

The matrix Sylvester equation for congruence

Froilán M. Dopico

ICMAT and Departamento de Matemáticas, Universidad Carlos III de Madrid, Spain

joint work with Fernando De Terán, Nathan Guillery, Daniel Montealegre, and
Nicolás Reyes

School of Mathematics, University of Edinburgh, Scotland
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The **matrix Sylvester equation**

$$AX - XB = C, \quad A \in \mathbb{C}^{m \times m}, B \in \mathbb{C}^{n \times n} \text{ are given}$$

is, probably, the most famous matrix equation. It arises

- as a step in algorithms for computing eigenvalues/vectors;
- in the perturbation theory of invariant subspaces of matrices;
- in the characterization of the matrices that commute with a given matrix $AX = XA$.
- Its particular case, the Lyapunov equation,

$$AX + XA^* = C$$

arises in control and linear system theory and in stability theory...

- ▶ Properties of Sylvester eq. are well-known and are presented in standard books on Matrix Analysis.
- ▶ Numerical methods for solution are also well-known.

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Abstract (II)

Recently, the **matrix Sylvester equation for congruence** or T-Sylvester equation

$$AX + X^T B = C, \quad A \in \mathbb{C}^{m \times n}, B \in \mathbb{C}^{n \times m}$$

has received considerable attention as a consequence of its relationship with **palindromic eigenvalue problems**

$$Gx = -\lambda G^T x, \quad G \in \mathbb{C}^{n \times n}.$$

These problems arise in a number of applications:

- the mathematical modelling and numerical simulation of the behavior of periodic surface acoustic wave filters (2002, 2006);
- the analysis of rail track vibrations produced by high speed trains (2004, 2006, 2009);
- discrete-time optimal control problems (2008).

The spectrum of palindromic eigenproblems has the symmetry $(\lambda, 1/\lambda)$.

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$$AX - XB = C \quad \text{vs.} \quad AX + X^T B = C$$

The “transposed second X ” makes the study of both equations very different.
Not many? references available for T-Sylvester equation.

In this talk, I will revise my research work on Sylvester equation for congruence that has been published in

- F. De Terán & D, *The solution of the equation $XA + AX^T = 0$ and its application to the theory of orbits*, Lin. Alg. Appl. 434 (2011) 44-67.
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Both interesting, both related, but DIFFERENT!!!

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In this talk for simplicity mostly **T-case** is considered,

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but sometimes both cases simultaneously: $\star = T$ or $*$

1 Previous and related work

2 The equation $AX^T + XA = 0$

- Motivation: Orbits and the computation of canonical forms
- Strategy for solving $AX^T + XA = 0$
- The canonical form for congruence
- The solution of $AX^T + XA = 0$
- Generic canonical structure for congruence

3 The general equation $AX + X^*B = C$

- Motivation
- Consistency of the Sylvester equation for $\star$ -congruence
- Uniqueness of solutions
- Efficient and stable algorithm to compute unique solutions

4 General solution of $AX + X^*B = 0$

5 Conclusions

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5 Conclusions

An important particular case of $AX + X^*B = C$

$$AX + X^*A^* = C \quad (\star = T \text{ or } *)$$

arises in time-invariant Hamiltonian systems and R-matrix treatment of completely integrable mechanical systems.

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- De Terán & D., Lin. Alg. Appl. and Elec. J. Lin. Alg., (2011):
 - General solution obtained in the spirit of classical methods of solution of standard Sylvester equation.
 - Related to the theory of orbits by the action of congruence.

References for general equation $AX + X^\star B = C$

- Wimmer, Lin. Alg. Appl., (1994): Necessary and sufficient conditions for **consistency**. Complex matrices
- Byers & Kressner, SIAM J. Matrix Anal. Appl., (2006): Necessary and sufficient conditions for **unique solution** for $\star = T$. Complex matrices.
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- Motivation: Orbits and the computation of canonical forms
- Strategy for solving $AX^T + XA = 0$
- The canonical form for congruence
- The solution of $AX^T + XA = 0$
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Due to roundoff errors, uncertainty in the data, ... , usually it **is not possible** to compute the **exact canonical forms of matrix eigenvalue problems**.

- $Ax = \lambda x$ (Jordan Canonical Form (JCF)).
- $Ax = \lambda Bx$ (Kronecker Canonical Form (KCF)).

Some related questions:

- Which are the **nearby** canonical structures (JCF, KCF) to a given one?
- Which is the **generic** canonical structure?

Same questions for matrices/matrix pencils in a particular subset (low-rank, palindromic, symmetric,...) and **structure preserving numerical methods**.

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Congruence, equivalence, and similarity. Orbits

Given $A, B \in \mathbb{C}^{n \times n}$

$$\mathcal{O}(A) = \{PAP^T : P \text{ nonsingular}\}$$

Congruence orbit of A

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Equivalency orbit of $A - \lambda B$

Similarity/equivalency orbits

- have been widely studied: Arnold (1971), Demmel-Edelman (1995), Edelman-Elmroth-Kågström (1997, 1999), Johansson (2006), ...
- correspond to matrices with the same Jordan Canonical Form (JCF) / Pencils with the same Kronecker Canonical Form (KCF).
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Codimension of the tangent space

$$\begin{aligned} T_{\mathcal{O}(A)}(A) &= \{XA + AX^T : X \in \mathbb{C}^{n \times n}\} && \text{Tangent space of } \mathcal{O}(A) \text{ at } A \\ T_{\mathcal{O}_s(A)}(A) &= \{XA - AX : X \in \mathbb{C}^{n \times n}\} && \text{Tangent space of } \mathcal{O}_s(A) \text{ at } A \end{aligned}$$

Then:

$$(a) \operatorname{codim} \mathcal{O}(A) = \operatorname{codim} T_{\mathcal{O}(A)}(A) = \dim(\text{solution space of } XA + AX^T = 0)$$

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General solution of $XA - AX = 0$: known since the 1950's ([Gantmacher](#)) and probably before. Depends on the JCF of A .

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Getting a simpler equation via congruence

Notation: $\mathcal{S}_A = \{X \in \mathbb{C}^{n \times n} : AX^T + XA = 0\}$ (solution space)

Consider $B := PAP^T$ (P nonsingular) then

$$B(PXP^{-1})^T + (PXP^{-1})B = 0$$

and $\mathcal{S}_A = P^{-1}\mathcal{S}_BP$.

In particular: $\dim \mathcal{S}_A = \dim \mathcal{S}_B$

Procedure to solve $AX^T + XA = 0$:

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Theorem (Canonical form for congruence (Horn & Sergeichuk, 2006))

Each matrix $A \in \mathbb{C}^{n \times n}$ is **congruent** to a direct sum, uniquely determined up to permutation of summands, of blocks of types 0, I and II.

$$(Type\ 0) \quad J_k(0) = \begin{bmatrix} 0 & 1 & & \\ & \ddots & \ddots & \\ & & 0 & 1 \\ & & & 0 \end{bmatrix}_{k \times k}$$

$$(Type\ I) \quad \Gamma_k = \begin{bmatrix} 0 & & & & & & (-1)^{k+1} \\ & & & & & \ddots & (-1)^k \\ & & & & \ddots & & \\ & & & -1 & \ddots & & \\ & & & & 1 & 1 & \\ & & -1 & -1 & & & \\ 1 & 1 & & & & & 0 \end{bmatrix}_{k \times k}, \quad \Gamma_1 = [1]$$

$$(Type\ II) \quad H_{2k}(\mu) = \begin{bmatrix} 0 & I_k \\ J_k(\mu) & 0 \end{bmatrix}_{2k \times 2k}, \quad H_2(\mu) = \begin{bmatrix} 0 & 1 \\ \mu & 0 \end{bmatrix}, \quad (0 \neq \mu \neq (-1)^{k+1})$$

The canonical form for congruence: a brief history (I)

- **Turnbull** (U. St. Andrews, Scotland) & **Aitken** (U. Edinburgh, Scotland), *An Introduction to the Theory of Canonical Matrices*, 1932.
For complex matrices. **Six types of blocks**.
- **Gabriel**, J. Algebra (1974), studied equivalence of bilinear forms in fields with characteristic $\neq 2$.
- **Riehm**, J. Algebra (1974), reduced the problem of equivalence of bilinear forms to equivalence of Hermitian forms.
- **Sergeichuk**, Math. USSR Izvestiya (1988) complete study via quivers and Hermitian forms in fields with characteristic $\neq 2$.
- **Thompson**, Linear Algebra and its Applications (1991). Complex and real matrices: Symmetric/Skew-Symmetric pencils.
- **Lee and Weinberg**, Linear Algebra and its Applications (1996). Complex and real matrices based on Thompson and $A = S + K$, with $S = S^T$ and $K = -K^T$. Six blocks for complex (Turnbull and Aitken). Eight blocks for real matrices.

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Proofs: first based on quivers and second constructive and based only on basic Matrix Analysis.

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Procedure to solve $AX^T + XA = 0$:

- 1 Set $C_A = PAP^T$, the **canonical form of A for congruence**.
- 2 Solve $C_A Y^T + Y C_A = 0$.
- 3 Undo the change: $X = P^{-1} Y P$.

Partition into blocks to solve $XC_A + C_A X^T = 0$ (I)

Set $C_A = D_1 \oplus \cdots \oplus D_s$, $D_i = J_k(0)$, Γ_k , or $H_{2k}(\mu)$ (Canonical form of A)

$$\text{Partition } X = \begin{bmatrix} X_{11} & \cdots & X_{1s} \\ \vdots & & \vdots \\ X_{s1} & \cdots & X_{ss} \end{bmatrix} \text{ conformally with } C_A.$$

Equating the (i, j) and (j, i) blocks of $XC_A + C_A X^T = 0$, we get:

- $i = j$: $X_{ii}D_i + D_iX_{ii}^T = 0 \rightarrow \text{codim } D_i$ (codimension)
- $i \neq j$: $\begin{matrix} (i, j) & X_{ij}D_j + D_iX_{ji}^T = 0 \\ (j, i) & X_{ji}D_i + D_jX_{ij}^T = 0 \end{matrix} \rightarrow \text{inter}(D_i, D_j)$ (interaction)

Then:

$$\dim \mathcal{S}_A = \text{codim } \mathcal{O}(A) = \sum_i \text{codim } D_i + \sum_{i \neq j} \text{inter}(D_i, D_j)$$

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Partition into blocks to solve $XC_A + C_AX^T = 0$ (II)

The problem reduces to solve matrix equations of the types:

(a) $XD + DX^T = 0$ (easier Sylvester equation for congruence)

with $D = J_k(0)$ (type 0), Γ_k (type I), or $H_{2k}(\mu)$ (type II)
(3 different types of eqs.)

(b) $\begin{aligned} XD_1 + D_2Y^T &= 0 \\ YD_2 + D_1X^T &= 0 \end{aligned}$ (system of two matrix equations)

with $D_1, D_2 = J_k(0)$ (type 0), Γ_ℓ (type I), or $H_{2m}(\mu)$ (type II)
(6 different types of eqs.)

Theorem (De Terán & D, Lin. Alg. Appl., 2011)

Let $A \in \mathbb{C}^{n \times n}$ be a matrix with canonical form for congruence

$$\begin{aligned} C_A = & J_{p_1}(0) \oplus J_{p_2}(0) \oplus \cdots \oplus J_{p_a}(0) \\ & \oplus \Gamma_{q_1} \oplus \Gamma_{q_2} \oplus \cdots \oplus \Gamma_{q_b} \\ & \oplus H_{2r_1}(\mu_1) \oplus H_{2r_2}(\mu_2) \oplus \cdots \oplus H_{2r_c}(\mu_c). \end{aligned}$$

Then the *codimension of the orbit of A* for the action of congruence, i.e., the dimension of the solution space of $XA + AX^T = 0$, depends only on C_A . It can be computed as the sum

$$c_{\text{Total}} = c_0 + c_1 + c_2 + i_{00} + i_{11} + i_{22} + i_{01} + i_{02} + i_{12}.$$

Codimensions and interactions of canonical blocks

Codimension

$$c_0 \rightarrow \left\lceil \frac{k}{2} \right\rceil \quad c_1 \rightarrow \left\lfloor \frac{k}{2} \right\rfloor \quad c_2 \rightarrow \begin{cases} k, & \text{if } \mu \neq (-1)^k \\ k + 2 \left\lceil \frac{k}{2} \right\rceil, & \text{if } \mu = (-1)^k \end{cases}$$

Interaction (same type)

$$i_{00} \rightarrow \begin{cases} \ell, & \ell \text{ even} \\ k, & \ell \text{ odd and } k \neq \ell \\ k + 1, & \ell \text{ odd and } k = \ell \end{cases}$$

$$i_{11} \rightarrow \begin{cases} 0, & k, \ell \text{ different parity} \\ \min\{k, \ell\}, & k, \ell \text{ same parity} \end{cases}$$

$$i_{22} \rightarrow \begin{cases} 4 \min\{k, \ell\}, & \mu = \tilde{\mu} = \pm 1 \\ 2 \min\{k, \ell\}, & \mu = \tilde{\mu} \neq \pm 1 \\ 2 \min\{k, \ell\}, & \mu \neq \tilde{\mu}, \mu \tilde{\mu} = 1 \\ 0, & \mu \neq \tilde{\mu}, \mu \tilde{\mu} \neq 1 \end{cases}$$

Interaction (different type)

$$i_{01} \rightarrow \begin{cases} 0, & k \text{ even} \\ \ell, & k \text{ odd} \end{cases}$$

$$i_{02} \rightarrow \begin{cases} 0, & k \text{ even} \\ 2\ell, & k \text{ odd} \end{cases}$$

$$i_{12} \rightarrow \begin{cases} 2 \min\{k, \ell\}, & \mu = (-1)^{k+1} \\ 0, & \mu \neq (-1)^{k+1} \end{cases}$$

► Explicit solution found by De Terán & D (LAA, 2011) in all cases, except for the case eq. corresp. to codim. of two special type II blocks:

$$X H_{2k}((-1)^k) + H_{2k}((-1)^k) X^T = 0.$$

This solved by S. R. García & A. L. Shoemaker, Lin. Alg. Appl. 2012.

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Generic = codimension zero

Theorem (De Terán & D, Lin. Alg. Appl., 2011)

The minimal codimension for a congruence orbit in $\mathbb{C}^{n \times n}$ is $\lfloor n/2 \rfloor$.

Generic canonical **structure** for congruence is not given by a single orbit!!

Similarity orbits (JCF): There is no generic JCF with fixed eigenvalues.

► The **generic Jordan structure** is $J_1(\lambda_1) \oplus \cdots \oplus J_1(\lambda_n)$, with $\lambda_1, \dots, \lambda_n$ different (not fixed)

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Definition (Arnold, 1971)

Given $A \in \mathbb{C}^{n \times n}$ with Jordan Canonical Form

$$J_A = J_{\lambda_1} \oplus \cdots \oplus J_{\lambda_d},$$

where

$$J_{\lambda_i} := J_{n_{i,1}}(\lambda_i) \oplus \cdots \oplus J_{n_{i,q_i}}(\lambda_i), \quad \text{for } i = 1, \dots, d \text{ and } \lambda_i \neq \lambda_j \text{ if } i \neq j,$$

the **similarity bundle** of A is

$$\mathcal{B}_s(A) = \bigcup_{\substack{\lambda'_i \in \mathbb{C}, i=1, \dots, d \\ \lambda'_i \neq \lambda'_j, i \neq j}} \mathcal{O}_s \left(J_{\lambda'_1} \oplus \cdots \oplus J_{\lambda'_d} \right)$$

Definition (De Terán & D, Lin. Alg. Appl., 2011)

Given $A \in \mathbb{C}^{n \times n}$ with canonical form for congruence

$$C_A = \bigoplus_{i=1}^a J_{p_i}(0) \oplus \bigoplus_{i=1}^b \Gamma_{q_i} \oplus \bigoplus_{i=1}^t \mathcal{H}(\mu_i), \quad \mu_i \neq \mu_j, \mu_i \neq 1/\mu_j \text{ if } i \neq j,$$

where

$$\mathcal{H}(\mu_i) = H_{2r_{i,1}}(\mu_i) \oplus H_{2r_{i,2}}(\mu_i) \oplus \cdots \oplus H_{2r_{i,g_i}}(\mu_i), \quad \text{for } i = 1, \dots, t,$$

the **congruence bundle** of A is

$$\mathcal{B}(A) = \bigcup_{\substack{\mu'_i \in \mathbb{C}, i=1, \dots, t \\ \mu'_i \neq \mu'_j, \mu'_i \mu'_j \neq 1, i \neq j}} \mathcal{O} \left(\bigoplus_{i=1}^a J_{p_i}(0) \oplus \bigoplus_{i=1}^b \Gamma_{q_i} \oplus \bigoplus_{i=1}^t \mathcal{H}(\mu'_i) \right).$$

(same structure as C_A but unfixed complex values μ in type II blocks)

The generic canonical structure for congruence

If t =number of different μ' s appearing in type II blocks of C_A , then $\text{codim}(\mathcal{B}(A)) = \text{codim}(\mathcal{O}(A)) - t$.

Theorem (De Terán & D, Lin. Alg. Appl., 2011)

The following bundles for congruence in $\mathbb{C}^{n \times n}$ have **codimension zero**

1 n even

$$G_n = \mathcal{B} \left(H_2(\mu_1) \oplus H_2(\mu_2) \oplus \cdots \oplus H_2(\mu_{n/2}) \right),$$

with $\mu_i \neq \pm 1, i = 1, \dots, n/2, \mu_i \neq \mu_j$ and $\mu_i \neq 1/\mu_j$ if $i \neq j$.

2 n odd

$$G_n = \mathcal{B} \left(H_2(\mu_1) \oplus H_2(\mu_2) \oplus \cdots \oplus H_2(\mu_{(n-1)/2}) \oplus \Gamma_1 \right),$$

with $\mu_i \neq \pm 1, i = 1, \dots, (n-1)/2, \mu_i \neq \mu_j$ and $\mu_i \neq 1/\mu_j$ if $i \neq j$.

Then G_n is the **generic canonical structure for congruence** in $\mathbb{C}^{n \times n}$ (with unspecified values $\mu_1, \mu_2, \dots$).

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Summary of section 3

Given $A \in \mathbb{C}^{m \times n}$, $B \in \mathbb{C}^{n \times m}$, and $C \in \mathbb{C}^{m \times m}$, we study the equations

$$AX + X^*B = C, \quad (X^* = X^T \text{ or } X^*),$$

where $X \in \mathbb{C}^{n \times m}$ is the unknown to be determined. More precisely:

- 1 Necessary and sufficient conditions for consistency (Wimmer 1994, De Terán & D., Elect. J. Lin. Alg., 2011 (2)).
- 2 Necessary and sufficient conditions for uniqueness of solutions (Byers, Kressner, Schröder, Watkins, 2006, 2009).
- 3 Efficient and stable numerical algorithm for computing the unique solution (De Terán & D., Elect. J. Lin. Alg., 2011 (2)).

We establish parallelisms/differences with well-known Sylvester equation

$$AX - XB = C, \quad A \in \mathbb{C}^{m \times m}, B \in \mathbb{C}^{n \times n}, C \in \mathbb{C}^{m \times n}.$$

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Summary of section 3

Given $A \in \mathbb{C}^{m \times n}$, $B \in \mathbb{C}^{n \times m}$, and $C \in \mathbb{C}^{m \times m}$, we study the equations

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$$A X - X B = C, \quad A \in \mathbb{C}^{m \times m}, B \in \mathbb{C}^{n \times n}, C \in \mathbb{C}^{m \times n}.$$

KEY: Canonical forms for congruence DO not work for $AX + X^T B = C$

- $AX + X^T B = C$, with $A \neq B$.
- $A = Q_A C_A Q_A^T$ and $B = Q_B C_B Q_B^T$.

$$Q_A C_A Q_A^T X + X^T Q_B C_B Q_B^T = C$$

$$C_A Q_A^T X Q_B^{-T} + Q_A^{-1} X^T Q_B C_B = Q_A^{-1} C Q_B^{-T}$$

$$C_A \boxed{Q_A^T X Q_B^{-T}} + \boxed{Q_A^{-1} X^T Q_B} C_B = Q_A^{-1} C Q_B^{-T}$$

- But,

$$\boxed{(Q_A^T X Q_B^{-T})^T} \neq \boxed{Q_A^{-1} X^T Q_B}$$

- with equality only if

$$Q_A = Q_B$$

KEY: Canonical forms for congruence DO not work for $AX + X^T B = C$

- $AX + X^T B = C$, with $A \neq B$.
- $A = Q_A C_A Q_A^T$ and $B = Q_B C_B Q_B^T$.

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The proper transformation for $AX + X^T B = C$

Equivalence of pencil $A - \lambda B^T$

- $AX + X^T B = C$, with $A \neq B$.
- $A - \lambda B^T = PRQ - \lambda PSQ = P(R - \lambda S)Q$, with P and Q nonsingular.

$$PRQX + X^T Q^T S^T P^T = C$$

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- Consistency of the Sylvester equation for $\star$ -congruence
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Motivation for studying $AX + X^*B = C$ (I)

It is well known that given a **block upper triangular matrix** (computed by the **QR-algorithm for eigenvalues**), then

$$\begin{bmatrix} I & X \\ 0 & I \end{bmatrix} \begin{bmatrix} A & C \\ 0 & B \end{bmatrix} \begin{bmatrix} I & X \\ 0 & I \end{bmatrix}^{-1} = \begin{bmatrix} A & C - (AX - XB) \\ 0 & B \end{bmatrix}.$$

Therefore, to find a solution of the **Sylvester equation** $AX - XB = C$ allows us to **block-diagonalize block-triangular matrices** via **similarity**

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This is indeed done in practice in **numerical algorithms** (**LAPACK, MATLAB**) to **compute bases of invariant subspaces** (**eigenvectors**) of matrices, via the classical **Bartels-Stewart algorithm** (**Comm ACM, 1972**) or level-3 BLAS variants of it **Jonsson-Kågström** (**ACM TMS, 2002**).

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Structured numerical algorithms for linear palindromic eigenproblems ($Z + \lambda Z^*$) compute an **anti-triangular Schur form** via unitary $\star$ -congruence:

Theorem (Kressner, Schröder, Watkins (Numer. Alg., 2009) and Mackey², Mehl, Mehrmann (NLAA, 2009))

Let $Z \in \mathbb{C}^{n \times n}$. Then there exists a unitary matrix $U \in \mathbb{C}^{n \times n}$ such that

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M can be computed via **structure-preserving methods** (Kressner, Schröder, Watkins (Numer. Alg., 2009)) or (Mackey², Mehl, Mehrmann (NLAA, 2009)) and **compute eigenvalues of $Z + \lambda Z^*$ with exact pairing $\lambda, 1/\lambda^*$** .

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Given a block upper **ANTI**-triangular matrix (computed via structured algorithms for linear palindromic eigenproblems, when the matrix is real or several eigenvalues form a cluster), then

$$\begin{bmatrix} I & 0 \\ -X & I \end{bmatrix}^* \begin{bmatrix} C & A \\ B & 0 \end{bmatrix} \begin{bmatrix} I & 0 \\ -X & I \end{bmatrix} = \begin{bmatrix} C - (AX + X^*B) & A \\ B & 0 \end{bmatrix}.$$

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Theorem (Wimmer (LAA, 1994), De Terán and D. (ELA, 2011))

Let $\mathbb{F}$ be a field of characteristic different from two and let $A \in \mathbb{F}^{m \times n}$, $B \in \mathbb{F}^{n \times m}$, $C \in \mathbb{F}^{m \times m}$ be given. There is some $X \in \mathbb{F}^{n \times m}$ such that

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Remarks:

- The implication $\Rightarrow$ very easy: done in previous slide.
- The implication $\Leftarrow$ more challenging.
- Wimmer proved in 1994 the result, for $\mathbb{F} = \mathbb{C}$ and $\star = *$, **without any reference to palindromic eigenproblems.**
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Theorem (Roth (Proc. AMS, 1952))

Let $\mathbb{F}$ be any field and let $A \in \mathbb{F}^{m \times m}$, $B \in \mathbb{F}^{n \times n}$, $C \in \mathbb{F}^{m \times n}$ be given. There is some $X \in \mathbb{F}^{m \times n}$ such that

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Uniqueness of solutions of $AX + X^*B = C$ (I)

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- If the matrices $A \in \mathbb{F}^{m \times n}$ and $B \in \mathbb{F}^{n \times m}$ are rectangular ($m \neq n$), then the equation does not have a unique solution for every right-hand side C ,
- that is, the operator

$$\begin{aligned}\mathbb{F}^{n \times m} &\longrightarrow \mathbb{F}^{m \times m} \\ X &\longmapsto AX + X^*B\end{aligned}$$

is never invertible.

- It is of course possible that $m > n$ and that for particular A , B and C , a solution exists and is unique,
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Uniqueness of solutions of $AX + X^*B = C$ (II)

Definition: a set $\{\lambda_1, \dots, \lambda_n\} \subset \mathbb{C}$ is *$\star$ -reciprocal free* if $\lambda_i \neq 1/\lambda_j^*$ for any $1 \leq i, j \leq n$. We admit 0 and/or ∞ as elements of $\{\lambda_1, \dots, \lambda_n\}$.

Theorem (Byers, Kressner (SIMAX, 2006), Kressner, Schröder, Watkins, (Num. Alg., 2009))

Let $A, B \in \mathbb{C}^{n \times n}$ be given. Then:

- $AX + X^T B = C$ has *a unique solution X for every right-hand side $C \in \mathbb{C}^{n \times n}$* if and only if the following conditions hold:
 - 1) The pencil $A - \lambda B^T$ is regular, and
 - 2) the set of eigenvalues of $A - \lambda B^T \setminus \{1\}$ is T -reciprocal free and if 1 is an eigenvalue of $A - \lambda B^T$, then it has algebraic multiplicity 1.
- $AX + X^* B = C$ has *a unique solution X for every right-hand side $C \in \mathbb{C}^{n \times n}$* if and only if the following conditions hold:
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Theorem

Let $A \in \mathbb{C}^{m \times m}$ and $B \in \mathbb{C}^{n \times n}$ be given. Then:

- $AX - XB = C$ has **a unique solution X for every right-hand side $C \in \mathbb{C}^{m \times n}$ if and only if A and B have no eigenvalues in common.**

Remark: Comparison of both results brings to our attention **a key difference** that appears always **between solution methods for $AX + X^*B = C$ and $AX - XB = C$:**

- In $AX + X^*B = C$, one starts by dealing with the eigenproblem of $A - \lambda B^*$, that is, one deals from the very beginning **simultaneously with A and B .**
- By contrast in $AX - XB = C$, one starts by dealing **independently** with the eigenproblems **of A and B .**

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1 Previous and related work

2 The equation $AX^T + XA = 0$

- Motivation: Orbits and the computation of canonical forms
- Strategy for solving $AX^T + XA = 0$
- The canonical form for congruence
- The solution of $AX^T + XA = 0$
- Generic canonical structure for congruence

3 The general equation $AX + X^*B = C$

- Motivation
- Consistency of the Sylvester equation for $\star$ -congruence
- Uniqueness of solutions
- Efficient and stable algorithm to compute unique solutions

4 General solution of $AX + X^*B = 0$

5 Conclusions

The fundamental transformation

- In this section in $AX + X^*B = C$ all matrices are in $\mathbb{C}^{n \times n}$ and the solution is unique for every C .
- $AX + X^*B = C$ is equivalent to a linear system for the n^2 entries of X if $\star = T$ and to a linear system for the $2n^2$ entries of $(\operatorname{Re} X, \operatorname{Im} X)$ if $\star = *$. From now on, we say simply “linear system” for X .
- Then, it is possible to use Gaussian elimination on the equivalent system, but it costs $O(n^6)$ flops, which is not feasible except for small n .
- **IDEA:** transform $AX + X^*B = C$ into an equation of the same type but with much simpler coefficients instead of A and B and that can be easily solved to get a total cost of $O(n^3)$ flops.
- To this purpose, use **QZ algorithm** to compute in $O(n^3)$ flops the **generalized Schur decomposition** of

$$A - \lambda B^* = U(R - \lambda S)V, \quad \text{where} \quad \begin{cases} R, S & \text{are upper triangular} \\ U, V & \text{are unitary matrices} \end{cases}$$

If A, B real matrices: use real arithmetic to get *quasi-triangular* R . We do not consider this for brevity.

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Algorithm to solve $AX + X^*B = C$ in $O(n^3)$ flops

INPUT: $A, B, C \in \mathbb{C}^{n \times n}$

OUTPUT: $X \in \mathbb{C}^{n \times n}$

Step 1. Compute via QZ algorithm R, S, U and V such that

$$A = URV, \quad B^* = USV, \quad \text{where } \begin{cases} R, S & \text{are upper triangular} \\ U, V & \text{are unitary matrices} \end{cases}$$

Step 2. Compute $E = U^* C (U^*)^*$

Step 3. Solve for $W \in \mathbb{C}^{n \times n}$ the transformed equation

$$RW + W^* S^* = E$$

Step 4. Compute $X = V^* W U^*$

Algorithm to solve $AX + X^* B = C$ in $O(n^3)$ flops

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Step 3. How to solve for $W \in \mathbb{C}^{n \times n}$ the transformed equation

$$RW + W^* S^* = E \quad ?$$

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Algorithm to solve the transformed equation $RW + W^* S^* = E$ (I)

We illustrate with 4×4 example for simplicity:

$$\begin{aligned}
 & \begin{bmatrix} r_{11} & r_{12} & r_{13} & r_{14} \\ 0 & r_{22} & r_{23} & r_{24} \\ 0 & 0 & r_{33} & r_{34} \\ 0 & 0 & 0 & r_{44} \end{bmatrix} \begin{bmatrix} w_{11} & w_{12} & w_{13} & w_{14} \\ w_{21} & w_{22} & w_{23} & w_{24} \\ w_{31} & w_{32} & w_{33} & w_{34} \\ w_{41} & w_{42} & w_{43} & w_{44} \end{bmatrix} \\
 & + \begin{bmatrix} w_{11}^* & w_{21}^* & w_{31}^* & w_{41}^* \\ w_{12}^* & w_{22}^* & w_{32}^* & w_{42}^* \\ w_{13}^* & w_{23}^* & w_{33}^* & w_{43}^* \\ w_{14}^* & w_{24}^* & w_{34}^* & w_{44}^* \end{bmatrix} \begin{bmatrix} s_{11}^* & 0 & 0 & 0 \\ s_{12}^* & s_{22}^* & 0 & 0 \\ s_{13}^* & s_{23}^* & s_{33}^* & 0 \\ s_{14}^* & s_{24}^* & s_{34}^* & s_{44}^* \end{bmatrix} = \begin{bmatrix} e_{11} & e_{12} & e_{13} & e_{14} \\ e_{21} & e_{22} & e_{23} & e_{24} \\ e_{31} & e_{32} & e_{33} & e_{34} \\ e_{41} & e_{42} & e_{43} & e_{44} \end{bmatrix}
 \end{aligned}$$

If we equate the (4,4)-entry, then we get

$$r_{44} w_{44} + w_{44}^* s_{44}^* = e_{44} ,$$

a scalar equation that allows us to determine w_{44} .

Algorithm to solve the transformed equation $RW + W^* S^* = E$ (I)

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 \end{aligned}$$

If we equate the (3,4) and (4,3) entries, then we get

$$\begin{aligned}
 s_{33} w_{34} + w_{43}^* r_{44}^* &= e_{43}^* - s_{34} w_{44} \\
 r_{33} w_{34} + w_{43}^* s_{44}^* &= e_{34} - r_{34} w_{44}
 \end{aligned}$$

a 2×2 system of scalar equations that allows us to determine w_{34} and w_{43} **simultaneously**.

Algorithm to solve the transformed equation $RW + W^* S^* = E$ (I)

We illustrate with 4×4 example for simplicity:

$$\begin{aligned}
 & \begin{bmatrix} r_{11} & r_{12} & r_{13} & r_{14} \\ 0 & r_{22} & r_{23} & r_{24} \\ 0 & 0 & r_{33} & r_{34} \\ 0 & 0 & 0 & r_{44} \end{bmatrix} \begin{bmatrix} w_{11} & w_{12} & w_{13} & w_{14} \\ w_{21} & w_{22} & w_{23} & w_{24} \\ w_{31} & w_{32} & w_{33} & w_{34} \\ w_{41} & w_{42} & w_{43} & w_{44} \end{bmatrix} \\
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If we equate the (3,4) and (4,3) entries, then we get

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a 2×2 system of scalar equations that allows us to determine w_{34} and w_{43} **simultaneously**.

Algorithm to solve the transformed equation $RW + W^* S^* = E$ (I)

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 \end{aligned}$$

If we equate the (2,4) and (4,2) entries, then we get

$$\begin{aligned}
 s_{22} w_{24} + w_{42}^* r_{44}^* &= e_{42}^* - s_{23} w_{34} - s_{24} w_{44} \\
 r_{22} w_{24} + w_{42}^* s_{44}^* &= e_{24} - r_{23} w_{34} - r_{24} w_{44}
 \end{aligned}$$

a 2×2 system of scalar equations that allows us to determine w_{24} and w_{42} **simultaneously**.

Algorithm to solve the transformed equation $RW + W^* S^* = E$ (I)

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$$\begin{aligned}
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 + & \begin{bmatrix} w_{11}^* & w_{21}^* & w_{31}^* & w_{41}^* \\ w_{12}^* & w_{22}^* & w_{32}^* & w_{42}^* \\ w_{13}^* & w_{23}^* & w_{33}^* & w_{43}^* \\ w_{14}^* & w_{24}^* & w_{34}^* & w_{44}^* \end{bmatrix} \begin{bmatrix} s_{11}^* & 0 & 0 & 0 \\ s_{12}^* & s_{22}^* & 0 & 0 \\ s_{13}^* & s_{23}^* & s_{33}^* & 0 \\ s_{14}^* & s_{24}^* & s_{34}^* & s_{44}^* \end{bmatrix} = \begin{bmatrix} e_{11} & e_{12} & e_{13} & e_{14} \\ e_{21} & e_{22} & e_{23} & e_{24} \\ e_{31} & e_{32} & e_{33} & e_{34} \\ e_{41} & e_{42} & e_{43} & e_{44} \end{bmatrix}
 \end{aligned}$$

If we equate the (2,4) and (4,2) entries, then we get

$$\begin{aligned}
 s_{22} w_{24} + w_{42}^* r_{44}^* &= e_{42}^* - s_{23} w_{34} - s_{24} w_{44} \\
 r_{22} w_{24} + w_{42}^* s_{44}^* &= e_{24} - r_{23} w_{34} - r_{24} w_{44}
 \end{aligned}$$

a 2×2 system of scalar equations that allows us to determine w_{24} and w_{42} **simultaneously**.

Algorithm to solve the transformed equation $RW + W^* S^* = E$ (I)

We illustrate with 4×4 example for simplicity:

$$\begin{bmatrix} r_{11} & r_{12} & r_{13} & r_{14} \\ 0 & r_{22} & r_{23} & r_{24} \\ 0 & 0 & r_{33} & r_{34} \\ 0 & 0 & 0 & r_{44} \end{bmatrix} \begin{bmatrix} w_{11} & w_{12} & w_{13} & w_{14} \\ w_{21} & w_{22} & w_{23} & w_{24} \\ w_{31} & w_{32} & w_{33} & w_{34} \\ w_{41} & w_{42} & w_{43} & w_{44} \end{bmatrix} + \begin{bmatrix} w_{11}^* & w_{21}^* & w_{31}^* & w_{41}^* \\ w_{12}^* & w_{22}^* & w_{32}^* & w_{42}^* \\ w_{13}^* & w_{23}^* & w_{33}^* & w_{43}^* \\ w_{14}^* & w_{24}^* & w_{34}^* & w_{44}^* \end{bmatrix} \begin{bmatrix} s_{11}^* & 0 & 0 & 0 \\ s_{12}^* & s_{22}^* & 0 & 0 \\ s_{13}^* & s_{23}^* & s_{33}^* & 0 \\ s_{14}^* & s_{24}^* & s_{34}^* & s_{44}^* \end{bmatrix} = \begin{bmatrix} e_{11} & e_{12} & e_{13} & e_{14} \\ e_{21} & e_{22} & e_{23} & e_{24} \\ e_{31} & e_{32} & e_{33} & e_{34} \\ e_{41} & e_{42} & e_{43} & e_{44} \end{bmatrix}$$

If we equate the (1,4) and (4,1) entries, then we get

$$\begin{aligned} s_{11} w_{14} + w_{41}^* r_{44}^* &= e_{41}^* - s_{12} w_{24} - s_{13} w_{34} - s_{14} w_{44} \\ r_{11} w_{14} + w_{41}^* s_{44}^* &= e_{14} - r_{12} w_{24} - r_{13} w_{34} - r_{14} w_{44} \end{aligned}$$

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If we equate the (1,4) and (4,1) entries, then we get

$$\begin{aligned} s_{11} w_{14} + w_{41}^* r_{44} &= e_{41} - s_{12} w_{24} - s_{13} w_{34} - s_{14} w_{44} \\ r_{11} w_{14} + w_{41}^* s_{44} &= e_{14} - r_{12} w_{24} - r_{13} w_{34} - r_{14} w_{44} \end{aligned}$$

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If we equate the (1:3,1:3) submatrices , then we get

$$\begin{bmatrix} r_{11} & r_{12} & r_{13} \\ 0 & r_{22} & r_{23} \\ 0 & 0 & r_{33} \end{bmatrix} \begin{bmatrix} w_{11} & w_{12} & w_{13} \\ w_{21} & w_{22} & w_{23} \\ w_{31} & w_{32} & w_{33} \end{bmatrix} + \begin{bmatrix} w_{11}^* & w_{21}^* & w_{31}^* \\ w_{12}^* & w_{22}^* & w_{32}^* \\ w_{13}^* & w_{23}^* & w_{33}^* \end{bmatrix} \begin{bmatrix} s_{11}^* & 0 & 0 \\ s_{12}^* & s_{22}^* & 0 \\ s_{13}^* & s_{23}^* & s_{33}^* \end{bmatrix} = \begin{bmatrix} e_{11} & e_{12} & e_{13} \\ e_{21} & e_{22} & e_{23} \\ e_{31} & e_{32} & e_{33} \end{bmatrix} - \begin{bmatrix} r_{14} \\ r_{24} \\ r_{34} \end{bmatrix} \begin{bmatrix} w_{41} & w_{42} & w_{43} \end{bmatrix} - \begin{bmatrix} w_{41}^* \\ w_{42}^* \\ w_{43}^* \end{bmatrix} \begin{bmatrix} s_{14}^* & s_{24}^* & s_{34}^* \end{bmatrix}$$

which is a 3×3 matrix equation of the same type as the original one.

Algorithm to solve the transformed equation $RW + W^* S^* = E$ (I)

We illustrate with 4×4 example for simplicity:

$$\begin{bmatrix} r_{11} & r_{12} & r_{13} & r_{14} \\ 0 & r_{22} & r_{23} & r_{24} \\ 0 & 0 & r_{33} & r_{34} \\ 0 & 0 & 0 & r_{44} \end{bmatrix} \begin{bmatrix} w_{11} & w_{12} & w_{13} & w_{14} \\ w_{21} & w_{22} & w_{23} & w_{24} \\ w_{31} & w_{32} & w_{33} & w_{34} \\ w_{41} & w_{42} & w_{43} & w_{44} \end{bmatrix} + \begin{bmatrix} w_{11}^* & w_{21}^* & w_{31}^* & w_{41}^* \\ w_{12}^* & w_{22}^* & w_{32}^* & w_{42}^* \\ w_{13}^* & w_{23}^* & w_{33}^* & w_{43}^* \\ w_{14}^* & w_{24}^* & w_{34}^* & w_{44}^* \end{bmatrix} \begin{bmatrix} s_{11}^* & 0 & 0 & 0 \\ s_{12}^* & s_{22}^* & 0 & 0 \\ s_{13}^* & s_{23}^* & s_{33}^* & 0 \\ s_{14}^* & s_{24}^* & s_{34}^* & s_{44}^* \end{bmatrix} = \begin{bmatrix} e_{11} & e_{12} & e_{13} & e_{14} \\ e_{21} & e_{22} & e_{23} & e_{24} \\ e_{31} & e_{32} & e_{33} & e_{34} \\ e_{41} & e_{42} & e_{43} & e_{44} \end{bmatrix}$$

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Remarks on algorithm to solve $AX + X^*B = C$

- **Cost:** $2n^3 + O(n^2)$ flops for simplified system and a total cost $76n^3 + O(n^2)$ flops for the whole algorithm for $AX + X^*B = C$.
- Forward stable algorithm.
- The algorithm should be compared with Bartels-Stewart algorithm for Sylvester equation $AX - XB = C$:
 - 1 Compute independently triang. Schur forms T_A and T_B of A and B .
 - 2 Solve $T_A Y - Y T_B = D$ for Y .
 - 3 Recover X from Y .
- Same flavor, but also relevant differences.

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- Consistency of the Sylvester equation for $\star$ -congruence
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5 Conclusions

Theoretical method to solve $AX + X^*B = 0$ (I)

- In case of consistency, but “nonuniqueness”, general solution of $AX + X^*B = C$ is $X = X_p + X_h$, where
 - 1 X_p is a particular solution and
 - 2 X_h is the general solution of $AX + X^*B = 0$.

The latter found by De Terán, D., Guillery, Montealegre, Reyes, Lin. Alg. Appl., 2013

- Summer REU program 2011, U. of California at S. Barbara, M.I. Bueno.
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Theoretical method to solve $AX + X^*B = 0$ (II)

- **KEY IDEA:** If $E - \lambda F^*$ is the *Kronecker Canonical form (KCF)* of $A - \lambda B^*$ for strict equivalence, i.e.,

$$A - \lambda B^* = P(E - \lambda F^*)Q, \quad \text{with } P \text{ and } Q \text{ nonsingular,}$$

then $AX + X^*B = 0$ can be transformed into

$$EY + Y^*F = 0, \quad \text{with } Y = QXP^{-*}.$$

- If $E = E_1 \oplus \cdots \oplus E_d$, $F^* = F_1^* \oplus \cdots \oplus F_d^*$, and $Y = [Y_{ij}]$ is partitioned into blocks accordingly, then this equation decouples in

$$E_i Y_{ii} + Y_{ii}^* F_i = 0 \quad \text{and} \quad \begin{cases} E_i Y_{ij} + Y_{ji}^* F_j = 0 \\ E_j Y_{ji} + Y_{ij}^* F_i = 0 \end{cases}, \quad (1 \leq i < j \leq d).$$

- Since KCF has 4 types of blocks, this produces **14** different types of matrix (systems) equations, whose explicit solutions have been found.
- Much more complicated general solution than standard Sylvester eq: $AX - XB = 0$, which depends on JCF of A and B and requires to solve only one type of equation.

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The Kronecker Canonical Form of a Matrix Pencil

Theorem

Let $G, H \in \mathbb{C}^{m \times n}$. Then $G - \lambda H$ is strictly equivalent to a direct sum of pencils of the following types

“Finite blocks”: $J_k(\lambda_i - \lambda) := \begin{bmatrix} \lambda_i - \lambda & 1 & & \\ & \lambda_i - \lambda & 1 & \\ & & \ddots & \ddots \\ & & & \lambda_i - \lambda \end{bmatrix}$ are $k \times k$.

“Infinite blocks”: $N_\ell = \begin{bmatrix} 1 & \lambda & & \\ & 1 & \lambda & \\ & & \ddots & \ddots \\ & & & 1 \end{bmatrix}$ are $\ell \times \ell$.

“Right singular blocks”: $L_p := \begin{bmatrix} \lambda & 1 & & \\ & \lambda & 1 & \\ & & \ddots & \ddots \\ & & & \lambda & 1 \end{bmatrix}$ are $p \times (p + 1)$.

“Left singular blocks”: *transposes of right singular blocks.*

Theorem

Let $A \in \mathbb{C}^{m \times n}$ and $B \in \mathbb{C}^{n \times m}$. If the pencil $A - \lambda B^T$ has the KCF

$$\begin{aligned} E - \lambda F^T &= L_{\varepsilon_1} \oplus L_{\varepsilon_2} \oplus \cdots \oplus L_{\varepsilon_a} \\ &\quad \oplus L_{\eta_1}^T \oplus L_{\eta_2}^T \oplus \cdots \oplus L_{\eta_b}^T \\ &\quad \oplus N_{u_1} \oplus N_{u_2} \oplus \cdots \oplus N_{u_c} \\ &\quad \oplus J_{k_1}(\lambda_1 - \lambda) \oplus J_{k_2}(\lambda_2 - \lambda) \oplus \cdots \oplus J_{k_d}(\lambda_d - \lambda). \end{aligned}$$

Then the dimension of the solution space of the matrix equation

$$AX + X^T B = 0$$

depends only on $E - \lambda F^T$ and is

Theorem

$$\begin{aligned}
 \text{dimension} = & \sum_{i=1}^a \varepsilon_i + \sum_{\lambda_i=1} [k_i/2] + \sum_{\lambda_j=-1} [k_j/2] \\
 & + \sum_{\substack{i,j=1 \\ i < j}}^a (\varepsilon_i + \varepsilon_j) + \sum_{\substack{i < j \\ \lambda_i \lambda_j = 1}} \min\{k_i, k_j\} \\
 & + \sum_{\varepsilon_i \leq \eta_j} (\eta_j - \varepsilon_i + 1) \\
 & + a \sum_{i=1}^c u_i + a \sum_{i=1}^d k_i + \sum_{\substack{i,j \\ \lambda_j=0}} \min\{u_i, k_j\}
 \end{aligned}$$

- The method in this section can be applied when $A = B$ (orbits).
- Same results are obtained but expressed in different ways.
- What method is better?

Last important comment on solution of $AX + X^T B = 0$

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- Many questions related to the Sylvester equation for $\star$ -congruence $AX + X^\star B = C$ are nowadays well-understood.
- This equation appears in several applications and is related to “congruence problems”.
- Connections with classical Sylvester equation $AX - XB = C$ but also relevant differences.
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