rank-nullity theorem $\dim \mathcal{N}_\ell(G(\lambda)) = p - \operatorname{rank} G(\lambda)$ and $\dim \mathcal{N}_r(G(\lambda)) = m - \operatorname{rank} G(\lambda)$. Therefore $\operatorname{rank} \mathcal{L}(\lambda) = q + \operatorname{rank} G(\lambda)$ if and only if $\dim \mathcal{N}_r(G(\lambda)) = \dim \mathcal{N}_r(\mathcal{L}(\lambda))$. And $\operatorname{rank} \mathcal{L}(\lambda) = q + \operatorname{rank} G(\lambda)$ if and only if $\dim \mathcal{N}_\ell(G(\lambda)) = \dim \mathcal{N}_\ell(\mathcal{L}(\lambda))$. \square

Lemma 2.7 follows from Definition 2.3. It shows an easy way to obtain strong linearizations for a rational matrix $G(\lambda)$ from a particular strong linearization $\mathcal{L}(\lambda)$ by multiplying $\mathcal{L}(\lambda)$ by some appropriate matrices. This simple result is fundamental in this paper, and we conjecture that it will be fundamental for constructing (in the future) other families of strong linearizations of rational matrices.

Lemma 2.7. *Let $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ be a rational matrix, and let*

$$\mathcal{L}_1(\lambda) = \begin{bmatrix} A_1\lambda + A_0 & B_1\lambda + B_0 \\ -(C_1\lambda + C_0) & D_1\lambda + D_0 \end{bmatrix} \in \mathbb{F}[\lambda]^{(n+(p+s)) \times (n+(m+s))}$$

be a strong linearization of $G(\lambda)$. Consider $Q_1, Q_3 \in \mathbb{F}^{n \times n}$, $Q_2 \in \mathbb{F}^{(p+s) \times (p+s)}$, $Q_4 \in \mathbb{F}^{(m+s) \times (m+s)}$ nonsingular matrices, $W \in \mathbb{F}^{(p+s) \times n}$, and $Z \in \mathbb{F}^{n \times (m+s)}$. Then the linear polynomial matrix

$$\mathcal{L}_2(\lambda) = \begin{bmatrix} Q_1 & 0 \\ W & Q_2 \end{bmatrix} \mathcal{L}_1(\lambda) \begin{bmatrix} Q_3 & Z \\ 0 & Q_4 \end{bmatrix}$$

is a strong linearization of $G(\lambda)$.

Proof. Let us write

$$\mathcal{L}_2(\lambda) = \begin{bmatrix} A_2\lambda + \tilde{A}_0 & B_2\lambda + \tilde{B}_0 \\ -(C_2\lambda + \tilde{C}_0) & D_2\lambda + \tilde{D}_0 \end{bmatrix}.$$

We have $\det(A_2\lambda + \tilde{A}_0) \neq 0$ if $n > 0$, since $A_2\lambda + \tilde{A}_0 = Q_1(A_1\lambda + A_0)Q_3$. Let us consider the transfer functions $\widehat{G}_1(\lambda)$, $\widehat{G}_2(\lambda)$ of $\mathcal{L}_1(\lambda)$, $\mathcal{L}_2(\lambda)$, respectively. They satisfy $\widehat{G}_2(\lambda) = Q_2\widehat{G}_1(\lambda)Q_4$. Let q_1 be the first invariant order at infinity of $G(\lambda)$ and $g = \min(0, q_1)$. For $i = 1, 2$, let $\widehat{g}_i = \min(0, \widehat{q}_i)$, where \widehat{q}_i is the first invariant order at infinity of $\widehat{G}_i(\lambda)$. Since $\mathcal{L}_1(\lambda)$ is a strong linearization of $G(\lambda)$, there exist unimodular matrices $U_1(\lambda)$ and $U_2(\lambda)$ such that $U_1(\lambda) \text{diag}(G(\lambda), I_s)U_2(\lambda) = \widehat{G}_1(\lambda)$, and biproper matrices $B_1(\lambda)$ and $B_2(\lambda)$ such that $B_1(\lambda) \text{diag}(\lambda^g G(\lambda), I_s)B_2(\lambda) = \lambda^{\widehat{g}_1} \widehat{G}_1(\lambda)$. By using the equality $\widehat{G}_2(\lambda) = Q_2\widehat{G}_1(\lambda)Q_4$, we have that $\widehat{g}_1 = \widehat{g}_2$, and by the same equality, we get

$$Q_2U_1(\lambda) \text{diag}(G(\lambda), I_s)U_2(\lambda)Q_4 = \widehat{G}_2(\lambda),$$

and

$$Q_2B_1(\lambda) \text{diag}(\lambda^g G(\lambda), I_s)B_2(\lambda)Q_4 = \lambda^{\widehat{g}_2} \widehat{G}_2(\lambda).$$

Then we obtain that conditions (a) and (b) in Definition 2.3 hold for $\mathcal{L}_2(\lambda)$. \square

Strong linearizations of a rational matrix $G(\lambda)$ expressed in the form (3) can be constructed from combining minimal state-space realizations of the strictly proper matrix $G_{sp}(\lambda)$ and strong linearizations of its polynomial part $D(\lambda)$. In particular, strong block minimal bases pencils associated to $D(\lambda)$ with sharp degree can be used (see [4]). The definition of this concept is taken from [13] and appears also in [4]. As in [13], we will say that a polynomial matrix $K(\lambda) \in \mathbb{F}[\lambda]^{p \times m}$ (with $p < m$) is a *minimal basis* if its rows form a minimal basis of the rational subspace they span (see [16]). Moreover, a minimal basis $N(\lambda) \in \mathbb{F}[\lambda]^{q \times m}$ is said to be *dual* to $K(\lambda)$ if $p + q = m$ and $K(\lambda)N(\lambda)^T = 0$ (see [13, Definition 2.5]).

Definition 2.8. [4, Definition 5.2] Let $D(\lambda) \in \mathbb{F}[\lambda]^{p \times m}$ be a polynomial matrix. A strong block minimal bases pencil associated to $D(\lambda)$ is a linear polynomial matrix with the following structure

$$\mathcal{L}(\lambda) = \begin{bmatrix} \underbrace{M(\lambda)}_{m+\widehat{m}} & \underbrace{K_2(\lambda)^T}_{\widehat{p}} \\ \underbrace{K_1(\lambda)}_{m+\widehat{m}} & 0 \end{bmatrix} \begin{array}{l} \} p+\widehat{p} \\ \} \widehat{m} \end{array}, \quad (8)$$

where $K_1(\lambda) \in \mathbb{F}[\lambda]^{\widehat{m} \times (m+\widehat{m})}$ (respectively $K_2(\lambda) \in \mathbb{F}[\lambda]^{\widehat{p} \times (p+\widehat{p})}$) is a minimal basis with all its row degrees equal to 1 and with the row degrees of a minimal basis $N_1(\lambda) \in \mathbb{F}[\lambda]^{m \times (m+\widehat{m})}$ (respectively $N_2(\lambda) \in \mathbb{F}[\lambda]^{p \times (p+\widehat{p})}$) dual to $K_1(\lambda)$ (respectively $K_2(\lambda)$) all equal, and such that

$$D(\lambda) = N_2(\lambda)M(\lambda)N_1(\lambda)^T. \quad (9)$$

If, in addition, $\deg(D(\lambda)) = \deg(N_2(\lambda)) + \deg(N_1(\lambda)) + 1$ then $\mathcal{L}(\lambda)$ is said to be a strong block minimal bases pencil associated to $D(\lambda)$ with sharp degree.

Remark 2.9. The following useful characterization of minimal bases will be used (see [16, Main Theorem] or [13, Theorem 2.2]). Namely, $K(\lambda) \in \mathbb{F}[\lambda]^{p \times m}$ is a minimal basis if and only if $K(\lambda_0)$ has full row rank for all $\lambda_0 \in \overline{\mathbb{F}}$ and $K(\lambda)$ is *row reduced*, i.e., its highest row degree coefficient matrix has full row rank (see [13, Definition 2.1]). \square

Remark 2.10. A first application of the key Lemma 2.7 is to construct strong linearizations of a rational matrix $G(\lambda)$ from any Fiedler-like strong linearization $L_F(\lambda)$ of its polynomial part $D(\lambda)$. For this purpose, note that [8, Theorems 3.8, 3.15, 3.16] guarantee that there exist permutation matrices Π_1 and Π_2 and a strong block minimal bases pencil $L(\lambda)$ associated to $D(\lambda)$ such that $L_F(\lambda) = \Pi_1 L(\lambda) \Pi_2$. In addition, Theorem 5.11 in [4] explains how to construct a strong linearization $\mathcal{L}(\lambda)$ of $G(\lambda)$ from $L(\lambda)$. Thus, according to Lemma 2.7, $\text{diag}(I_n, \Pi_1) \mathcal{L}(\lambda) \text{diag}(I_n, \Pi_2)$ is a strong linearization of $G(\lambda)$ based on $L_F(\lambda)$. \square

In what follows, the Kronecker product of two matrices A and B , denoted by $A \otimes B$, will be used (see [21, Chapter 4]).

3. \mathbb{M}_1 -strong linearizations

In this section and in Section 4 we present strong linearizations of square rational matrices $G(\lambda)$ with polynomial part $D(\lambda)$ expressed in an orthogonal basis. More precisely, we consider strong linearizations of $D(\lambda)$ that belong to the ansatz spaces $\mathbb{M}_1(D)$ or $\mathbb{M}_2(D)$, recently developed by H. Faßbender and P. Saltenberger in [15], and based on them, we construct strong linearizations of $G(\lambda)$ by using Lemma 2.7 and the strong linearizations presented in [4, Section 5.2].

As said in the preliminaries, we consider an arbitrary field \mathbb{F} throughout this paper, although the results in [15] are stated only for the real field \mathbb{R} . Nevertheless, the results of [15] that are used in this paper are also valid for any field \mathbb{F} . We consider a polynomial basis $\{\phi_j(\lambda)\}_{j=0}^\infty$ of $\mathbb{F}[\lambda]$, viewed as an \mathbb{F} -vector space, with $\phi_j(\lambda)$ a polynomial of degree j , that satisfies the following three-term recurrence relation:

$$\alpha_j \phi_{j+1}(\lambda) = (\lambda - \beta_j) \phi_j(\lambda) - \gamma_j \phi_{j-1}(\lambda) \quad j \geq 0 \quad (10)$$

where $\alpha_j, \beta_j, \gamma_j \in \mathbb{F}$, $\alpha_j \neq 0$, $\phi_{-1}(\lambda) = 0$, and $\phi_0(\lambda) = 1$. Let $P(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ be a polynomial matrix of degree k written in terms of this basis as follows

$$P(\lambda) = P_k \phi_k(\lambda) + P_{k-1} \phi_{k-1}(\lambda) + \cdots + P_1 \phi_1(\lambda) + P_0 \phi_0(\lambda). \quad (11)$$

We define $\Phi_k(\lambda) = [\phi_{k-1}(\lambda) \cdots \phi_1(\lambda) \phi_0(\lambda)]^T$ and $V_P = \{v \otimes P(\lambda) : v \in \mathbb{F}^k\}$, and we consider the set of pencils

$$\mathbb{M}_1(P) = \{L(\lambda) = \lambda X + Y : X, Y \in \mathbb{F}^{km \times km}, L(\lambda)(\Phi_k(\lambda) \otimes I_m) \in V_P\}.$$

A pencil $L(\lambda) \in \mathbb{M}_1(P)$, which verifies $L(\lambda)(\Phi_k(\lambda) \otimes I_m) = v \otimes P(\lambda)$ for some vector $v \in \mathbb{F}^k$, is said to have *right ansatz vector* v . A particular pencil in $\mathbb{M}_1(P)$ introduced in [15, page 63] is

$$F_\Phi^P(\lambda) = \begin{bmatrix} m_\Phi^P(\lambda) \\ M_\Phi(\lambda) \otimes I_m \end{bmatrix} \in \mathbb{F}[\lambda]^{km \times km}, \quad (12)$$

where

$$m_\Phi^P(\lambda) = \begin{bmatrix} (\lambda - \beta_{k-1}) P_k + P_{k-1} & P_{k-2} - \frac{\gamma_{k-1}}{\alpha_{k-1}} P_k & P_{k-3} & \cdots & P_1 & P_0 \end{bmatrix},$$

Note that $\Phi_k(\lambda)^T$ is a minimal basis because $\phi_0(\lambda) = 1$, so $\Phi_k(\lambda_0)$ has rank 1 for all $\lambda_0 \in \overline{\mathbb{F}}$, and

$$[\Phi_k^T]_{hr} = \begin{bmatrix} \frac{1}{\alpha_0 \alpha_1 \cdots \alpha_{k-2}} & 0 & \cdots & 0 \end{bmatrix}$$

has also rank 1. Therefore, $N(\lambda) = \Phi_k(\lambda)^T \otimes I_m$ is also a minimal basis. Since $K(\lambda)N(\lambda)^T = (M_\Phi(\lambda) \otimes I_m)(\Phi_k(\lambda) \otimes I_m) = 0$ and $\begin{bmatrix} K(\lambda) \\ N(\lambda) \end{bmatrix}$ is a square matrix, we have that $K(\lambda)$ and $N(\lambda)$ are dual minimal bases. In addition, it is obvious that all the row degrees of $K(\lambda)$ are equal to 1 and all the row degrees of $\Phi_k(\lambda)^T \otimes I_m$ are equal to $k - 1$. Hence, $F_\Phi^P(\lambda)$ is a strong block minimal bases pencil associated to the polynomial matrix $M(\lambda)N(\lambda)^T = m_\Phi^P(\lambda)(\Phi_k(\lambda) \otimes I_m) = P(\lambda)$ and $\deg(P(\lambda)) = 1 + \deg(N(\lambda))$, which means that $F_\Phi^P(\lambda)$ has sharp degree. \square

Since every strong block minimal bases pencil is a strong linearization (see [13, Theorem 3.3]), the following corollary is straightforward.

Corollary 3.4. $F_\Phi^P(\lambda)$ is a strong linearization for $P(\lambda)$.

The proof of the next result is trivial because if $L(\lambda) = [v \otimes I_m \quad H]F_\Phi^P(\lambda)$ with $[v \otimes I_m \quad H]$ nonsingular then $L(\lambda)$ is strictly equivalent to $F_\Phi^P(\lambda)$.

Corollary 3.5. [15, Corollary 2.1] Let $L(\lambda) = [v \otimes I_m \quad H]F_\Phi^P(\lambda) \in \mathbb{M}_1(P)$. If $[v \otimes I_m \quad H]$ is nonsingular then $L(\lambda)$ is a strong linearization for $P(\lambda)$.

Remark 3.6. Although $F_\Phi^P(\lambda)$ is a strong block minimal bases pencil associated to $P(\lambda)$ this structure is not preserved in general when we multiply on the left by a nonsingular matrix $[v \otimes I_m \quad H]$. For example, consider the polynomial matrix $P(\lambda) = I\lambda^3 + 2I\lambda^2 + I\lambda + S \in \mathbb{R}[\lambda]^{2 \times 2}$ expressed in the monomial basis, where $S = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}$ and I stands for

I_2 . In this case, the matrix $F_\Phi^P(\lambda)$ is $F_\Phi^P(\lambda) = \begin{bmatrix} \lambda I + 2I & I & S \\ -I & \lambda I & 0 \\ 0 & -I & \lambda I \end{bmatrix}$. Let $v = [1 \quad 1 \quad 0]^T$

and $H = \begin{bmatrix} 0 & 0 \\ I & 0 \\ 0 & I \end{bmatrix}$ and let $L(\lambda) = [v \otimes I \quad H]F_\Phi^P(\lambda) = \begin{bmatrix} \lambda I + 2I & I & S \\ \lambda I + I & \lambda I + I & S \\ 0 & -I & \lambda I \end{bmatrix}$. Notice

that if $L(\lambda)$ were a strong block minimal bases pencil associated to $P(\lambda)$, one of these two different situations would happen in (8):

1. $M(\lambda) = [\lambda I + 2I \quad I \quad S]$, $K_1(\lambda) = \begin{bmatrix} \lambda I + I & \lambda I + I & S \\ 0 & -I & \lambda I \end{bmatrix}$ and $K_2(\lambda)$ empty.

2. $M(\lambda) = \begin{bmatrix} \lambda I + 2I \\ \lambda I + I \\ 0 \end{bmatrix}$, $K_2(\lambda)^T = \begin{bmatrix} I & S \\ \lambda I + I & S \\ -I & \lambda I \end{bmatrix}$ and $K_1(\lambda)$ empty.

In the first case, the matrix $K_1(\lambda)$ has not full row rank for $\lambda = -1$. In the second case, the matrix $K_2(\lambda)$ has not full row rank for $\lambda = 0$. Therefore, $L(\lambda)$ is not a strong block minimal bases pencil associated to $P(\lambda)$. \square

From the fact that $F_{\Phi}^P(\lambda)$ is a strong block minimal bases pencil, we can obtain strong linearizations for rational matrices by applying Theorem 5.11 in [4]. For this purpose, we prove first the following lemma.

Lemma 3.7. *The matrix*

$$U(\lambda) = \begin{bmatrix} M_{\Phi}(\lambda) \otimes I_m \\ e_k^T \otimes I_m \end{bmatrix} = \begin{bmatrix} M_{\Phi}(\lambda) \\ e_k^T \end{bmatrix} \otimes I_m$$

is unimodular, and its inverse has the form $U(\lambda)^{-1} = [\widehat{\Phi}_k(\lambda) \quad \Phi_k(\lambda) \otimes I_m]$ with $\widehat{\Phi}_k(\lambda) \in \mathbb{F}[\lambda]^{km \times (k-1)m}$.

Proof. Let us consider the matrix

$$\tilde{U}(\lambda) = \begin{bmatrix} M_{\Phi}(\lambda) \\ e_k^T \end{bmatrix} = \begin{bmatrix} -\alpha_{k-2} & (\lambda - \beta_{k-2}) & -\gamma_{k-2} & & & & \\ & -\alpha_{k-3} & (\lambda - \beta_{k-3}) & -\gamma_{k-3} & & & \\ & & \ddots & \ddots & \ddots & & \\ & & & -\alpha_1 & (\lambda - \beta_1) & -\gamma_1 & \\ & & & & -\alpha_0 & (\lambda - \beta_0) & \\ 0 & \dots & & & 0 & 1 & \end{bmatrix}.$$

Since $\tilde{U}(\lambda)$ is upper triangular, its determinant is $(-\alpha_{k-2}) \cdots (-\alpha_0)$, i.e., a constant different from zero. Therefore, $\tilde{U}(\lambda)$ is unimodular. Finally, note that $\tilde{U}(\lambda)\Phi_k(\lambda) = e_k \in \mathbb{F}^k$. Thus $\Phi_k(\lambda)$ is the last column of $\tilde{U}(\lambda)^{-1}$. \square

Theorem 3.8. *Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix, let $G(\lambda) = D(\lambda) + G_{sp}(\lambda)$ be its unique decomposition into its polynomial part $D(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ and its strictly proper part $G_{sp}(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$, and let $G_{sp}(\lambda) = C(\lambda I_n - A)^{-1}B$ be a minimal state-space realization of $G_{sp}(\lambda)$. Assume that $\deg(D(\lambda)) \geq 2$. Write $D(\lambda)$ in terms of the polynomial basis $\{\phi_j(\lambda)\}_{j=0}^{\infty}$ satisfying the three-term recurrence relation (10), as*

$$D(\lambda) = D_k \phi_k(\lambda) + D_{k-1} \phi_{k-1}(\lambda) + \cdots + D_1 \phi_1(\lambda) + D_0 \phi_0(\lambda) \quad (13)$$

with $D_k \neq 0$, and let $F_{\Phi}^D(\lambda)$ be the matrix pencil defined as in (12). Then, for any nonsingular matrices $X, Y \in \mathbb{F}^{n \times n}$ the linear polynomial matrix

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|cc} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} & XB \\ \hline -CY & & F_{\Phi}^D(\lambda) \\ 0_{(k-1)m \times n} & & \end{array} \right]$$

is a strong linearization of $G(\lambda)$.

Proof. Lemmas 3.3 and 3.7 allow us to apply [4, Theorem 5.11], with $K_1(\lambda) = M_{\Phi}(\lambda) \otimes I_m$, $\widehat{K}_1 = e_k^T \otimes I_m$, $K_2(\lambda)^T$ empty and $\widehat{K}_2^T = I_m$. \square

Then, from combining Lemma 2.7 and Theorem 3.8 we obtain strong linearizations of a rational matrix from strong linearizations in $\mathbb{M}_1(D)$ of its polynomial part.

Theorem 3.9. *Under the same assumptions as in Theorem 3.8, let $v \in \mathbb{F}^k$, $H \in \mathbb{F}^{km \times (k-1)m}$ with $[v \otimes I_m \ H]$ nonsingular and let $L(\lambda) = [v \otimes I_m \ H] F_{\Phi}^D(\lambda) \in \mathbb{M}_1(D)$. Then, the linear polynomial matrix*

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} \ XB \\ \hline -(v \otimes I_m)CY & L(\lambda) \end{array} \right]$$

is a strong linearization of $G(\lambda)$.

Proof. Set $K = [v \otimes I_m \ H]$. If K is nonsingular then, by Lemma 2.7 and Theorem 3.8,

$$\begin{aligned} \mathcal{L}(\lambda) &= \begin{bmatrix} I_n & 0 \\ 0 & K \end{bmatrix} \left[\begin{array}{c|c} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} \ XB \\ \hline -CY & F_{\Phi}^D(\lambda) \\ 0_{(k-1)m \times n} & \end{array} \right] \\ &= \left[\begin{array}{c|c} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} \ XB \\ \hline -(v \otimes I_m)CY & L(\lambda) \end{array} \right]. \end{aligned}$$

is a strong linearization of $G(\lambda)$. \square

The strong linearizations of square rational matrices constructed in Theorem 3.9 will be called \mathbb{M}_1 -strong linearizations.

4. \mathbb{M}_2 -strong linearizations

In this section we obtain strong linearizations of a square rational matrix from the transposed version of $\mathbb{M}_1(P)$, where $P(\lambda)$ is the polynomial matrix in (11). Since the proofs of the results are similar to those in Section 3, they are omitted for brevity. We define $W_P = \{w^T \otimes P(\lambda) : w \in \mathbb{F}^k\}$, and we consider the set of pencils

$$\mathbb{M}_2(P) = \{L(\lambda) = \lambda X + Y : X, Y \in \mathbb{F}^{km \times km}, (\Phi_k(\lambda)^T \otimes I_m)L(\lambda) \in W_P\}.$$

A pencil $L(\lambda) \in \mathbb{M}_2(P)$, which verifies $(\Phi_k(\lambda)^T \otimes I_m)L(\lambda) = w^T \otimes P(\lambda)$ for some vector $w \in \mathbb{F}^k$, is said to have *left ansatz vector* w . Pencils in $\mathbb{M}_2(P)$ are characterized in [15, Theorem 2]. We need the definition of the block-transpose of a $km \times lm$ pencil $L(\lambda)$.

If we express $L(\lambda)$ as $L(\lambda) = \sum_{i=1}^k \sum_{j=1}^l e_i e_j^T \otimes L_{ij}(\lambda)$ for certain $m \times m$ pencils $L_{ij}(\lambda)$,

where e_i denotes the i th canonical vector in \mathbb{F}^k , and e_j the j th canonical vector in \mathbb{F}^l ,

$L(\lambda)^{\mathcal{B}} = \sum_{i=1}^k \sum_{j=1}^l e_j e_i^T \otimes L_{ij}(\lambda)$ is called the *block-transpose* of $L(\lambda)$ (see [20]). Notice that

$$F_{\Phi}^P(\lambda)^{\mathcal{B}} = [m_{\Phi}^P(\lambda)^{\mathcal{B}} \ M_{\Phi}(\lambda)^T \otimes I_m].$$

Theorem 4.1. [15, Theorem 2] *Let $P(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ be a polynomial matrix with degree $k \geq 2$. Then $L(\lambda) \in \mathbb{M}_2(P)$ with left ansatz vector $w \in \mathbb{F}^k$ if and only if*

$$L(\lambda) = F_{\Phi}^P(\lambda)^{\mathcal{B}} \begin{bmatrix} w^T \otimes I_m \\ H^{\mathcal{B}} \end{bmatrix}$$

for some matrix $H \in \mathbb{F}^{km \times (k-1)m}$ partitioned into $k \times (k-1)$ blocks each of size $m \times m$.

The vector space $\mathbb{M}_2(P)$ reduces to the well-known space $\mathbb{L}_2(P)$ when $\{\phi_k(\lambda)\}_{k=0}^\infty$ is the monomial basis, see [26]. Lemma 4.2 is for $\mathbb{M}_2(P)$ the counterpart of Lemma 3.3 for $\mathbb{M}_1(P)$ and can be used to proceed with $\mathbb{M}_2(P)$ analogously as we did with $\mathbb{M}_1(P)$.

Lemma 4.2. $F_\Phi^P(\lambda)^\mathcal{B}$ is a strong block minimal bases pencil with only one block row associated to $P(\lambda)$ with sharp degree.

In particular, Lemma 4.2 allows us to apply [4, Theorem 5.11] to the strong linearization $F_\Phi^D(\lambda)^\mathcal{B}$ of the polynomial part of a square rational matrix, with $K_2(\lambda) = M_\Phi(\lambda) \otimes I_m$, $\hat{K}_2 = e_k^T \otimes I_m$, $K_1(\lambda)$ empty and $\hat{K}_1 = I_m$. Thus, we get the following results to obtain strong linearizations of a square rational matrix $G(\lambda) = D(\lambda) + G_{sp}(\lambda)$ expressed as in (3) from strong linearizations in $\mathbb{M}_2(D)$.

Theorem 4.3. Under the same assumptions as in Theorem 3.8, the linear polynomial matrix

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & XB \quad 0_{n \times (k-1)m} \\ \hline 0_{(k-1)m \times n} & F_\Phi^D(\lambda)^\mathcal{B} \\ -CY & \end{array} \right]$$

is a strong linearization of $G(\lambda)$.

Theorem 4.4. Under the same assumptions as in Theorem 3.8, let $w \in \mathbb{F}^k$, $H \in \mathbb{F}^{km \times (k-1)m}$ with $\begin{bmatrix} w^T \otimes I_m \\ H^\mathcal{B} \end{bmatrix}$ nonsingular and let $L(\lambda) = F_\Phi^D(\lambda)^\mathcal{B} \begin{bmatrix} w^T \otimes I_m \\ H^\mathcal{B} \end{bmatrix} \in \mathbb{M}_2(D)$. Then the linear polynomial matrix

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & XB(w^T \otimes I_m) \\ \hline 0_{(k-1)m \times n} & L(\lambda) \\ -CY & \end{array} \right]$$

is a strong linearization of $G(\lambda)$.

Proof. We apply Lemma 2.7 by multiplying on the right the matrix $\mathcal{L}(\lambda)$ in Theorem 4.3 by the matrix $\begin{bmatrix} I_n & 0 \\ 0 & K \end{bmatrix}$ with $K = \begin{bmatrix} w^T \otimes I_m \\ H^\mathcal{B} \end{bmatrix}$ nonsingular. \square

The strong linearizations of rational matrices constructed in Theorem 4.4 will be called \mathbb{M}_2 -strong linearizations.

5. Recovering eigenvectors from \mathbb{M}_1 - and \mathbb{M}_2 - strong linearizations of rational matrices

In this section we will recover right and left eigenvectors of a rational matrix. These eigenvectors will be obtained without essentially computational cost from the right and left eigenvectors of the strong linearizations that we have constructed in Theorems 3.9 and 4.4. Previously, and due to the fact that we can see strong linearizations as polynomial system matrices, in Subsection 5.1 we will see the relation between the eigenvectors of a polynomial system matrix and the eigenvectors of its transfer function matrix. For the sake of brevity, in this section the following nomenclature is adopted: “ (λ_0, x_0) is a solution of the REP $G(\lambda)x = 0$ ” means that λ_0 is a finite eigenvalue of $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ and x_0 is a right

eigenvector corresponding to λ_0 , and “ (λ_0, x_0) is a solution of the REP $x^T G(\lambda) = 0$ ” means that λ_0 is a finite eigenvalue of $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ and x_0 is a left eigenvector corresponding to λ_0 . An analogous notation is adopted for polynomial eigenvalue problems (PEPs) and the particular case of linear eigenvalue problems (LEPs).

5.1. Eigenvectors of polynomial system and transfer function matrices

We know from [5, Proposition 3.1] how to recover right eigenvectors of a polynomial system matrix $P(\lambda)$ from those of its transfer function $G(\lambda)$, and conversely. In Proposition 5.1 we state an extended version of [5, Proposition 3.1] that includes a result about the null-spaces of $P(\lambda)$ and $G(\lambda)$ evaluated at the eigenvalue of interest. That is, for a finite eigenvalue λ_0 of a rational matrix $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ we denote by $\mathcal{N}_r(G(\lambda_0))$ the right null-space over $\overline{\mathbb{F}}$ of $G(\lambda_0)$, i.e., $\mathcal{N}_r(G(\lambda_0)) = \{x \in \overline{\mathbb{F}}^{m \times 1} : G(\lambda_0)x = 0\}$. We state without proof the analogous result for left eigenvectors and null-spaces in Proposition 5.2.

In what follows, we assume that eigenvectors of the form $\begin{bmatrix} y \\ x \end{bmatrix}$ are partitioned conformable to the corresponding polynomial system matrix.

Proposition 5.1. *Let $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ be a rational matrix and*

$$P(\lambda) = \begin{bmatrix} A(\lambda) & B(\lambda) \\ -C(\lambda) & D(\lambda) \end{bmatrix} \in \mathbb{F}[\lambda]^{(n+p) \times (n+m)}$$

be any polynomial system matrix with $G(\lambda)$ as transfer function matrix.

- a) *If $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix}\right)$ is a solution of the PEP $P(\lambda)z = 0$ such that $\det A(\lambda_0) \neq 0$, then (λ_0, x_0) is a solution of the REP $G(\lambda)x = 0$.*
- b) *Moreover, if $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_r(P(\lambda_0))$, with $\det A(\lambda_0) \neq 0$, then $\{x_1, \dots, x_t\}$ is a basis of $\mathcal{N}_r(G(\lambda_0))$.*
- c) *Conversely, if (λ_0, x_0) is a solution of the REP $G(\lambda)x = 0$ such that $\det A(\lambda_0) \neq 0$ and y_0 is defined as the unique solution of $A(\lambda_0)y_0 + B(\lambda_0)x_0 = 0$, then $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix}\right)$ is a solution of the PEP $P(\lambda)z = 0$.*
- d) *Moreover, if $\{x_1, \dots, x_t\}$ is a basis of $\mathcal{N}_r(G(\lambda_0))$, with $\det A(\lambda_0) \neq 0$, and, for $i = 1, \dots, t$, y_i is defined as the unique solution of $A(\lambda_0)y_i + B(\lambda_0)x_i = 0$, then $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_r(P(\lambda_0))$.*

Proof. The statements a) and c) are the results in [5, Proposition 3.1] stated here for a rectangular matrix $G(\lambda)$. The proofs are exactly the same as in [5] and, therefore, are omitted. To prove b) and d) we write

$$\begin{bmatrix} A(\lambda_0) & B(\lambda_0) \\ -C(\lambda_0) & D(\lambda_0) \end{bmatrix} = \begin{bmatrix} I_n & 0 \\ -C(\lambda_0)A(\lambda_0)^{-1} & I_p \end{bmatrix} \begin{bmatrix} A(\lambda_0) & 0 \\ 0 & G(\lambda_0) \end{bmatrix} \begin{bmatrix} I_n & A(\lambda_0)^{-1}B(\lambda_0) \\ 0 & I_m \end{bmatrix}.$$

Since $\det A(\lambda_0) \neq 0$, $\text{rank} P(\lambda_0) = n + \text{rank} G(\lambda_0)$. Therefore

$$\dim \mathcal{N}_r(P(\lambda_0)) = \dim \mathcal{N}_r(G(\lambda_0)). \quad (14)$$

Then *b*) and *d*) are obtained by using *a*) and *c*), respectively, taking (14) and the linear independence of the considered sets into account, and observing that $P(\lambda_0) \begin{bmatrix} y_0 \\ x_0 \end{bmatrix} = 0$ if and only if $y_0 = -A(\lambda_0)^{-1}B(\lambda_0)x_0$ and $G(\lambda_0)x_0 = 0$. \square

Proposition 5.2 is an analogous result to Proposition 5.1 for left eigenvectors and left null-spaces as we announced, and it can be proved in a similar way. The left null-space of $G(\lambda_0) \in \overline{\mathbb{F}}^{p \times m}$ is denoted and defined as $\mathcal{N}_\ell(G(\lambda_0)) = \{x \in \overline{\mathbb{F}}^{p \times 1} : x^T G(\lambda_0) = 0\}$.

Proposition 5.2. *Let $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ be a rational matrix and*

$$P(\lambda) = \begin{bmatrix} A(\lambda) & B(\lambda) \\ -C(\lambda) & D(\lambda) \end{bmatrix} \in \mathbb{F}[\lambda]^{(n+p) \times (n+m)}$$

be any polynomial system matrix with $G(\lambda)$ as transfer function matrix.

- a) If $(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix})$ is a solution of the PEP $z^T P(\lambda) = 0$ such that $\det A(\lambda_0) \neq 0$, then (λ_0, x_0) is a solution of the REP $x^T G(\lambda) = 0$.*
- b) Moreover, if $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_q \\ x_q \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_\ell(P(\lambda_0))$, with $\det A(\lambda_0) \neq 0$, then $\{x_1, \dots, x_q\}$ is a basis of $\mathcal{N}_\ell(G(\lambda_0))$.*
- c) Conversely, if (λ_0, x_0) is a solution of the REP $x^T G(\lambda) = 0$ such that $\det A(\lambda_0) \neq 0$, and y_0 is defined as the unique solution of $y_0^T A(\lambda_0) - x_0^T C(\lambda_0) = 0$, then $(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix})$ is a solution of the PEP $z^T P(\lambda) = 0$.*
- d) Moreover, if $\{x_1, \dots, x_q\}$ is a basis of $\mathcal{N}_\ell(G(\lambda_0))$, with $\det A(\lambda_0) \neq 0$, and, for $i = 1, \dots, q$, y_i is defined as the unique solution of $y_i^T A(\lambda_0) - x_i^T C(\lambda_0) = 0$, then $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_q \\ x_q \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_\ell(P(\lambda_0))$.*

Remark 5.3. If $G(\lambda)$ is singular, then for any $\lambda_0 \in \overline{\mathbb{F}}$ that is not a pole of $G(\lambda)$, including those λ_0 that are not eigenvalues of $G(\lambda)$, $\mathcal{N}_r(G(\lambda_0)) \neq \{0\}$ or $\mathcal{N}_\ell(G(\lambda_0)) \neq \{0\}$. The reader can check easily that Propositions 5.1 and 5.2 remain valid for any $\lambda_0 \in \overline{\mathbb{F}}$ that is not a pole of $G(\lambda)$ in the case $G(\lambda)$ is singular. \square

5.2. Eigenvectors from \mathbb{M}_1 -strong linearizations

We consider in this subsection the linearizations that we have constructed in Theorem 3.9, which we called \mathbb{M}_1 -strong linearizations. We will recover the eigenvectors of a rational matrix $G(\lambda)$ from those of its \mathbb{M}_1 -strong linearizations, and conversely. Lemma 5.4 will be used for this purpose.

Lemma 5.4. Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix with polynomial part of degree $k \geq 2$, let

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} \quad XB \\ \hline -(v \otimes I_m)CY & L(\lambda) \end{array} \right]$$

be an \mathbb{M}_1 -strong linearization of $G(\lambda)$, and let $\widehat{G}(\lambda)$ be the transfer function of $\mathcal{L}(\lambda)$. Then

$$\widehat{G}(\lambda)(\Phi_k(\lambda) \otimes I_m) = v \otimes G(\lambda). \quad (15)$$

Proof. We consider the transfer function of the matrix $\mathcal{L}(\lambda)$,

$$\widehat{G}(\lambda) = L(\lambda) + [0_{km \times (k-1)m} \quad (v \otimes I_m)C(\lambda I_n - A)^{-1}B].$$

Let $D(\lambda)$ be the polynomial part of $G(\lambda)$. Since $L(\lambda)$ belongs to $\mathbb{M}_1(D)$, $L(\lambda)(\Phi_k(\lambda) \otimes I_m) = v \otimes D(\lambda) = (v \otimes I_m)D(\lambda)$. Therefore, we obtain

$$\begin{aligned} \widehat{G}(\lambda)(\Phi_k(\lambda) \otimes I_m) &= (L(\lambda) + [0_{km \times (k-1)m} \quad (v \otimes I_m)C(\lambda I_n - A)^{-1}B])(\Phi_k(\lambda) \otimes I_m) \\ &= (v \otimes I_m)D(\lambda) + (v \otimes I_m)C(\lambda I_n - A)^{-1}B \\ &= (v \otimes I_m)G(\lambda). \end{aligned}$$

□

Remark 5.5. Since $\mathcal{L}(\lambda)$ is a strong linearization of the rational matrix $G(\lambda)$ we have, by Definition 2.3, that there are unimodular matrices $U(\lambda), V(\lambda) \in \mathbb{F}[\lambda]^{km \times km}$ such that

$$U(\lambda)\widehat{G}(\lambda)V(\lambda) = \text{diag}(G(\lambda), I_{(k-1)m}). \quad (16)$$

Thus, if we consider a finite eigenvalue λ_0 of $G(\lambda)$ then it is also an eigenvalue of the transfer function $\widehat{G}(\lambda)$ and

$$\dim \mathcal{N}_r(G(\lambda_0)) = \dim \mathcal{N}_r(\widehat{G}(\lambda_0)). \quad (17)$$

By [4, Theorem 3.10], $\det(\lambda_0 I_n - A) \neq 0$. Thus, by Proposition 5.1,

$$\dim \mathcal{N}_r(\widehat{G}(\lambda_0)) = \dim \mathcal{N}_r(\mathcal{L}(\lambda_0)). \quad (18)$$

By (16) and Proposition 5.2, we have the same equalities for the dimensions of the left null-spaces, i.e.,

$$\dim \mathcal{N}_\ell(G(\lambda_0)) = \dim \mathcal{N}_\ell(\widehat{G}(\lambda_0)) \quad \text{and} \quad \dim \mathcal{N}_\ell(\widehat{G}(\lambda_0)) = \dim \mathcal{N}_\ell(\mathcal{L}(\lambda_0)). \quad (19)$$

Moreover, notice that since $G(\lambda)$ is square, $\dim \mathcal{N}_r(G(\lambda_0)) = \dim \mathcal{N}_\ell(G(\lambda_0))$. □

A consequence of Lemma 5.4 is that we can recover very easily right eigenvectors of a rational matrix $G(\lambda)$ from the eigenvectors of the transfer function $\widehat{G}(\lambda)$ of any \mathbb{M}_1 -strong linearization of $G(\lambda)$. We state that in Theorem 5.6, and we emphasize that this result is in the spirit of the one presented in [15, Proposition 3.1] for polynomial matrices $P(\lambda)$ and their strong linearizations in $\mathbb{M}_1(P)$.

Theorem 5.6. *Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix with polynomial part of degree $k \geq 2$, and let $\widehat{G}(\lambda)$ be the transfer function of an \mathbb{M}_1 -strong linearization*

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} \quad XB \\ \hline -(v \otimes I_m)CY & L(\lambda) \end{array} \right]$$

of $G(\lambda)$. Let λ_0 be a finite eigenvalue of $G(\lambda)$. Then, $u \in \mathcal{N}_r(G(\lambda_0))$ if and only if $\Phi_k(\lambda_0) \otimes u \in \mathcal{N}_r(\widehat{G}(\lambda_0))$. Moreover, $\{u_1, \dots, u_t\}$ is a basis of $\mathcal{N}_r(G(\lambda_0))$ if and only if $\{\Phi_k(\lambda_0) \otimes u_1, \dots, \Phi_k(\lambda_0) \otimes u_t\}$ is a basis of $\mathcal{N}_r(\widehat{G}(\lambda_0))$.

Proof. By Lemma 5.4, $\widehat{G}(\lambda_0)(\Phi_k(\lambda_0) \otimes I_m) = v \otimes G(\lambda_0)$. Thus, it is easy to see that $u \in \mathcal{N}_r(G(\lambda_0))$ if and only if $\Phi_k(\lambda_0) \otimes u \in \mathcal{N}_r(\widehat{G}(\lambda_0))$. Consider $\{u_1, \dots, u_t\}$ a basis of $\mathcal{N}_r(G(\lambda_0))$. Therefore, as $\dim \mathcal{N}_r(G(\lambda_0)) = \dim \mathcal{N}_r(\widehat{G}(\lambda_0))$, an immediate linear independence argument proves that $\{\Phi_k(\lambda_0) \otimes u_1, \dots, \Phi_k(\lambda_0) \otimes u_t\}$ is a basis of $\mathcal{N}_r(\widehat{G}(\lambda_0))$, and conversely. \square

In addition, by using Proposition 5.1, we can recover the right eigenvectors of the transfer function $\widehat{G}(\lambda)$ from the right eigenvectors of the linearization $\mathcal{L}(\lambda)$, and conversely. In particular, if $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix} \right)$ is a solution of the polynomial eigenvalue problem $\mathcal{L}(\lambda)z = 0$ such that $\det(\lambda_0 I_n - A) \neq 0$, then (λ_0, x_0) is a solution of the rational eigenvalue problem $\widehat{G}(\lambda)x = 0$.

In what follows, if we have a vector $\begin{bmatrix} y \\ x \end{bmatrix}$, with $y \in \overline{\mathbb{F}}^{n \times 1}$ and $x \in \overline{\mathbb{F}}^{km \times 1}$, we will consider the vector x partitioned as $x = [x^{(1)} \quad x^{(2)} \quad \dots \quad x^{(k)}]^T$ with $x^{(j)} \in \overline{\mathbb{F}}^{m \times 1}$ for $j = 1, \dots, k$. Recall also in Theorem 5.7 that, as we have explained in Remark 5.5, if $\lambda_0 \in \overline{\mathbb{F}}$ is a finite eigenvalue of $G(\lambda)$ then $\det(\lambda_0 I_n - A) \neq 0$. However, if λ_0 is an eigenvalue of $\mathcal{L}(\lambda)$, then, according to [4, Theorem 3.10], λ_0 might be a zero of $G(\lambda)$ that is simultaneously a pole and, therefore, $\det(\lambda_0 I_n - A) = 0$, and λ_0 is not an eigenvalue of $G(\lambda)$. This is the reason why the condition $\det(\lambda_0 I_n - A) \neq 0$ is assumed in parts a) and b) of Theorem 5.7.

Theorem 5.7. (Recovery of right eigenvectors from \mathbb{M}_1 -strong linearizations)
Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix with polynomial part of degree $k \geq 2$, and let

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} \quad XB \\ \hline -(v \otimes I_m)CY & L(\lambda) \end{array} \right]$$

be an \mathbb{M}_1 -strong linearization of $G(\lambda)$.

- a) *If $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix} \right)$ is a solution of the LEP $\mathcal{L}(\lambda)z = 0$ such that $\det(\lambda_0 I_n - A) \neq 0$, then $(\lambda_0, x_0^{(k)})$ is a solution of the REP $G(\lambda)x = 0$.*
- b) *Moreover, if $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_r(\mathcal{L}(\lambda_0))$, with $\det(\lambda_0 I_n - A) \neq 0$, then $\{x_1^{(k)}, \dots, x_t^{(k)}\}$ is a basis of $\mathcal{N}_r(G(\lambda_0))$.*

- c) Conversely, if (λ_0, u_0) is a solution of the REP $G(\lambda)x = 0$, $x_0 = \Phi_k(\lambda_0) \otimes u_0$ and y_0 is defined as the unique solution of $(\lambda_0 I_n - A)Yy_0 + Bu_0 = 0$, then $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix}\right)$ is a solution of the LEP $\mathcal{L}(\lambda)z = 0$.
- d) Moreover, if $\{u_1, \dots, u_t\}$ is a basis of $\mathcal{N}_r(G(\lambda_0))$ and, for $i = 1, \dots, t$, $x_i = \Phi_k(\lambda_0) \otimes u_i$ and y_i is defined as the unique solution of $(\lambda_0 I_n - A)Yy_i + Bu_i = 0$, then $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_r(\mathcal{L}(\lambda_0))$.

Proof. By Proposition 5.1, if $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix}\right)$ is a solution of the LEP $\mathcal{L}(\lambda)z = 0$ such that $\det(\lambda_0 I_n - A) \neq 0$, then (λ_0, x_0) is a solution of the REP $\widehat{G}(\lambda)x = 0$, where $\widehat{G}(\lambda)$ is the transfer function matrix of $\mathcal{L}(\lambda)$. By Theorem 5.6, x_0 has the form $x_0 = \Phi_k(\lambda_0) \otimes u$ for some $u \in \mathcal{N}_r(G(\lambda_0))$. Since $\phi_0(\lambda) = 1$ we have that $u = x_0^{(k)}$, which proves a). The converse c) is proved analogously. The implications b) and d) are consequences of (17), (18), basic arguments of linear independence, and the fact that $\mathcal{L}(\lambda_0) \begin{bmatrix} y_0 \\ x_0 \end{bmatrix} = 0$ if and only if $(\lambda_0 I_n - A)Yy_0 + XBx_0^{(k)} = 0$ and $\widehat{G}(\lambda_0)x_0 = 0$. \square

Next, we pay attention to the recovery of left eigenvectors.

Theorem 5.8. (Recovery of left eigenvectors from \mathbb{M}_1 -strong linearizations)

Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix with polynomial part of degree $k \geq 2$, let

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} \quad XB \\ \hline -(v \otimes I_m)CY & L(\lambda) \end{array} \right]$$

be an \mathbb{M}_1 -strong linearization of $G(\lambda)$, and let $\widehat{G}(\lambda)$ be the transfer function of $\mathcal{L}(\lambda)$.

- a) If $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix}\right)$ is a solution of the LEP $z^T \mathcal{L}(\lambda) = 0$ such that $\det(\lambda_0 I_n - A) \neq 0$, then $(\lambda_0, (v^T \otimes I_m)x_0)$ is a solution of the REP $x^T G(\lambda) = 0$.
- b) Moreover, if $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_\ell(\mathcal{L}(\lambda_0))$, with $\det(\lambda_0 I_n - A) \neq 0$, then $\{(v^T \otimes I_m)x_1, \dots, (v^T \otimes I_m)x_t\}$ is a basis of $\mathcal{N}_\ell(G(\lambda_0))$.
- c) Conversely, if (λ_0, u_0) is a solution of the REP $x^T G(\lambda) = 0$, then there exists $x_0 \in \mathcal{N}_\ell(\widehat{G}(\lambda_0))$ such that $u_0 = (v^T \otimes I_m)x_0$ and if y_0 is defined as the unique solution of $y_0^T X(\lambda_0 I_n - A) - u_0^T C = 0$, then $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix}\right)$ is a solution of the LEP $z^T \mathcal{L}(\lambda) = 0$.
- d) Moreover, if $\{u_1, \dots, u_t\}$ is a basis of $\mathcal{N}_\ell(G(\lambda_0))$ then, for $i = 1, \dots, t$, there exists $x_i \in \mathcal{N}_\ell(\widehat{G}(\lambda_0))$ such that $u_i = (v^T \otimes I_m)x_i$ and if y_i is defined as the unique solution of $y_i^T X(\lambda_0 I_n - A) - u_i^T C = 0$, then $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_\ell(\mathcal{L}(\lambda_0))$.

Proof. We consider the transfer function of $\mathcal{L}(\lambda)$, $\widehat{G}(\lambda) = L(\lambda) + [0_{km \times (k-1)m} \quad (v \otimes I_m)C(\lambda I_n - A)^{-1}B]$. If $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix}\right)$ is a solution of the LEP $z^T \mathcal{L}(\lambda) = 0$ such that $\det(\lambda_0 I_n - A) \neq 0$, by using Proposition 5.2 a) applied to $\mathcal{L}(\lambda)$, we get

$$x_0^T \widehat{G}(\lambda_0) = x_0^T L(\lambda_0) + [0_{1 \times (k-1)m} \quad x_0^T (v \otimes I_m)C(\lambda_0 I_n - A)^{-1}B] = 0, \quad (20)$$

where $x_0 \neq 0$ since (λ_0, x_0) is a solution of the REP $x^T \widehat{G}(\lambda) = 0$ ³. In addition, by Lemma 5.4, $x_0^T \widehat{G}(\lambda_0)(\Phi_k(\lambda_0) \otimes I_m) = x_0^T (v \otimes I_m)G(\lambda_0)$. Therefore $x_0^T (v \otimes I_m)G(\lambda_0) = 0$. To see that $(v^T \otimes I_m)x_0$ is a left eigenvector of $G(\lambda_0)$, we only need to prove that $x_0^T (v \otimes I_m) \neq 0$. Let us suppose that $x_0^T (v \otimes I_m) = 0$, and let us get a contradiction. In this case $x_0^T (v \otimes I_m)C(\lambda_0 I_n - A)^{-1}B = 0$ and, therefore, $x_0^T L(\lambda_0) = x_0^T [v \otimes I_m \quad H]F_{\Phi}^D(\lambda_0) = 0$ by (20). We call $w^T = x_0^T [v \otimes I_m \quad H]$ and we consider w partitioned as $w = (w_i)_{i=1}^k$ with $w_i \in \overline{\mathbb{F}}^{m \times 1}$. We have that $w_1^T = x_0^T (v \otimes I_m) = 0$. Therefore $[0 \quad w_2^T \quad \cdots \quad w_k^T]F_{\Phi}^D(\lambda_0) = 0$. This implies $-\alpha_{k-2}w_2^T = 0$ and thus $w_2 = 0$, since $\alpha_{k-2} \neq 0$. Therefore $[0 \quad 0 \quad w_3^T \quad \cdots \quad w_k^T]F_{\Phi}^D(\lambda_0) = 0$ and $w_3 = 0$. Proceeding in this way it is easy to prove that $w_i = 0$ for $i = 2, \dots, k$. Thus $x_0^T [v \otimes I_m \quad H] = 0$ which is a contradiction because $[v \otimes I_m \quad H]$ is assumed to be regular and $x_0 \neq 0$. This proves a).

The implication b) is proved as follows. From part a), the vectors $(v^T \otimes I_m)x_1, \dots, (v^T \otimes I_m)x_t$ belong to $\mathcal{N}_{\ell}(G(\lambda_0))$. Therefore, as a consequence of (19), if we prove that $\{(v^T \otimes I_m)x_1, \dots, (v^T \otimes I_m)x_t\}$ is linearly independent, then b) is proved. For this purpose, let $\alpha_1, \dots, \alpha_t \in \overline{\mathbb{F}}$ be arbitrary scalars such that at least one is different from zero. Thus $0 \neq \begin{bmatrix} \alpha_1 y_1 + \cdots + \alpha_t y_t \\ \alpha_1 x_1 + \cdots + \alpha_t x_t \end{bmatrix} \in \mathcal{N}_{\ell}(\mathcal{L}(\lambda_0))$, and, from part a), $x = (v^T \otimes I_m)(\alpha_1 x_1 + \cdots + \alpha_t x_t) \neq 0$ and $x \in \mathcal{N}_{\ell}(G(\lambda_0))$.

For proving c), we prove first that there exists a basis of $\mathcal{N}_{\ell}(G(\lambda_0))$ of the form $\{(v^T \otimes I_m)x_1, \dots, (v^T \otimes I_m)x_t\}$, where $\{x_1, \dots, x_t\}$ is a basis of $\mathcal{N}_{\ell}(\widehat{G}(\lambda_0))$. To this purpose, let $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ be a basis of $\mathcal{N}_{\ell}(\mathcal{L}(\lambda_0))$. Then, Proposition 5.2 b) applied to $\mathcal{L}(\lambda)$ implies that $\{x_1, \dots, x_t\}$ is a basis of $\mathcal{N}_{\ell}(\widehat{G}(\lambda_0))$ and Theorem 5.8 b) that $\{(v^T \otimes I_m)x_1, \dots, (v^T \otimes I_m)x_t\}$ is a basis of $\mathcal{N}_{\ell}(G(\lambda_0))$. Then, if (λ_0, u_0) is a solution of the REP $x^T G(\lambda) = 0$, u_0 can be written as $u_0 = (v^T \otimes I_m) \sum_{i=1}^t a_i x_i$ with $a_i \in \overline{\mathbb{F}}$, and we define $x_0 = \sum_{i=1}^t a_i x_i \in \mathcal{N}_{\ell}(\widehat{G}(\lambda_0))$.

Finally, Proposition 5.2 c) applied to the solution (λ_0, x_0) of the REP $x^T \widehat{G}(\lambda) = 0$ and to $\mathcal{L}(\lambda)$, and the fact that $\det(\lambda_0 I_n - A) \neq 0$ imply that if y_0 is the unique solution of $y_0^T X(\lambda_0 I_n - A) - x_0^T (v \otimes I_m)C = 0$, which is equivalent to $y_0^T X(\lambda_0 I_n - A) - u_0^T C = 0$, then $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix}\right)$ is a solution of the LEP $z^T \mathcal{L}(\lambda) = 0$.

Finally, the proof of d) proceeds as follows. From part c), we obtain that the vectors x_1, \dots, x_t satisfying $u_i = (v^T \otimes I_m)x_i$ exist, and that the vectors $\begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix}$ belong to $\mathcal{N}_{\ell}(\mathcal{L}(\lambda_0))$. Therefore, taking (19) into account, it only remains to prove that

³With the notation of Proposition 5.2, it is easy to see that $[y_0^T \quad x_0^T]P(\lambda_0) = 0$ if and only if $y_0^T A(\lambda_0) - x_0^T C(\lambda_0) = 0$ and $x_0^T G(\lambda_0) = 0$. Thus, $x_0 = 0$ and $\det A(\lambda_0) \neq 0$ imply $y_0 = 0$. Therefore, any left eigenvector of $P(\lambda)$ corresponding to the finite eigenvalue λ_0 must have $x_0 \neq 0$.

$\begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix}$ are linearly independent. This is easily proved by contradiction:

If $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is linearly dependent, then $\{x_1, \dots, x_t\}$ is linearly dependent, and $\{u_1, \dots, u_t\}$ is linearly dependent, which is a contradiction since $\{u_1, \dots, u_t\}$ is a basis. \square

Remark 5.9. Analogously to Remark 5.3, if $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ is singular, then the results on null-spaces proved so far in Section 5.2 are valid for any $\lambda_0 \in \overline{\mathbb{F}}$ that satisfies $\det(\lambda_0 I_n - A) \neq 0$. \square

Finally, we study the recovery of the eigenvectors corresponding to the infinite eigenvalue from \mathbb{M}_1 -strong linearizations.

Theorem 5.10. (Recovery of eigenvectors associated to infinity from \mathbb{M}_1 -strong linearizations) Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix with polynomial part of degree $k \geq 2$, let

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} \quad XB \\ \hline -(v \otimes I_m)CY & L(\lambda) \end{array} \right]$$

be an \mathbb{M}_1 -strong linearization of $G(\lambda)$, and let D_k be the leading matrix coefficient of the polynomial part of $G(\lambda)$ as in (13). Then the following statements hold:

a) $\mathcal{N}_r(\text{rev } G(0)) = \mathcal{N}_r(D_k)$ and $x_0 \in \mathcal{N}_r(D_k)$ if and only if $\begin{bmatrix} 0 \\ e_1 \otimes x_0 \end{bmatrix} \in \mathcal{N}_r(\text{rev } \mathcal{L}(0))$.

Moreover, $\{x_1, \dots, x_q\}$ is a basis of $\mathcal{N}_r(\text{rev } G(0))$ if and only if $\left\{ \begin{bmatrix} 0 \\ e_1 \otimes x_1 \end{bmatrix}, \dots, \begin{bmatrix} 0 \\ e_1 \otimes x_q \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_r(\text{rev } \mathcal{L}(0))$.

b) $\mathcal{N}_\ell(\text{rev } G(0)) = \mathcal{N}_\ell(D_k)$ and $\begin{bmatrix} 0 \\ x_0 \end{bmatrix} \in \mathcal{N}_\ell(\text{rev } \mathcal{L}(0))$ if and only if $(v^T \otimes I_m)x_0 \in \mathcal{N}_\ell(D_k)$.

Moreover, $\left\{ \begin{bmatrix} 0 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} 0 \\ x_q \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_\ell(\text{rev } \mathcal{L}(0))$ if and only if $\{(v^T \otimes I_m)x_1, \dots, (v^T \otimes I_m)x_q\}$ is a basis of $\mathcal{N}_\ell(\text{rev } G(0))$.

Proof. Notice that from (12),

$$F_\Phi^D(\lambda) = \lambda \begin{bmatrix} \alpha_{k-1}^{-1} D_k & 0 \\ 0 & I_{(k-1)m} \end{bmatrix} + F_\Phi^D(0).$$

We consider

$$L(\lambda) = [v \otimes I_m \quad H] F_\Phi^D(\lambda) = [\alpha_{k-1}^{-1}(v \otimes D_k) \quad H] \lambda + L(0) =: L_1 \lambda + L_0$$

and let $\widehat{G}(\lambda)$ be the transfer function matrix of $\mathcal{L}(\lambda)$. We have that $\text{rev } \mathcal{L}(0) = \left[\begin{array}{c|c} XY & 0 \\ \hline 0 & L_1 \end{array} \right]$

and $\text{rev } \widehat{G}(0) = \text{rev } L(0) = L_1$. Moreover, $\text{rev } G(0) = \alpha_0^{-1} \alpha_1^{-1} \dots \alpha_{k-1}^{-1} D_k$, that is, the coefficient of λ^k in $D(\lambda)$. Therefore, $\mathcal{N}_r(\text{rev } G(0)) = \mathcal{N}_r(D_k)$, $\mathcal{N}_\ell(\text{rev } G(0)) = \mathcal{N}_\ell(D_k)$ and ∞ is an eigenvalue of $G(\lambda)$ if and only if D_k is singular. In addition, every right

(respectively left) eigenvector w of $\text{rev } \mathcal{L}(0)$ has the form $w = \begin{bmatrix} 0 \\ x_0 \end{bmatrix}$ for some $x_0 \in \mathcal{N}_r(L_1)$ (respectively $x_0 \in \mathcal{N}_\ell(L_1)$). By Lemma 5.4, we have

$$\lambda \widehat{G} \left(\frac{1}{\lambda} \right) \left(\lambda^{k-1} \Phi_k \left(\frac{1}{\lambda} \right) \otimes I_m \right) = v \otimes \lambda^k G \left(\frac{1}{\lambda} \right).$$

Therefore,

$$\text{rev } \widehat{G}(0) (\text{rev } \Phi_k(0) \otimes I_m) = (v \otimes I_m) \text{rev } G(0).$$

Since $\text{rev } \Phi_k(0) = \alpha_0^{-1} \alpha_1^{-1} \cdots \alpha_{k-2}^{-1} e_1$, we obtain

$$\alpha_0^{-1} \alpha_1^{-1} \cdots \alpha_{k-2}^{-1} \text{rev } \widehat{G}(0) (e_1 \otimes I_m) = (v \otimes I_m) \text{rev } G(0).$$

In addition, by (7), there exist unimodular matrices $W_1(\lambda)$ and $W_2(\lambda)$ such that

$$W_1(0) \text{diag} (\text{rev } G(0), I_{(k-1)m}) W_2(0) = \text{rev } \widehat{G}(0),$$

which implies that $\dim \mathcal{N}_r(\text{rev } G(0)) = \dim \mathcal{N}_r(\text{rev } \widehat{G}(0))$ and $\dim \mathcal{N}_\ell(\text{rev } G(0)) = \dim \mathcal{N}_\ell(\text{rev } \widehat{G}(0))$. Finally *a*) and *b*) follow from the results above by using similar arguments to the ones we used in the recovery of eigenvectors associated to finite eigenvalues. \square

5.3. Eigenvectors from \mathbb{M}_2 -strong linearizations

If we proceed analogously as we did with \mathbb{M}_1 -strong linearizations, and we use Lemma 5.11, then we get Theorems 5.12, 5.13 and 5.15 to recover right and left eigenvectors of a rational matrix from those of its \mathbb{M}_2 -strong linearizations. The proofs are essentially the same as those in Section 5.2 by interchanging the roles of left and right eigenvectors, and they are omitted for brevity.

Lemma 5.11. *Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix with polynomial part of degree $k \geq 2$, let*

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & XB(w^T \otimes I_m) \\ \hline \begin{matrix} 0_{(k-1)m \times n} \\ -CY \end{matrix} & L(\lambda) \end{array} \right]$$

be an \mathbb{M}_2 -strong linearization of $G(\lambda)$, and let $\widehat{G}(\lambda)$ be the transfer function of $\mathcal{L}(\lambda)$. Then

$$(\Phi_k(\lambda)^T \otimes I_m) \widehat{G}(\lambda) = w^T \otimes G(\lambda). \quad (21)$$

Theorem 5.12. (Recovery of right eigenvectors from \mathbb{M}_2 -strong linearizations)
Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix with polynomial part of degree $k \geq 2$, let

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & XB(w^T \otimes I_m) \\ \hline \begin{matrix} 0_{(k-1)m \times n} \\ -CY \end{matrix} & L(\lambda) \end{array} \right]$$

be an \mathbb{M}_2 -strong linearization of $G(\lambda)$, and let $\widehat{G}(\lambda)$ be the transfer function of $\mathcal{L}(\lambda)$.

- a) If $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix} \right)$ is a solution of the LEP $\mathcal{L}(\lambda)z = 0$ such that $\det(\lambda_0 I_n - A) \neq 0$ then, $(\lambda_0, (w^T \otimes I_m)x_0)$ is a solution of the REP $G(\lambda)x = 0$.*

- b) Moreover, if $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_r(\mathcal{L}(\lambda_0))$, with $\det(\lambda_0 I_n - A) \neq 0$, then $\{(w^T \otimes I_m)x_1, \dots, (w^T \otimes I_m)x_t\}$ is a basis of $\mathcal{N}_r(G(\lambda_0))$.
- c) Conversely, if (λ_0, u_0) is a solution of the REP $G(\lambda)x = 0$, then there exists $x_0 \in \mathcal{N}_r(\widehat{G}(\lambda_0))$ such that $u_0 = (w^T \otimes I_m)x_0$ and if y_0 is defined as the unique solution of $(\lambda_0 I_n - A)Yy_0 + Bu_0 = 0$, then $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix} \right)$ is a solution of the LEP $\mathcal{L}(\lambda)z = 0$.
- d) Moreover, if $\{u_1, \dots, u_t\}$ is a basis of $\mathcal{N}_r(G(\lambda_0))$ then, for $i = 1, \dots, t$, there exists $x_i \in \mathcal{N}_r(\widehat{G}(\lambda_0))$ such that $u_i = (w^T \otimes I_m)x_i$ and if y_i is defined as the unique solution of $(\lambda_0 I_n - A)Yy_i + Bu_i = 0$, then $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_r(\mathcal{L}(\lambda_0))$.

Theorem 5.13. (Recovery of left eigenvectors from \mathbb{M}_2 -strong linearizations) Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix with polynomial part of degree $k \geq 2$, and let

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & XB(w^T \otimes I_m) \\ \hline \begin{array}{c} 0_{(k-1)m \times n} \\ -CY \end{array} & L(\lambda) \end{array} \right]$$

be an \mathbb{M}_2 -strong linearization of $G(\lambda)$.

- a) If $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix} \right)$ is a solution of the LEP $z^T \mathcal{L}(\lambda) = 0$ such that $\det(\lambda_0 I_n - A) \neq 0$, then $(\lambda_0, x_0^{(k)})$ is a solution of the REP $x^T G(\lambda) = 0$.
- b) Moreover, if $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_\ell(\mathcal{L}(\lambda_0))$, with $\det(\lambda_0 I_n - A) \neq 0$, then $\{x_1^{(k)}, \dots, x_t^{(k)}\}$ is a basis of $\mathcal{N}_\ell(G(\lambda_0))$.
- c) Conversely, if (λ_0, u_0) is a solution of the REP $x^T G(\lambda) = 0$, $x_0 = \Phi_k(\lambda_0) \otimes u_0$ and y_0 is defined as the unique solution of $y_0^T X(\lambda_0 I_n - A) - u_0^T C = 0$, then $\left(\lambda_0, \begin{bmatrix} y_0 \\ x_0 \end{bmatrix} \right)$ is a solution of the LEP $z^T \mathcal{L}(\lambda) = 0$.
- d) Moreover, if $\{u_1, \dots, u_t\}$ is a basis of $\mathcal{N}_\ell(G(\lambda_0))$ and, for $i = 1, \dots, t$, $x_i = \Phi_k(\lambda_0) \otimes u_i$ and y_i is defined as the unique solution of $y_i^T X(\lambda_0 I_n - A) - u_i^T C = 0$, then $\left\{ \begin{bmatrix} y_1 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} y_t \\ x_t \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_\ell(\mathcal{L}(\lambda_0))$.

Remark 5.14. Analogously to Remarks 5.3 and 5.9, if $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ is singular, then the results on null-spaces in Theorems 5.12 and 5.13 hold for any $\lambda_0 \in \overline{\mathbb{F}}$ such that $\det(\lambda_0 I_n - A) \neq 0$. \square

Theorem 5.15. (Recovery of eigenvectors associated to infinity from \mathbb{M}_2 -strong linearizations) Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix with polynomial part of degree $k \geq 2$, let

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & XB(w^T \otimes I_m) \\ \hline \begin{array}{c} 0_{(k-1)m \times n} \\ -CY \end{array} & L(\lambda) \end{array} \right]$$

be an \mathbb{M}_2 -strong linearization of $G(\lambda)$, and let D_k be the leading matrix coefficient of the polynomial part of $G(\lambda)$ as in (13). Then the following statements hold:

- a) $\mathcal{N}_r(\text{rev } G(0)) = \mathcal{N}_r(D_k)$ and $\begin{bmatrix} 0 \\ x_0 \end{bmatrix} \in \mathcal{N}_r(\text{rev } \mathcal{L}(0))$ if and only if $(w^T \otimes I_m)x_0 \in \mathcal{N}_r(D_k)$. Moreover, $\left\{ \begin{bmatrix} 0 \\ x_1 \end{bmatrix}, \dots, \begin{bmatrix} 0 \\ x_q \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_r(\text{rev } \mathcal{L}(0))$ if and only if $\{(w^T \otimes I_m)x_1, \dots, (w^T \otimes I_m)x_q\}$ is a basis of $\mathcal{N}_r(\text{rev } G(0))$.
- b) $\mathcal{N}_\ell(\text{rev } G(0)) = \mathcal{N}_\ell(D_k)$ and $x_0 \in \mathcal{N}_\ell(D_k)$ if and only if $\begin{bmatrix} 0 \\ e_1 \otimes x_0 \end{bmatrix} \in \mathcal{N}_\ell(\text{rev } \mathcal{L}(0))$. Moreover, $\{x_1, \dots, x_q\}$ is a basis of $\mathcal{N}_\ell(\text{rev } G(0))$ if and only if $\left\{ \begin{bmatrix} 0 \\ e_1 \otimes x_1 \end{bmatrix}, \dots, \begin{bmatrix} 0 \\ e_1 \otimes x_q \end{bmatrix} \right\}$ is a basis of $\mathcal{N}_\ell(\text{rev } \mathcal{L}(0))$.

6. Symmetric realizations of symmetric rational matrices

In this section and in the next one our aim is to obtain a strong linearization of a symmetric rational matrix $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$, i.e., $G(\lambda)^T = G(\lambda)$, that preserves its symmetric structure. We write $G(\lambda)$ as

$$G(\lambda) = D(\lambda) + G_{sp}(\lambda) \quad (22)$$

with $D(\lambda)$ its polynomial part and $G_{sp}(\lambda)$ its strictly proper part. Since (22) is a unique decomposition we obtain the following result just by taking transposes.

Proposition 6.1. *Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a symmetric rational matrix. Then the matrices $D(\lambda)$ and $G_{sp}(\lambda)$ in (22) are also symmetric.*

Proposition 6.5 is the main result in this section and shows that any symmetric strictly proper rational matrix admits a state-space realization that reveals transparently the symmetry. In order to state concisely Proposition 6.5, we will use the following definition.

Definition 6.2. *Let $G_{sp}(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a symmetric strictly proper rational matrix and let $n = \nu(G_{sp}(\lambda))$ be the least order of $G_{sp}(\lambda)$. A symmetric minimal state-space realization of $G_{sp}(\lambda)$ is an expression of the form*

$$G_{sp}(\lambda) = W(S_1\lambda - S_2)^{-1}W^T$$

where $S_1, S_2 \in \mathbb{F}^{n \times n}$ are symmetric matrices with S_1 nonsingular and $W \in \mathbb{F}^{m \times n}$.

We remark that the realization described in Definition 6.2 is equivalent to [6, Definition 4.44] for a minimal state-space realization. However, in Definition 6.2 we express strictly proper matrices in a form more convenient for the goals of this paper. In particular, we will see in Section 7 that by combining a symmetric minimal state-space realization of the matrix $G_{sp}(\lambda)$ in (22) and a symmetric strong block minimal bases pencil associated to $D(\lambda)$, we can construct symmetric strong linearizations of $G(\lambda)$. The next technical lemma is used in the proof of Proposition 6.5.

Lemma 6.3. *Let $G_{sp}(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a symmetric strictly proper rational matrix and let $G_{sp}(\lambda) = C(\lambda I_n - A)^{-1}B$ be a minimal state-space realization of $G_{sp}(\lambda)$. Then there exists a unique nonsingular and symmetric matrix $S \in \mathbb{F}^{n \times n}$ such that $A^T = S^{-1}AS$ and $C^T = S^{-1}B$.*

Proof. As $G_{sp}(\lambda)$ is symmetric, $G_{sp}(\lambda) = B^T(\lambda I_n - A^T)^{-1}C^T$ is also a minimal state-space realization of $G_{sp}(\lambda)$ since both have the same minimal order n . Therefore, by [19, Proposition 3.3.2], the realizations (A, B, C) and (A^T, C^T, B^T) are similar and there exists a unique nonsingular matrix $S \in \mathbb{F}^{n \times n}$ such that

$$A^T = S^{-1}AS, \quad C^T = S^{-1}B, \quad B^T = CS. \quad (23)$$

The fact that (A, B, C) is a minimal realization of $G_{sp}(\lambda)$ is equivalent to that (A, B) and (A, C) are controllable and observable, respectively (see [29, Chapter 3]). That means that the controllability matrix of (A, B) and the observability matrix of (A, C) , i.e.,

$$\mathcal{C}(A, B) = [B \quad AB \quad A^2B \quad \cdots \quad A^{n-1}B] \quad \text{and} \quad \mathcal{O}(A, C) = \begin{bmatrix} C \\ CA \\ CA^2 \\ \vdots \\ CA^{n-1} \end{bmatrix},$$

have both rank n . From the equalities in (23) it is easy to see that $S^{-1}\mathcal{C}(A, B) = \mathcal{O}(A, C)^T$, and $S^{-T}\mathcal{C}(A, B) = \mathcal{O}(A, C)^T$. As $\mathcal{C}(A, B)$ has full row rank, we deduce that $S = S^T$. \square

Remark 6.4. Notice that the system similarity matrix S between the realizations in Lemma 6.3 is given by $S = \mathcal{O}(A, C)^+\mathcal{C}(A, B)^T = \mathcal{C}(A, B)(\mathcal{O}(A, C)^T)^\dagger$ where $+$ denotes any left inverse and \dagger denotes any right inverse. Notice also that these left and right inverses exist because (A, B, C) is a minimal realization of $G_{sp}(\lambda)$ and that they can be taken to be the Moore–Penrose inverses. Thus S can be efficiently computed when $\mathbb{F} = \mathbb{R}, \mathbb{C}$. \square

Proposition 6.5. *Any symmetric strictly proper rational matrix has a symmetric minimal state-space realization.*

Proof. As said in Section 2, any strictly proper rational matrix $G_{sp}(\lambda)$ admits a minimal state-space realization, that is, $G_{sp}(\lambda) = C(\lambda I_n - A)^{-1}B$ [29]. By Lemma 6.3, there exists a unique nonsingular and symmetric matrix S such that $A^T = S^{-1}AS$ and $C^T = S^{-1}B$. Thus, $G_{sp}(\lambda) = C(\lambda I_n - A)^{-1}SC^T = C(\lambda S^{-1} - S^{-1}A)^{-1}C^T$, and $S^{-1}A$ is symmetric, as $(S^{-1}A)^T = A^T S^{-1} = S^{-1}A$. \square

Remark 6.6. We can construct a symmetric minimal state-space realization of a symmetric strictly proper rational matrix $G_{sp}(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ without previously considering a non-symmetric minimal state-space realization of $G_{sp}(\lambda)$, in contrast to what we have done in the proof of Proposition 6.5. For this purpose we require \mathbb{F} not to be a field of characteristic 2. Let $G_{sp}(\lambda) = G_1\lambda^{-1} + G_2\lambda^{-2} + \cdots$ be the Laurent series of $G_{sp}(\lambda)$, which converges for $|\lambda|$ large enough. Let $n = \nu(G_{sp}(\lambda))$ be the least order of $G_{sp}(\lambda)$. We

consider the block Hankel matrix

$$H_n = \begin{bmatrix} G_1 & G_2 & \cdots & G_n \\ G_2 & G_3 & \cdots & G_{n+1} \\ \vdots & \vdots & \ddots & \vdots \\ G_n & G_{n+1} & \cdots & G_{2n-1} \end{bmatrix} \quad (24)$$

and follow in a symmetric way the three steps of the algorithm in [19, Section 3.4] to get a symmetric minimal state-space realization from the Hankel matrix. Notice that the Hankel matrix is symmetric since $G_{sp}(\lambda)$ is symmetric, which implies $G_i = G_i^T$ for all $i \geq 1$, and $\text{rank } H_n = n$ by [19, Proposition 3.3.2]. Therefore we can write

$$H_n = X \begin{bmatrix} K & 0 \\ 0 & 0 \end{bmatrix} X^T = X \begin{bmatrix} K \\ 0 \end{bmatrix} [I_n \quad 0] X^T$$

with X nonsingular and $K \in \mathbb{F}^{n \times n}$ diagonal (see [25, Theorem 34.1]). Let us denote

$$\Gamma = X \begin{bmatrix} K \\ 0 \end{bmatrix} \quad \text{and} \quad \Lambda = [I_n \quad 0] X^T.$$

We have that $H_n = \Gamma \Lambda$. We write $X = [X_1 \quad X_2]$, where $X_1 = [X_{i1}]_{i=1}^n$ with $X_{i1} \in \mathbb{F}^{m \times n}$ for $i = 1, \dots, n$. Thus

$$\Gamma = \begin{bmatrix} X_{11}K \\ \vdots \\ X_{n1}K \end{bmatrix} \quad \text{and} \quad \Lambda = [X_{11}^T \quad \cdots \quad X_{n1}^T].$$

We define

$$R = \begin{bmatrix} G_2 & G_3 & \cdots & G_{n+1} \\ G_3 & G_4 & \cdots & G_{n+2} \\ \vdots & \vdots & \ddots & \vdots \\ G_{n+1} & G_{n+2} & \cdots & G_{2n} \end{bmatrix}$$

and we set $C = X_{11}K$, $B = X_{11}^T$ and $A = \Gamma^+ R \Lambda^+$, with $\Gamma^+ = [K^{-1} \quad 0] X^{-1}$ and $\Lambda^+ = X^{-T} \begin{bmatrix} I_n \\ 0 \end{bmatrix}$. Thus $A = [K^{-1} \quad 0] X^{-1} R X^{-T} \begin{bmatrix} I_n \\ 0 \end{bmatrix}$ and, by [19, Theorem 3.4.1], (A, B, C) is a minimal realization for $G_{sp}(\lambda)$. Therefore

$$\begin{aligned} G_{sp}(\lambda) &= X_{11}K \left(\lambda I_n - [K^{-1} \quad 0] X^{-1} R X^{-T} \begin{bmatrix} I_n \\ 0 \end{bmatrix} \right)^{-1} X_{11}^T \\ &= X_{11} \left(\lambda K^{-1} - [K^{-1} \quad 0] X^{-1} R X^{-T} \begin{bmatrix} K^{-1} \\ 0 \end{bmatrix} \right)^{-1} X_{11}^T. \end{aligned}$$

Finally we set $W = X_{11}$, $S_1 = K^{-1}$ and $S_2 = [K^{-1} \quad 0] X^{-1} R X^{-T} \begin{bmatrix} K^{-1} \\ 0 \end{bmatrix}$, and we obtain a symmetric minimal state-space realization of $G_{sp}(\lambda)$.

In the particular case $G_{sp}(\lambda) \in \mathbb{R}(\lambda)^{m \times m}$, which is of significant importance in applications, the Hankel matrix H_n is real and symmetric. Therefore, we can write

$$H_n = P \begin{bmatrix} K & 0 \\ 0 & 0 \end{bmatrix} P^T$$

with P orthogonal, i.e., $P^{-1} = P^T$, and K a diagonal matrix that has the eigenvalues different from zero of H_n at the diagonal elements. In this case, let $P = \begin{bmatrix} P_1 & P_2 \end{bmatrix}$, where $P_1 = [P_{i1}]_{i=1}^n$ with $P_{i1} \in \mathbb{F}^{m \times n}$ for $i = 1, \dots, n$. Then we obtain $G_{sp}(\lambda) = P_{11}(\lambda K^{-1} - K^{-1}P_1^T R P_1 K^{-1})^{-1}P_{11}^T$. That is, $G_{sp}(\lambda)$ has a symmetric minimal state-space realization $G_{sp}(\lambda) = W(\lambda S_1 - S_2)^{-1}W^T$ where $W = P_{11}$, $S_1 = K^{-1}$ and $S_2 = K^{-1}P_1^T R P_1 K^{-1}$. \square

From Proposition 6.5 and Remark 6.6 we know how to write the strictly proper part $G_{sp}(\lambda)$ of a symmetric rational matrix $G(\lambda)$ as a symmetric minimal state-space realization with or without having in advance a particular non-symmetric minimal state-space realization of $G_{sp}(\lambda)$. Moreover, it is worth to emphasize that in many applications of symmetric REPs, this can be done very easily from the data of the model without any computational cost (see [4, Section 5.3] or [30, Section 4]).

7. Symmetric strong linearizations for symmetric rational matrices

In this section symmetric strong linearizations for symmetric rational matrices will be constructed. We start with Example 7.1 in which we construct a symmetric strong linearization of a symmetric rational matrix when the polynomial part has odd degree. We will use Proposition 6.5 and a particular symmetric strong block minimal bases pencil associated to its polynomial part with sharp degree. After that, we present symmetric strong linearizations for symmetric rational matrices in which the polynomial part may have even or odd degree but the leading coefficient must be nonsingular. In order to get these results, we need to study symmetric strong linearizations in the polynomial case.

Example 7.1. Let $G(\lambda) = D(\lambda) + G_{sp}(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a symmetric rational matrix. Consider the polynomial part $D(\lambda)$ written in terms of the monomial basis $D(\lambda) = D_k \lambda^k + D_{k-1} \lambda^{k-1} + \dots + D_0 \in \mathbb{F}[\lambda]^{m \times m}$, with $k > 1$ and $D_k \neq 0$, and the matrices

$$L_p(\lambda) = \begin{bmatrix} -1 & \lambda & & & \\ & -1 & \lambda & & \\ & & \ddots & \ddots & \\ & & & -1 & \lambda \end{bmatrix} \in \mathbb{F}[\lambda]^{p \times (p+1)}, \quad (25)$$

and

$$\Lambda_p(\lambda)^T = [\lambda^p \quad \dots \quad \lambda \quad 1] \in \mathbb{F}[\lambda]^{1 \times (p+1)}. \quad (26)$$

A block Kronecker linearization of $D(\lambda)$ is a pencil

$$L(\lambda) = \left[\underbrace{\begin{array}{c|c} M(\lambda) & L_\eta(\lambda)^T \otimes I_m \\ \hline L_\varepsilon(\lambda) \otimes I_m & 0 \end{array}}_{(\varepsilon+1)m} \right] \left. \begin{array}{l} \vphantom{L(\lambda)} \\ \vphantom{L(\lambda)} \\ \vphantom{L(\lambda)} \end{array} \right\} \begin{array}{l} (\eta+1)m \\ \varepsilon m \end{array} \quad (27)$$

such that $D(\lambda) = (\Lambda_\eta(\lambda)^T \otimes I_m) M(\lambda) (\Lambda_\varepsilon(\lambda) \otimes I_m)$ (see [13, Definition 4.1]). Recall that block Kronecker linearizations are particular cases of strong block minimal bases pencils [13]. If the polynomial part $D(\lambda)$ has odd degree $k = 2q + 1$ we can consider the symmetric block Kronecker linearization in which

$$M(\lambda) = \begin{bmatrix} D_{2q+1}\lambda + D_{2q} & & & \\ & D_{2q-1}\lambda + D_{2q-2} & & \\ & & \ddots & \\ & & & D_1\lambda + D_0 \end{bmatrix}$$

and $\varepsilon = \eta = q$. Proposition 6.5 allows us to write $G_{sp}(\lambda) = W(\lambda S_1 - S_2)^{-1} W^T$ with S_1 and S_2 symmetric and S_1 nonsingular. Applying [4, Theorem 5.11] with $Y = -S_1 X^T$ for any nonsingular matrix $X \in \mathbb{F}^{n \times n}$, $C = W S_1^{-1}$, $A = S_2 S_1^{-1}$, $B = W^T$, and $\widehat{K}_1 = \widehat{K}_2 = e_{q+1}^T \otimes I_m$, we obtain that the linear polynomial matrix

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c|c} X(S_2 - \lambda S_1)X^T & 0 & XW^T & 0 \\ \hline 0 & M(\lambda) & & L_q(\lambda)^T \otimes I_m \\ WX^T & & & \\ \hline 0 & L_q(\lambda) \otimes I_m & & 0 \end{array} \right]$$

is a symmetric strong linearization of $G(\lambda)$.

Remark 7.2. The approach in Example 7.1 can be extended to other symmetric strong block minimal bases pencils of the symmetric polynomial part $D(\lambda)$ of $G(\lambda) = G(\lambda)^T$ to construct other symmetric strong linearizations of $G(\lambda)$, as long as $D(\lambda)$ has odd-degree. See, for instance, the pencils considered in [14]. However, the linearization in Example 7.1 is particularly simple and, in view of the results in [7], we expect that it will have favourable numerical properties. \square

Let $P(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ be a polynomial matrix of degree k . A $km \times km$ pencil $L(\lambda)$ is called *block-symmetric* if $L(\lambda) = L(\lambda)^{\mathcal{B}}$, where $L(\lambda)$ is viewed as a block partitioned pencil with $k \times k$ blocks each of them of size $m \times m$. Notice that a pencil $L(\lambda)$ satisfies $L(\lambda)(\Phi_k(\lambda) \otimes I_m) = v \otimes P(\lambda)$ for some vector $v \in \mathbb{F}^k$ if and only if $L(\lambda)^{\mathcal{B}}$ satisfies $(\Phi_k(\lambda)^T \otimes I_m)L(\lambda)^{\mathcal{B}} = v^T \otimes P(\lambda)$. Thus, if $L(\lambda) \in \mathbb{M}_1(P)$ is block-symmetric, then $L(\lambda) \in \mathbb{M}_1(P) \cap \mathbb{M}_2(P)$. This intersection space was introduced in [15], it is called *double generalized ansatz space*, and it is denoted by

$$\mathbb{DM}(P) = \mathbb{M}_1(P) \cap \mathbb{M}_2(P).$$

If $\{\phi_j(\lambda)\}_{j=0}^\infty$ is the monomial basis, the space $\mathbb{DM}(P)$ is denoted $\mathbb{DL}(P)$ and was introduced originally in [26]. In [15, Corollary 6] it is shown that if a pencil $L(\lambda)$ belongs to $\mathbb{DM}(P)$ then its right and left ansatz vectors are the same, which is called simply *ansatz vector*, and that

$$\mathbb{DM}(P) = \{L(\lambda) \in \mathbb{M}_1(P) : L(\lambda) = L(\lambda)^{\mathcal{B}}\}.$$

In fact, if $P(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ is a symmetric polynomial matrix we obtain that any pencil in $\mathbb{DM}(P)$ must be symmetric. This result is not in [15], and we state it in Theorem 7.5. For its proof, we use Lemmas 7.3 and 7.4.

Let $P(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ be a symmetric polynomial matrix, and let us define the set

$$\mathbb{S}(P) = \{L(\lambda) \in \mathbb{M}_1(P) : L(\lambda) = L(\lambda)^T\}.$$

The elements in $\mathbb{S}(P)$ are in $\mathbb{DM}(P)$ because if $L(\lambda) = [v \otimes I_m \quad H]F_{\Phi}^P(\lambda) \in \mathbb{S}(P)$ then $L(\lambda)^T = F_{\Phi}^P(\lambda)^{\mathcal{B}} \begin{bmatrix} v^T \otimes I_m \\ H^T \end{bmatrix} \in \mathbb{M}_2(P)$, since in the case $P(\lambda)$ is symmetric $F_{\Phi}^P(\lambda)^T = F_{\Phi}^P(\lambda)^{\mathcal{B}}$, and $L(\lambda) = L(\lambda)^T$. Moreover, Theorem 7.5 shows that $\mathbb{S}(P)$ and $\mathbb{DM}(P)$ are equal.

Theorem 7.5. *Let $P(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ be a symmetric polynomial matrix of degree $k \geq 2$. Then*

$$\mathbb{DM}(P) = \mathbb{S}(P).$$

Proof. We have already seen that $\mathbb{S}(P) \subseteq \mathbb{DM}(P)$. To see the other inclusion we only have to use Lemma 7.4 and [15, Corollary 6], and notice that if $L(\lambda) \in \mathbb{DM}(P)$ with $P(\lambda)$ symmetric then

$$L(\lambda) = [v \otimes I_m \quad H]F_{\Phi}^P(\lambda) = F_{\Phi}^P(\lambda)^{\mathcal{B}} \begin{bmatrix} v^T \otimes I_m \\ H^{\mathcal{B}} \end{bmatrix},$$

and

$$L(\lambda)^T = [v \otimes I_m \quad (H^{\mathcal{B}})^T]F_{\Phi}^P(\lambda) = F_{\Phi}^P(\lambda)^{\mathcal{B}} \begin{bmatrix} v^T \otimes I_m \\ H^T \end{bmatrix},$$

which implies that $L(\lambda)^T \in \mathbb{DM}(P)$ and that $L(\lambda)$ and $L(\lambda)^T$ have the same ansatz vector. So, by Lemma 7.4, $L(\lambda) = L(\lambda)^T$ and $L(\lambda) \in \mathbb{S}(P)$. \square

Therefore, if $P(\lambda)$ is a symmetric polynomial matrix all the pencils in $\mathbb{DM}(P)$ are also symmetric. In order to find linearizations in $\mathbb{DM}(P)$ we have to consider only regular polynomials $P(\lambda)$ because by [15, Theorem 7] if $P(\lambda)$ is a singular polynomial matrix then none of the pencils in $\mathbb{DM}(P)$ is a linearization for $P(\lambda)$.

In Theorem 7.9, we construct symmetric strong linearizations for a symmetric rational matrix from a particular symmetric strong linearization of its polynomial part $D(\lambda)$ when the leading coefficient D_k of $D(\lambda)$ is nonsingular. This particular strong linearization is the pencil in $\mathbb{DM}(D)$ with ansatz vector e_k , i.e., the last vector in the canonical basis of \mathbb{F}^k . Some properties of this pencil are studied in Lemma 7.6 and Corollary 7.7.

Lemma 7.6. *Let $D(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ be a polynomial matrix with degree $k \geq 2$ and let $L(\lambda) = [e_k \otimes I_m \quad H]F_{\Phi}^D(\lambda) \in \mathbb{DM}(D)$. Then $[e_k \otimes I_m \quad H]$ is nonsingular if and only if the leading matrix coefficient D_k of $D(\lambda)$ is nonsingular.*

Proof. Let $L(\lambda) = [e_k \otimes I_m \quad H]F_{\Phi}^D(\lambda) \in \mathbb{DM}(D)$. We write, by using (28) and [15, Corollary 6],

$$\begin{aligned} L(\lambda) &= [e_k \otimes I_m \quad H]F_{\Phi}^D(\lambda) = [e_k \otimes \alpha_{k-1}^{-1}D_k \quad H]\lambda + [e_k \otimes I_m \quad H]F_{\Phi}^D(0) \\ &= \begin{bmatrix} e_k^T \otimes \alpha_{k-1}^{-1}D_k \\ H^{\mathcal{B}} \end{bmatrix} \lambda + F_{\Phi}^D(0)^{\mathcal{B}} \begin{bmatrix} e_k^T \otimes I_m \\ H^{\mathcal{B}} \end{bmatrix}. \end{aligned}$$

Then $H = \begin{bmatrix} 0_{m \times (k-2)m} & \alpha_{k-1}^{-1} D_k \\ & H' \end{bmatrix}$ for some $(k-1)m \times (k-1)m$ block symmetric matrix H' . Let $H' = [H'_{ij}]_{i,j=1}^{k-1}$ with $H'_{ij} \in \mathbb{F}^{m \times m}$. If we calculate the first block row and block column of the product $[e_k \otimes I_m \quad H]F_{\Phi}^D(0)$ we obtain

$$\begin{bmatrix} 0 & 0 & \cdots & 0 & -\frac{\alpha_0}{\alpha_{k-1}} D_k & -\frac{\beta_0}{\alpha_{k-1}} D_k \\ -\alpha_{k-2} H'_{11} & * & \cdots & * & * & * \\ -\alpha_{k-2} H'_{21} & * & \cdots & * & * & * \\ \vdots & & & & & \\ -\alpha_{k-2} H'_{(k-2)1} & * & \cdots & * & * & * \\ -\frac{\beta_{k-1}}{\alpha_{k-1}} D_k + D_{k-1} - \alpha_{k-2} H'_{(k-1)1} & * & \cdots & * & * & * \end{bmatrix}.$$

Since $[e_k \otimes I_m \quad H]F_{\Phi}^D(0)$ is block symmetric we obtain

$$H'_{1i} = H'_{i1} = 0 \text{ for } i = 1, \dots, k-3 \quad (29)$$

and

$$-\alpha_{k-2} H'_{(k-2)1} = -\frac{\alpha_0}{\alpha_{k-1}} D_k.$$

Thus,

$$H'_{(k-2)1} = H'_{1(k-2)} = \frac{\alpha_0}{\alpha_{k-1} \alpha_{k-2}} D_k. \quad (30)$$

Using (29) and (30) and calculating the second block row and block column of the product $[e_k \otimes I_m \quad H]F_{\Phi}^D(0)$ as before, we obtain

$$H'_{2i} = H'_{i2} = 0 \text{ for } i = 1, \dots, k-4$$

and

$$-\alpha_{k-3} H'_{(k-3)2} = -\alpha_1 H'_{1(k-2)}.$$

Thus,

$$H'_{(k-3)2} = H'_{2(k-3)} = \frac{\alpha_0 \alpha_1}{\alpha_{k-1} \alpha_{k-2} \alpha_{k-3}} D_k.$$

In general, an induction argument proves that

$$H'_{(k-j)i} = H'_{i(k-j)} = \frac{\alpha_0 \alpha_1 \cdots \alpha_{i-1}}{\alpha_{k-1} \alpha_{k-2} \cdots \alpha_{k-j}} D_k \text{ for } j-i=1,$$

and the matrix $[e_k \otimes I_m \quad H]$ has the following block anti-triangular form

$$\begin{bmatrix} 0 & 0 & \cdots & 0 & 0 & 0 & \alpha_{k-1}^{-1} D_k \\ 0 & 0 & \cdots & 0 & 0 & \frac{\alpha_0}{\alpha_{k-1} \alpha_{k-2}} D_k & * \\ 0 & 0 & \cdots & 0 & \frac{\alpha_0 \alpha_1}{\alpha_{k-1} \alpha_{k-2} \alpha_{k-3}} D_k & * & * \\ \vdots & & & \ddots & & & \\ 0 & 0 & \frac{\alpha_0 \alpha_1}{\alpha_{k-1} \alpha_{k-2} \alpha_{k-3}} D_k & * & * & * & * \\ 0 & \frac{\alpha_0}{\alpha_{k-1} \alpha_{k-2}} D_k & * & * & * & * & * \\ I_m & * & * & * & * & * & * \end{bmatrix}.$$

Therefore $[e_k \otimes I_m \quad H]$ is nonsingular if and only if D_k is nonsingular. \square

Corollary 7.7. *Let $D(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ be a polynomial matrix with degree $k \geq 2$ and leading matrix coefficient D_k , and let $L(\lambda) = [e_k \otimes I_m \quad H]F_{\Phi}^D(\lambda) \in \mathbb{DM}(D)$. Then the following statements hold:*

1. $L(\lambda)$ is a strong linearization of $D(\lambda)$ if D_k is nonsingular.
2. If $D(\lambda)$ is regular, $L(\lambda)$ is a strong linearization of $D(\lambda)$ if and only if D_k is nonsingular.

Proof. Item 1. follows from Lemma 7.6 and Corollary 3.5. Item 2. follows from Lemma 7.6 and [15, Theorem 3]. \square

Computing the pencil in $\mathbb{DM}(P)$ with ansatz vector e_k , or with any other ansatz vector v , may be difficult. In general, one can follow the procedure in [15, Section 7] or use the MATLAB code in [28, Subsection 7.1]. However, if the recurrence relation (10) is simple and k is low, then the computation can be performed easily by hand, as we illustrate in Example 7.8.

Example 7.8. For a second degree polynomial matrix $D(\lambda) = D_2\phi_2(\lambda) + D_1\phi_1(\lambda) + D_0\phi_0(\lambda)$ expressed in terms of a polynomial basis satisfying (10), the pencil $L(\lambda) \in \mathbb{DM}(D)$ with ansatz vector e_2 is

$$L(\lambda) = \begin{bmatrix} -\frac{\alpha_0}{\alpha_1}D_2 & \frac{\lambda-\beta_0}{\alpha_1}D_2 \\ \frac{\lambda-\beta_0}{\alpha_1}D_2 & \left(\frac{\beta_0-\beta_1}{\alpha_0\alpha_1}(\lambda-\beta_0) - \frac{\gamma_1}{\alpha_1}\right)D_2 + \frac{\lambda-\beta_0}{\alpha_0}D_1 + D_0 \end{bmatrix}.$$

This can be obtained, for instance, by computing the matrix H' as in the proof of Lemma 7.6. For example, Chebyshev polynomials of the first kind $\{\phi_j(\lambda)\}_{j=0}^{\infty}$ satisfy the following three-term recurrence relation:

$$\frac{1}{2}\phi_{j+1}(\lambda) = \lambda\phi_j(\lambda) - \frac{1}{2}\phi_{j-1}(\lambda) \quad j \geq 1 \quad (31)$$

where $\phi_{-1}(\lambda) = 0$, $\phi_0(\lambda) = 1$ and $\phi_1(\lambda) = \lambda$. Therefore, $\alpha_0 = 1$, $\alpha_j = \gamma_j = \frac{1}{2}$ for $j \geq 1$, $\beta_j = 0$ $j \geq 0$ and

$$L(\lambda) = \begin{bmatrix} -2D_2 & 2\lambda D_2 \\ 2\lambda D_2 & \lambda D_1 + D_0 - D_2 \end{bmatrix}.$$

Chebyshev polynomials of the second kind satisfy the same recurrence relation with $\phi_1(\lambda) = 2\lambda$. Thus, $\alpha_j = \gamma_j = \frac{1}{2}$, $\beta_j = 0$ for $j \geq 0$ and

$$L(\lambda) = \begin{bmatrix} -D_2 & 2\lambda D_2 \\ 2\lambda D_2 & 2\lambda D_1 + D_0 - D_2 \end{bmatrix}.$$

For a cubic polynomial matrix $D(\lambda) = D_3\phi_3(\lambda) + D_2\phi_2(\lambda) + D_1\phi_1(\lambda) + D_0\phi_0(\lambda)$ expressed in terms of Chebyshev polynomials of the first kind, the pencil $L(\lambda) \in \mathbb{DM}(D)$ with ansatz vector e_3 is

$$L(\lambda) = \begin{bmatrix} 0 & -2D_3 & 2\lambda D_3 \\ -2D_3 & 4\lambda D_3 - 2D_2 & 2\lambda D_2 - 2D_3 \\ 2\lambda D_3 & 2\lambda D_2 - 2D_3 & \lambda(D_1 + D_3) + D_0 - D_2 \end{bmatrix}.$$

If $D(\lambda)$ is expressed in terms of Chebyshev polynomials of the second kind we obtain

$$L(\lambda) = \begin{bmatrix} 0 & -D_3 & 2\lambda D_3 \\ -D_3 & 2\lambda D_3 - D_2 & 2\lambda D_2 - D_3 \\ 2\lambda D_3 & 2\lambda D_2 - D_3 & 2\lambda D_1 + D_0 - D_2 \end{bmatrix}.$$

By using Theorems 3.9 or 4.4, Theorem 7.5, Lemma 7.6 and Proposition 6.5, we obtain in Theorem 7.9 symmetric strong linearizations of a symmetric rational matrix when the leading coefficient of its polynomial part is nonsingular as we announced.

Theorem 7.9. *Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a symmetric rational matrix and let $G(\lambda) = D(\lambda) + G_{sp}(\lambda)$ be its unique decomposition into its polynomial part $D(\lambda) \in \mathbb{F}[\lambda]^{m \times m}$ and its strictly proper part $G_{sp}(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$. Assume that $\deg(D(\lambda)) = k \geq 2$ and let $n = \nu(G(\lambda))$. Consider a symmetric minimal state-space realization of $G_{sp}(\lambda)$, i.e., $G_{sp}(\lambda) = W(\lambda S_1 - S_2)^{-1}W^T$ as in Definition 6.2, and $L(\lambda) \in \mathbb{DM}(D)$ with ansatz vector e_k . If the leading matrix coefficient D_k of $D(\lambda)$ is nonsingular then, for any nonsingular matrix $Z \in \mathbb{F}^{n \times n}$, the linear polynomial matrix*

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} Z(S_2 - \lambda S_1)Z^T & \begin{array}{c} 0_{n \times (k-1)m} \\ ZW^T \end{array} \\ \hline \begin{array}{c} 0_{(k-1)m \times n} \\ WZ^T \end{array} & L(\lambda) \end{array} \right] \quad (32)$$

is a symmetric strong linearization of $G(\lambda)$.

Proof. Let $L(\lambda) = [e_k \otimes I_m \quad H]F_{\Phi}^D(\lambda)$ be the pencil in $\mathbb{DM}(D)$ with ansatz vector e_k . Since D_k is nonsingular, the matrix $[e_k \otimes I_m \quad H]$ is also nonsingular by using Lemma 7.6. Notice that if $G_{sp}(\lambda) = W(\lambda S_1 - S_2)^{-1}W^T$ is a symmetric minimal state-space realization of $G_{sp}(\lambda)$ then $G_{sp}(\lambda) = W(\lambda I_n - S_1^{-1}S_2)^{-1}S_1^{-1}W^T$ is a minimal state-space realization. It only remains to consider Theorem 3.9 with $X = ZS_1$ and $Y = -Z^T$. Equivalently, we can consider Theorem 4.4 with $X = ZS_1$ and $Y = -Z^T$. \square

Example 7.10. Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a symmetric rational matrix and write $G(\lambda) = D(\lambda) + G_{sp}(\lambda)$ as the sum of its polynomial part and its strictly proper part. Suppose that

$$D(\lambda) = D_k \lambda^k + D_{k-1} \lambda^{k-1} + \dots + D_1 \lambda + D_0,$$

with $k \geq 2$ and D_k nonsingular, and write $G_{sp}(\lambda) = W(\lambda S_1 - S_2)^{-1}W^T$ as a symmetric minimal state-space realization. For the monomial basis we obtain by [20, Theorem 3.5] that the pencil $L(\lambda) \in \mathbb{DL}(D)$ with ansatz vector e_k is

$$L(\lambda) = \lambda \begin{bmatrix} & & & & D_k \\ & & & \ddots & D_{k-1} \\ & & & \ddots & \vdots \\ & & \ddots & \ddots & D_2 \\ D_k & D_{k-1} & \cdots & D_2 & D_1 \end{bmatrix} - \begin{bmatrix} & & & & D_k \\ & & & \ddots & D_{k-1} \\ & & \ddots & \ddots & \vdots \\ D_k & D_{k-1} & \cdots & D_2 & -D_0 \end{bmatrix}.$$

Then, by Theorem 7.9 with $Z = I_n$, the linear polynomial matrix

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|cc} S_2 - \lambda S_1 & 0_{n \times (k-1)m} & W^T \\ \hline 0_{(k-1)m \times n} & & L(\lambda) \\ W & & \end{array} \right]$$

is a symmetric strong linearization of $G(\lambda)$.

We can obtain infinitely many symmetric strong linearizations by using Theorem 7.9 and Lemma 2.7.

Corollary 7.11. *Under the same assumptions as in Theorem 7.9, consider the symmetric strong linearization $\mathcal{L}(\lambda)$ in (32). Let $Q \in \mathbb{F}^{n \times n}$, $P \in \mathbb{F}^{km \times km}$ be nonsingular matrices and $R \in \mathbb{F}^{km \times n}$. Then*

$$\widehat{\mathcal{L}}(\lambda) = \begin{bmatrix} Q & 0 \\ R & P \end{bmatrix} \mathcal{L}(\lambda) \begin{bmatrix} Q^T & R^T \\ 0 & P^T \end{bmatrix}$$

is a symmetric strong linearization of $G(\lambda)$.

8. Hermitian strong linearizations for Hermitian rational matrices

In this section we extend the results in Sections 6 and 7 from symmetric to Hermitian rational matrices. Since most of the arguments are similar to those in the symmetric case, we limit ourselves to state the main results, and most of the proofs are omitted. We consider the ring of polynomials $\mathbb{C}[\lambda]$ and a polynomial basis $\{\phi_j(\lambda)\}_{j=0}^{\infty}$ that satisfies the three-term recurrence relation:

$$\alpha_j \phi_{j+1}(\lambda) = (\lambda - \beta_j) \phi_j(\lambda) - \gamma_j \phi_{j-1}(\lambda) \quad j \geq 0$$

as in (10), with $\alpha_j, \beta_j, \gamma_j \in \mathbb{R}$, $\alpha_j \neq 0$, $\phi_{-1}(\lambda) = 0$, and $\phi_0(\lambda) = 1$. Let $P(\lambda) \in \mathbb{C}[\lambda]^{m \times m}$ be a polynomial matrix of degree k written in terms of this basis, i.e., $P(\lambda) = \sum_{i=0}^k P_i \phi_i(\lambda)$ with $P_i \in \mathbb{C}^{m \times m}$. Suppose that $P(\lambda)$ is *Hermitian*, i.e., $P(\lambda)^* = P(\bar{\lambda})$ or, equivalently, $P^*(\lambda) = P(\lambda)$, where $P^*(\lambda)$ is defined as $P^*(\lambda) = \sum_{i=0}^k P_i^* \phi_i(\lambda)$ with P_i^* the conjugate transpose of $P_i \in \mathbb{C}^{m \times m}$. We also consider the set of pencils

$$\mathbb{H}(P) = \{\lambda X + Y \in \mathbb{M}_1(P) : X^* = X, Y^* = Y\}.$$

That is, $\mathbb{H}(P)$ is the set of pencils in $\mathbb{M}_1(P)$ that are Hermitian. Theorem 8.1 shows that the elements of $\mathbb{H}(P)$ are in $\mathbb{DM}(P)$, and that, in fact, they are the pencils in $\mathbb{DM}(P)$ with real ansatz vector. The proof of Theorem 8.1 is omitted for brevity since it is similar to the proof of [20, Theorem 6.1], which is Theorem 8.1 in the particular case $\phi_j(\lambda) = \lambda^j$ for $j \geq 0$.

Theorem 8.1. *Let $P(\lambda) \in \mathbb{C}[\lambda]^{m \times m}$ be a Hermitian polynomial matrix. Then $\mathbb{H}(P)$ is the subset of all pencils in $\mathbb{DM}(P)$ with real ansatz vector.*

Let $G(\lambda) \in \mathbb{C}(\lambda)^{m \times m}$ be a Hermitian rational matrix, i.e., a rational matrix satisfying $G(\lambda)^* = G(\bar{\lambda})$. Consider $G(\lambda) = D(\lambda) + G_{sp}(\lambda)$ as in (3). Then $D(\lambda)$ and $G_{sp}(\lambda)$ are also Hermitian. For Hermitian strictly proper rational matrices we introduce the notion of Hermitian minimal state-space realizations, in the spirit of Definition 6.2.

Definition 8.2. Let $G_{sp}(\lambda) \in \mathbb{C}(\lambda)^{m \times m}$ be a Hermitian strictly proper rational matrix and let $n = \nu(G_{sp}(\lambda))$. A Hermitian minimal state-space realization of $G_{sp}(\lambda)$ is an expression of the form

$$G_{sp}(\lambda) = W(\lambda H_1 - H_2)^{-1} W^*$$

where $H_1, H_2 \in \mathbb{C}^{n \times n}$ are Hermitian matrices, with H_1 nonsingular, and $W \in \mathbb{C}^{m \times n}$.

Following arguments similar to those in Lemma 6.3 and Proposition 6.5, it is easy to see that the strictly proper part of a Hermitian rational matrix has a Hermitian minimal state-space realization.

Proposition 8.3. Any Hermitian strictly proper rational matrix has a Hermitian minimal state-space realization.

Proof. In order to obtain a Hermitian minimal state-space realization of $G_{sp}(\lambda)$, we can consider a minimal state-space realization $G_{sp}(\lambda) = C(\lambda I_n - A)^{-1} B$. We prove analogously to Lemma 6.3 that there exists a unique nonsingular and Hermitian matrix $H \in \mathbb{C}^{n \times n}$ such that $A^* = H^{-1} A H$ and $C^* = H^{-1} B$. Therefore, $G_{sp}(\lambda) = C(\lambda H^{-1} - H^{-1} A)^{-1} C^*$ is a Hermitian minimal state-space realization of $G_{sp}(\lambda)$. \square

Remark 8.4. Another constructive way to prove Proposition 8.3 is to consider the Hankel matrix H_n of $G_{sp}(\lambda)$ defined in (24), that is also Hermitian, and write

$$H_n = U \begin{bmatrix} K & 0 \\ 0 & 0 \end{bmatrix} U^*$$

with U unitary, i.e., $U^{-1} = U^*$, and K a diagonal matrix that has the eigenvalues different from zero of H_n at the diagonal elements. Then proceed as in the last paragraph of Remark 6.6 to get a Hermitian minimal state-space realization. Notice that K is Hermitian because the eigenvalues of H_n are real. \square

By using Proposition 8.3 and Theorem 8.1, we obtain in Theorem 8.5 Hermitian strong linearizations of a Hermitian rational matrix when the leading coefficient of its polynomial part is nonsingular, analogously as we did in Theorem 7.9 for the symmetric case.

Theorem 8.5. Let $G(\lambda) \in \mathbb{C}(\lambda)^{m \times m}$ be a Hermitian rational matrix and let $G(\lambda) = D(\lambda) + G_{sp}(\lambda)$ be its unique decomposition into its polynomial part $D(\lambda) \in \mathbb{C}[\lambda]^{m \times m}$ and its strictly proper part $G_{sp}(\lambda) \in \mathbb{C}(\lambda)^{m \times m}$. Assume that $\deg(D(\lambda)) = k \geq 2$ and let $n = \nu(G(\lambda))$. Consider a Hermitian minimal state-space realization of $G_{sp}(\lambda)$, i.e., $G_{sp}(\lambda) = W(\lambda H_1 - H_2)^{-1} W^*$ as in Definition 8.2, and $L(\lambda) \in \mathbb{DM}(D)$ with ansatz vector e_k . If the leading matrix coefficient D_k of $D(\lambda)$ is nonsingular then, for any nonsingular matrix $Z \in \mathbb{C}^{n \times n}$, the linear polynomial matrix

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} Z(H_2 - \lambda H_1)Z^* & \begin{array}{c} 0_{n \times (k-1)m} \\ ZW^* \end{array} \\ \hline \begin{array}{c} 0_{(k-1)m \times n} \\ WZ^* \end{array} & L(\lambda) \end{array} \right] \quad (33)$$

is a Hermitian strong linearization of $G(\lambda)$.

As in Corollary 7.11, we can obtain infinitely many Hermitian strong linearizations by using Theorem 8.5 and Lemma 2.7.

Corollary 8.6. *Under the same assumptions as in Theorem 8.5, consider the Hermitian strong linearization $\mathcal{L}(\lambda)$ in (33). Let $Q \in \mathbb{C}^{n \times n}$, $P \in \mathbb{C}^{km \times km}$ be nonsingular matrices and $R \in \mathbb{C}^{km \times n}$. Then*

$$\widehat{\mathcal{L}}(\lambda) = \begin{bmatrix} Q & 0 \\ R & P \end{bmatrix} \mathcal{L}(\lambda) \begin{bmatrix} Q^* & R^* \\ 0 & P^* \end{bmatrix}$$

is a Hermitian strong linearization of $G(\lambda)$.

9. Strong linearizations of rational matrices with polynomial part expressed in other polynomial bases

Polynomial bases $\{\phi_j(\lambda)\}_{j=0}^{\infty}$ satisfying a three-term recurrence relation as in (10) are by far the most useful in applications. However, from a theoretical point of view, a natural question is whether or not the results in this paper can be extended to other polynomial bases. The goal of this section is to show that this can be done by using exactly the same tools that we have used in previous sections, that is, [4, Theorem 5.11], our key Lemma 2.7, and the results in [15]. Since the arguments in this section are very similar to the ones previously used, we will simply sketch the main ideas.

Let $D(\lambda)$ be the polynomial part of a rational matrix $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$, with $\deg(D(\lambda)) = k \geq 2$. Let us consider, motivated by (12) and its properties, a polynomial basis $\{\psi_j(\lambda)\}_{j=0}^{\infty}$ of $\mathbb{F}[\lambda]$, with $\psi_j(\lambda)$ a polynomial of degree j , that satisfies a linear relation:

$$M_{\Psi}(\lambda)\Psi_k(\lambda) = 0, \quad (34)$$

where $M_{\Psi}(\lambda) \in \mathbb{F}[\lambda]^{(k-1) \times k}$ is a minimal basis with all its row degrees equal to 1, and $\Psi_k(\lambda) = [\psi_{k-1}(\lambda) \cdots \psi_1(\lambda) \psi_0(\lambda)]^T$ with $\Psi_k(\lambda_0) \neq 0$ for all $\lambda_0 \in \overline{\mathbb{F}}$. Then there exists a vector $w \in \mathbb{F}^k$ such that

$$U(\lambda) = \begin{bmatrix} M_{\Psi}(\lambda) \otimes I_m \\ w^T \otimes I_m \end{bmatrix} \quad (35)$$

is unimodular, and its inverse has the form $U(\lambda)^{-1} = [\widehat{\Psi}_k(\lambda) \quad \Psi_k(\lambda) \otimes I_m]$ with $\widehat{\Psi}_k(\lambda) \in \mathbb{F}[\lambda]^{km \times (k-1)m}$ (see [4, Lemma 5.5]). Let

$$F_{\Psi}^D(\lambda) = \begin{bmatrix} m_{\Psi}^D(\lambda) \\ M_{\Psi}(\lambda) \otimes I_m \end{bmatrix} \in \mathbb{F}[\lambda]^{km \times km} \quad (36)$$

be a pencil such that $m_{\Psi}^D(\lambda)(\Psi_k(\lambda) \otimes I_m) = D(\lambda)$. Then, $F_{\Psi}^D(\lambda)$ is a strong block minimal bases pencil associated to $D(\lambda)$ with sharp degree which verifies $F_{\Psi}^D(\lambda)(\Psi_k(\lambda) \otimes I_m) = e_1 \otimes D(\lambda)$. Thus, we can apply [4, Theorem 5.11] and Lemma 2.7 in order to construct strong linearizations of $G(\lambda)$ from pencils of the form $L(\lambda) = [v \otimes I_m \quad H]F_{\Psi}^D(\lambda)$ with $v \in \mathbb{F}^k$ and $[v \otimes I_m \quad H]$ nonsingular. Notice that pencils $L(\lambda)$ of this form verify the ansatz relation $L(\lambda)(\Psi_k(\lambda) \otimes I_m) = v \otimes D(\lambda)$. In summary, with those arguments, we obtain the following result that is the generalization of Theorem 3.9 for polynomial bases as in (34).

Theorem 9.1. Let $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ be a rational matrix written as in (3), and let $G_{sp}(\lambda) = C(\lambda I_n - A)^{-1}B$ be a minimal state-space realization of $G_{sp}(\lambda)$. Assume that $\deg(D(\lambda)) \geq 2$ and write $D(\lambda)$ in terms of a polynomial basis $\{\psi_j(\lambda)\}_{j=0}^{\infty}$ satisfying (34), as

$$D(\lambda) = D_k \psi_k(\lambda) + D_{k-1} \psi_{k-1}(\lambda) + \cdots + D_1 \psi_1(\lambda) + D_0 \psi_0(\lambda) \quad (37)$$

with $D_k \neq 0$. Let $L(\lambda) = [v \otimes I_m \quad H] F_{\Psi}^D(\lambda)$ with $[v \otimes I_m \quad H]$ nonsingular and $F_{\Psi}^D(\lambda)$ as in (36). Let $w \in \mathbb{F}^k$ be the vector in (35). Then, for any nonsingular matrices $X, Y \in \mathbb{F}^{n \times n}$ the linear polynomial matrix

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & XB(w^T \otimes I_m) \\ \hline -(v \otimes I_m)CY & L(\lambda) \end{array} \right]$$

is a strong linearization of $G(\lambda)$.

In a similar manner, the results in Sections 4 and 5 can be extended to square rational matrices with polynomial parts expressed in terms of polynomial bases as in (34).

In Example 9.2 we consider degree-graded polynomial bases presented in [15, Section 9], and we construct strong linearizations of square rational matrices by expressing the polynomial parts in terms of these bases and using Theorem 9.1.

Example 9.2. Let $\{\psi_j(\lambda)\}_{j=0}^{\infty}$ be a degree-graded polynomial basis of $\mathbb{F}[\lambda]$ that satisfies the following recurrence relation:

$$\psi_j(\lambda) = (\lambda - \alpha_j) \psi_{j-1}(\lambda) + \sum_{i=0}^{j-2} \beta_j^i \psi_i(\lambda) \quad j \geq 1$$

where $\alpha_j \in \mathbb{F}$ for $j \geq 1$, $\beta_j^i \in \mathbb{F}$ for $j \geq 2$, $0 \leq i \leq j-2$ and $\psi_0(\lambda) = 1$. Let $G(\lambda) = D(\lambda) + C(\lambda I_n - A)^{-1}B$ be an $m \times m$ rational matrix written as in Theorem 9.1. We express the polynomial part $D(\lambda)$ in terms of the polynomial basis $\{\psi_j(\lambda)\}_{j=0}^{\infty}$, as in (37). Let us denote $\Psi_k(\lambda) = [\psi_{k-1}(\lambda) \cdots \psi_1(\lambda) \psi_0(\lambda)]^T$ and consider the following pencil $G_{\Psi}^D(\lambda)$ introduced in [15, Section 9]:

$$G_{\Psi}^D(\lambda) = \left[\begin{array}{c} m_{\Psi}^D(\lambda) \\ M_{\Psi}(\lambda) \otimes I_m \end{array} \right] \in \mathbb{F}[\lambda]^{km \times km},$$

where

$$m_{\Psi}^D(\lambda) = \left[(\lambda - \alpha_k)D_k + D_{k-1} \quad \beta_k^{k-2}D_k + D_{k-2} \quad \cdots \quad \beta_k^1 D_k + D_1 \quad \beta_k^0 D_k + D_0 \right],$$

and

$$M_{\Psi}(\lambda) = \left[\begin{array}{cccccccc} -1 & (\lambda - \alpha_{k-1}) & \beta_{k-1}^{k-3} & \beta_{k-1}^{k-4} & \cdots & \beta_{k-1}^2 & \beta_{k-1}^1 & \beta_{k-1}^0 \\ & -1 & (\lambda - \alpha_{k-2}) & \beta_{k-2}^{k-4} & \cdots & \beta_{k-2}^2 & \beta_{k-2}^1 & \beta_{k-2}^0 \\ & & \ddots & \ddots & \ddots & \vdots & \vdots & \vdots \\ & & & & & -1 & (\lambda - \alpha_2) & \beta_2^0 \\ & & & & & & -1 & (\lambda - \alpha_1) \end{array} \right].$$

The matrix $G_{\Psi}^D(\lambda)$ verifies that $G_{\Psi}^D(\lambda)(\Psi_k(\lambda) \otimes I_m) = e_1 \otimes D(\lambda)$. Moreover, $G_{\Psi}^D(\lambda)$ is a strong block minimal bases pencil associated to $D(\lambda)$ with sharp degree. It can be proved, as in [15, Theorem 1], that any pencil $L(\lambda)$ that verifies $L(\lambda)(\Psi_k(\lambda) \otimes I_m) = v \otimes D(\lambda)$ for some vector $v \in \mathbb{F}^k$ can be written as $L(\lambda) = [v \otimes I_m \quad H]G_{\Psi}^D(\lambda)$ for some matrix $H \in \mathbb{F}^{km \times (k-1)m}$. If we consider a pencil $L(\lambda)$ of this form with $[v \otimes I_m \quad H]$ nonsingular we can obtain strong linearizations for $G(\lambda)$. In particular, we have that conditions in Theorem 9.1 hold, and we can apply it with $w = e_k$. Then, we have that for any nonsingular matrices $X, Y \in \mathbb{F}^{n \times n}$ the linear polynomial matrix

$$\mathcal{L}(\lambda) = \left[\begin{array}{c|c} X(\lambda I_n - A)Y & 0_{n \times (k-1)m} \quad XB \\ \hline -(v \otimes I_m)CY & L(\lambda) \end{array} \right]$$

is a strong linearization of $G(\lambda)$.

10. Conclusions and future work

As a consequence of the definitions and the theory developed in [5], we have proved the simple Lemma 2.7, which allows us to construct infinitely many strong linearizations of any rational matrix $G(\lambda) \in \mathbb{F}(\lambda)^{p \times m}$ from any given strong linearization of $G(\lambda)$. This result has been combined with some of the strong linearizations of a rational matrix $G(\lambda) \in \mathbb{F}(\lambda)^{m \times m}$ constructed in [4, Theorem 5.11] and with the strong linearizations of its polynomial part presented in [15] to create new families of strong linearizations of square rational matrices. The recovery of the eigenvectors of the rational matrix from those of the linearizations in these families has been thoroughly investigated, as well as the preservation of symmetric and Hermitian structures of the rational matrix in the linearizations.

We are convinced that the techniques developed in this paper together with the results in [5] can be applied to solve essentially all the following problems: How to construct a strong linearization of a rational matrix $G(\lambda)$ expressed as the sum of its polynomial part $D(\lambda)$ and its strictly proper part $G_{sp}(\lambda)$, given a strong linearization of $D(\lambda)$ in any of the families of strong linearizations of polynomial matrices developed in the last years and a minimal state-space realization of $G_{sp}(\lambda)$. In particular, we hope that these techniques will allow to construct strong linearizations of rational matrices preserving structures that are different from the symmetric and Hermitian structures. However, we emphasize that, although any rational matrix can be expressed as the sum of its polynomial and strictly proper parts, this expression may not be easily available from the applications and/or may not be the best representation in a particular problem. Therefore, the development of strong linearizations of rational matrices starting from other representations is a problem that will be investigated in the future.

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